

# BGP Techniques for Internet Service Providers

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### **Presentation Slides**

- Are available on
  - ftp://ftp-eng.cisco.com
  - /pfs/seminars/NANOG40-BGP-Techniques.pdf
  - And on the NANOG website
- Please feel free to ask questions any time

### **BGP Techniques for Internet Service Providers**

- BGP Basics
- Scaling BGP
- Using Communities
- Deploying BGP in an ISP network



# **BGP Basics**

What is **BGP**?

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### **Border Gateway Protocol**

 A Routing Protocol used to exchange routing information between different networks

**Exterior gateway protocol** 

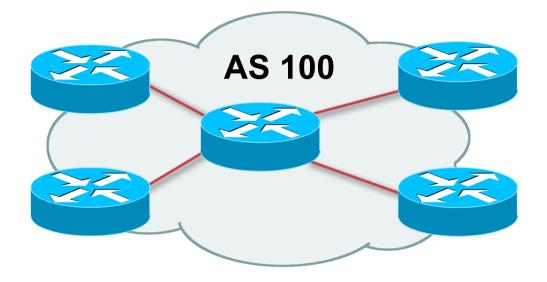
Described in RFC4271

RFC4276 gives an implementation report on BGP RFC4277 describes operational experiences using BGP

 The Autonomous System is BGP's fundamental operating unit

It is used to uniquely identify networks with a common routing policy

# Autonomous System (AS)



- Collection of networks with same routing policy
- Single routing protocol
- Usually under single ownership, trust and administrative control
- Identified by a unique number (ASN)

# Autonomous System Number (ASN)

#### • An ASN is a 16 bit integer

1-64511 are for use on the public Internet64512-65534 are for private use only0 and 65535 are reserved

#### ASNs are now extended to 32 bit!

**RFC4893 is standards document describing 32-bit ASNs** 

**Representation still under discussion:** 

32-bit notation or "16.16" notation

Now expired Internet Draft:

draft-michaelson-4byte-as-representation-02.txt

AS 23456 is used to represent 32-bit ASNs in 16-bit ASN world

# Autonomous System Number (ASN)

 ASNs are distributed by the Regional Internet Registries

They are also available from upstream ISPs who are members of one of the RIRs

 Current 16-bit ASN allocations up to 44031 have been made to the RIRs

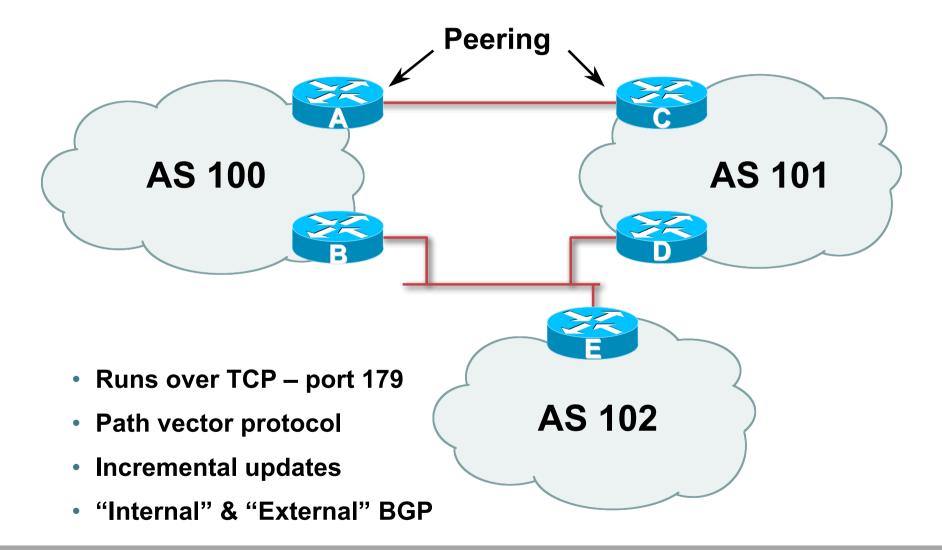
Around 25200 are visible on the Internet

The RIRs also have received 1024 32-bit ASNs each

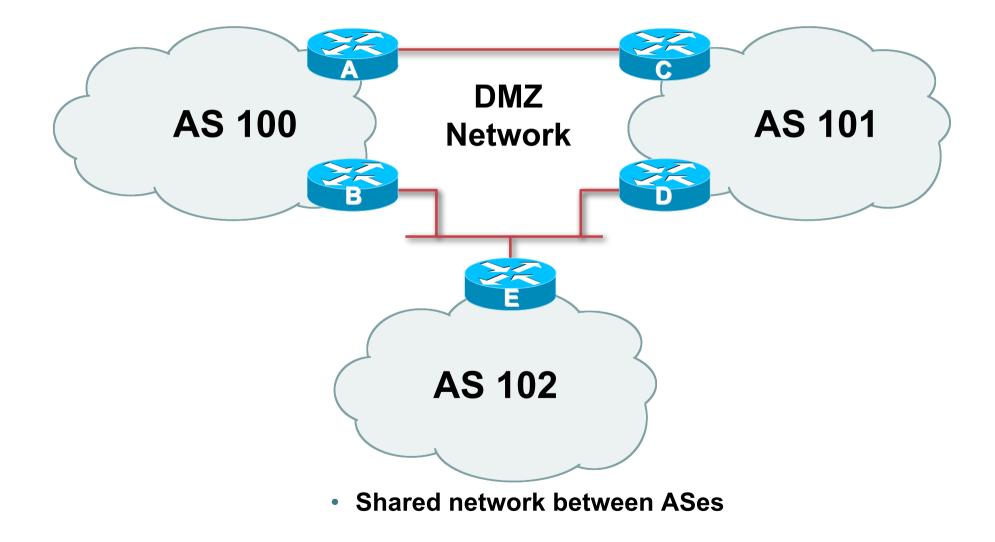
Around 5 are visible on the Internet (early adopters)

See www.iana.org/assignments/as-numbers

#### **BGP** Basics



# **Demarcation Zone (DMZ)**



# **BGP General Operation**

- Learns multiple paths via internal and external BGP speakers
- Picks the best path and installs in the forwarding table
- Best path is sent to external BGP neighbours
- Policies are applied by influencing the best path selection

### eBGP & iBGP

BGP used internally (iBGP) and externally (eBGP)

#### iBGP used to carry

some/all Internet prefixes across ISP backbone

**ISP's customer prefixes** 

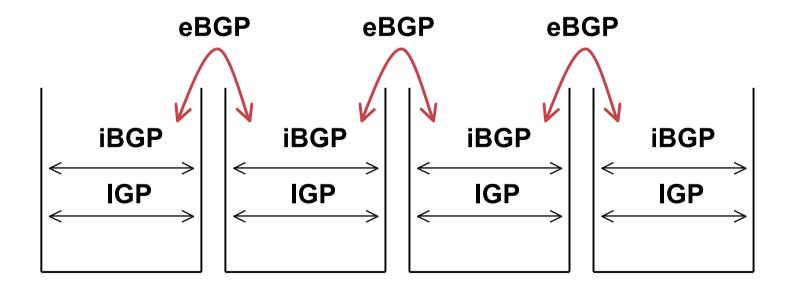
#### eBGP used to

exchange prefixes with other ASes

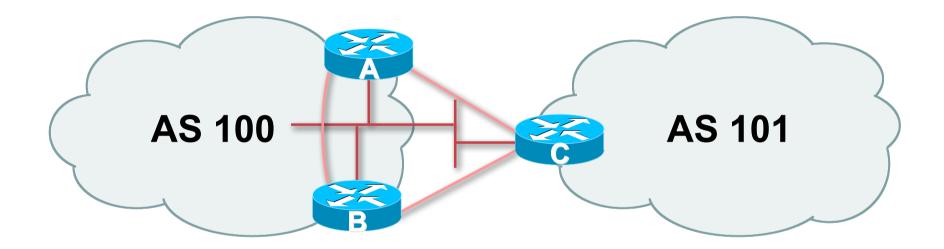
implement routing policy

# **BGP/IGP model used in ISP networks**

Model representation



### **External BGP Peering (eBGP)**



- Between BGP speakers in different AS
- Should be directly connected
- Never run an IGP between eBGP peers

# Internal BGP (iBGP)

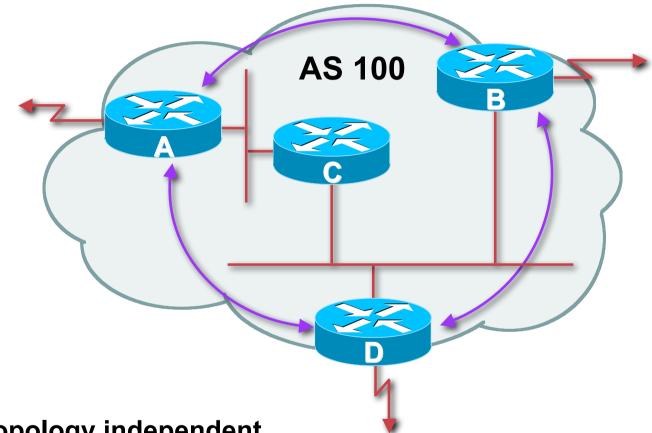
- BGP peer within the same AS
- Not required to be directly connected
   IGP takes care of inter-BGP speaker connectivity
- iBGP speakers must to be fully meshed:

They originate connected networks

They pass on prefixes learned from outside the ASN

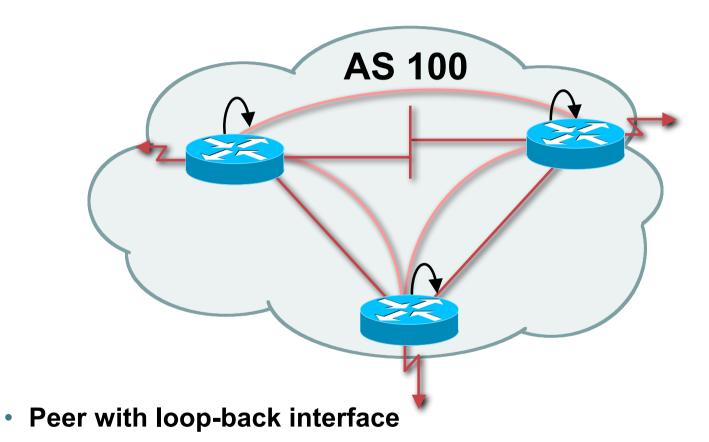
They do not pass on prefixes learned from other iBGP speakers

# Internal BGP Peering (iBGP)



- Topology independent
- Each iBGP speaker must peer with every other iBGP speaker in the AS

# **Peering to Loopback Interfaces**



Loop-back interface does not go down – ever!

• Do not want iBGP session to depend on state of a single interface or the physical topology



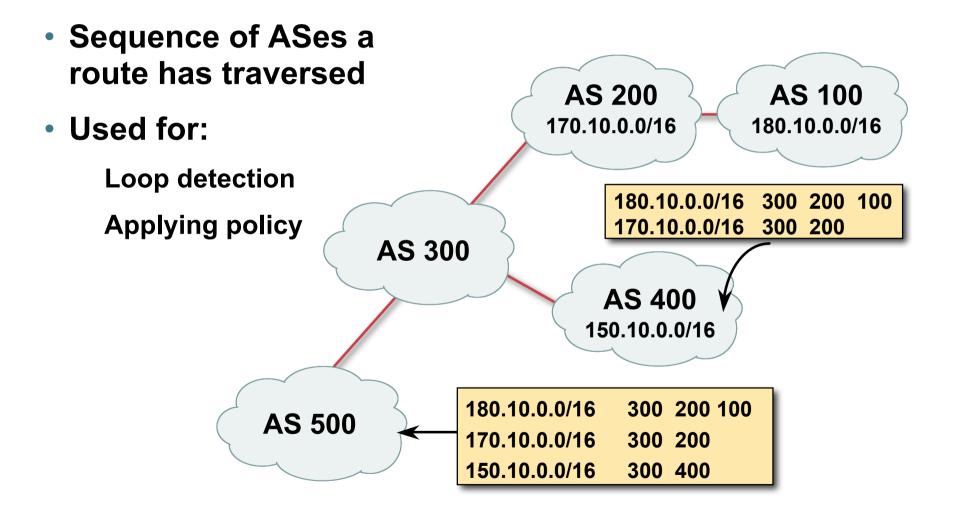
# **BGP** Attributes

#### **Information about BGP**

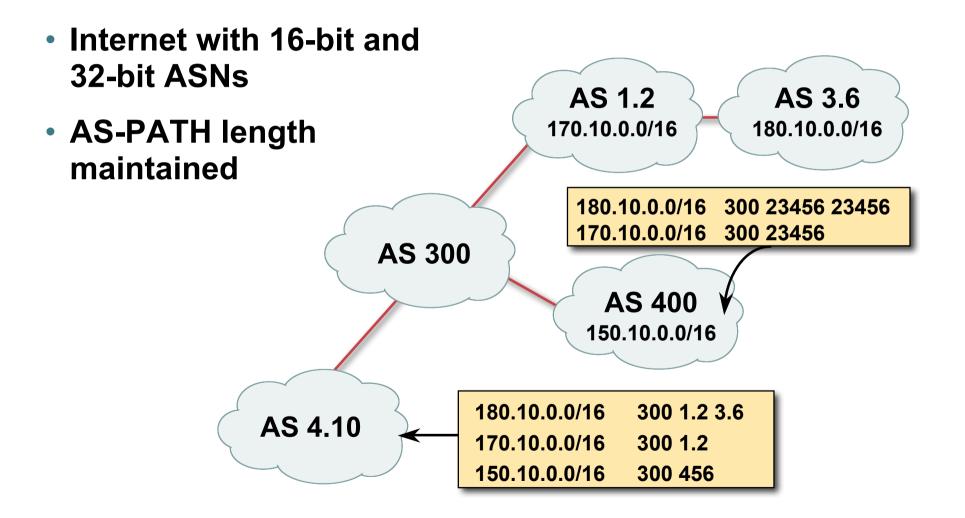
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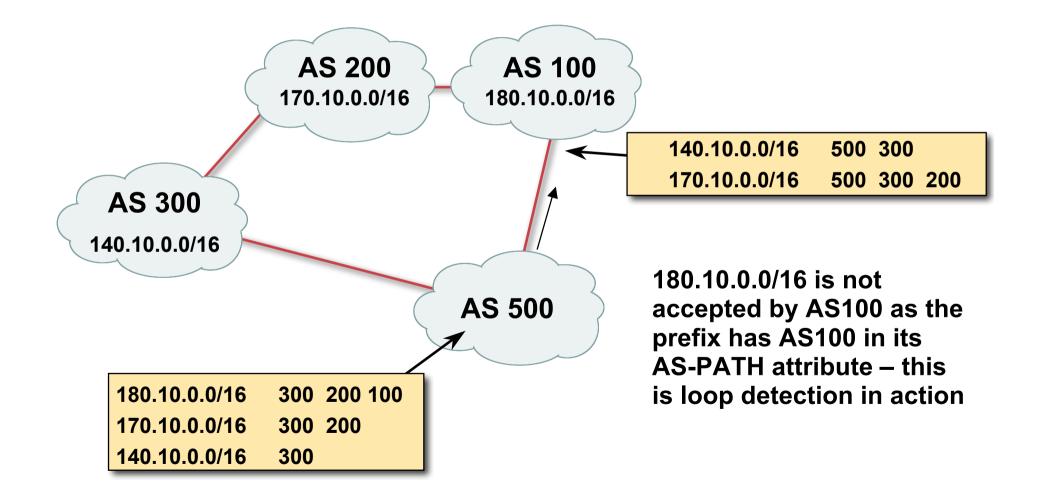
#### **AS-Path**



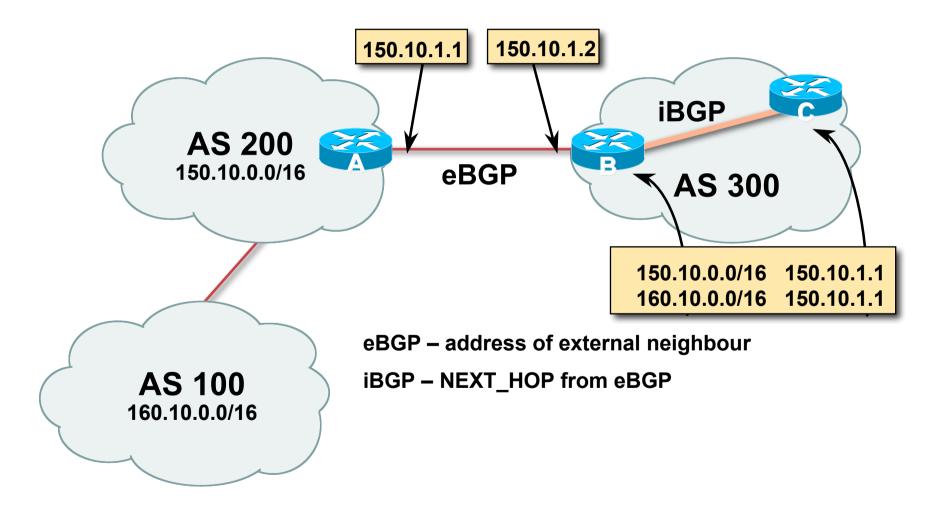
## AS-Path (with 16 and 32-bit ASNs)



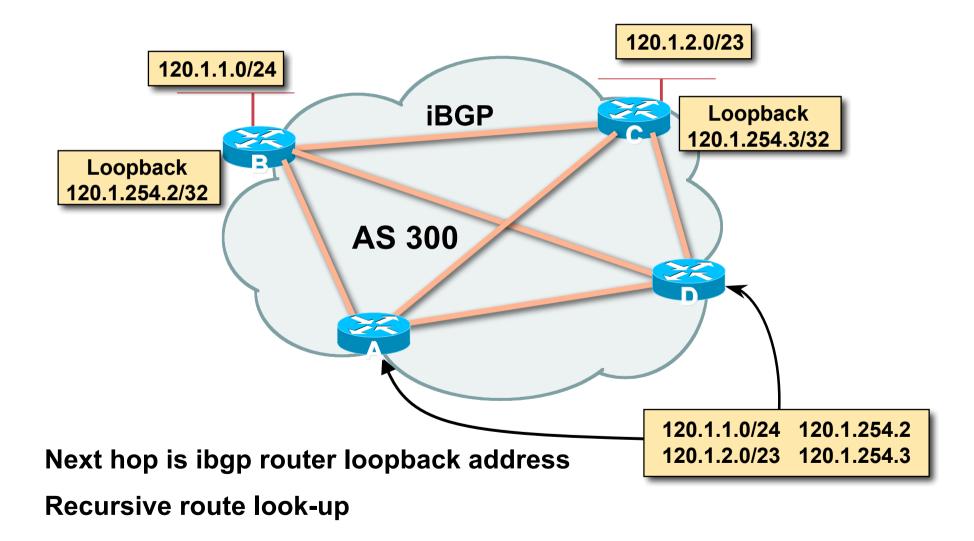
### **AS-Path loop detection**



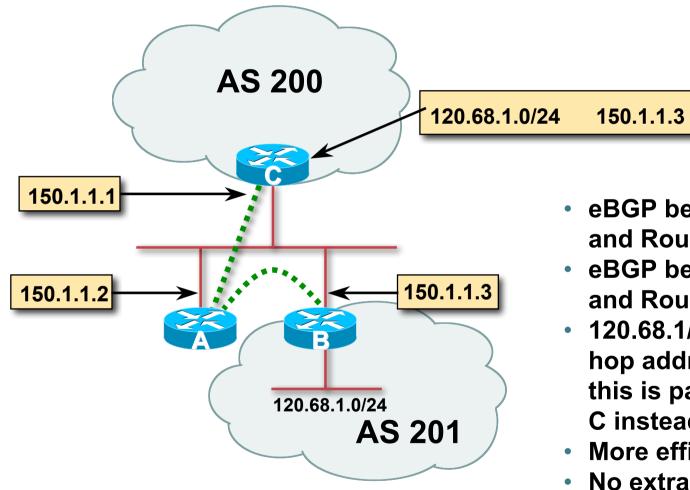
# **Next Hop**



### **iBGP** Next Hop



# **Third Party Next Hop**



- eBGP between Router A and Router C
- eBGP between Router A and Router B
- 120.68.1/24 prefix has next hop address of 150.1.1.3 this is passed on to Router C instead of 150.1.1.2
- More efficient
- No extra config needed

# Next Hop (Summary)

- IGP should carry route to next hops
- Recursive route look-up
- Unlinks BGP from actual physical topology
- Allows IGP to make intelligent forwarding decision

# Origin

- Conveys the origin of the prefix
- Historical attribute

Used in transition from EGP to BGP

- Influences best path selection
- Three values: IGP, EGP, incomplete

**IGP** – generated by BGP network statement

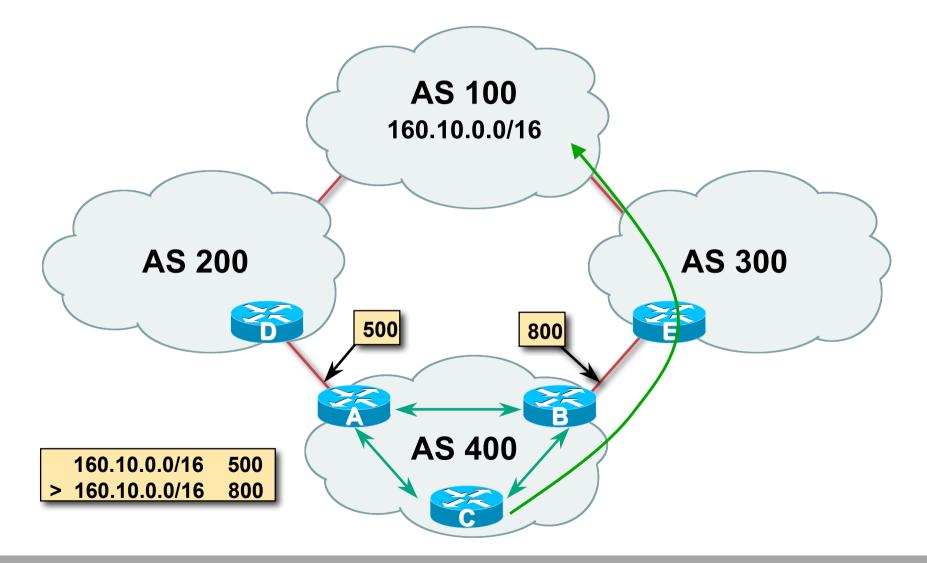
EGP – generated by EGP

incomplete – redistributed from another routing protocol

# Aggregator

- Conveys the IP address of the router or BGP speaker generating the aggregate route
- Useful for debugging purposes
- Does not influence best path selection

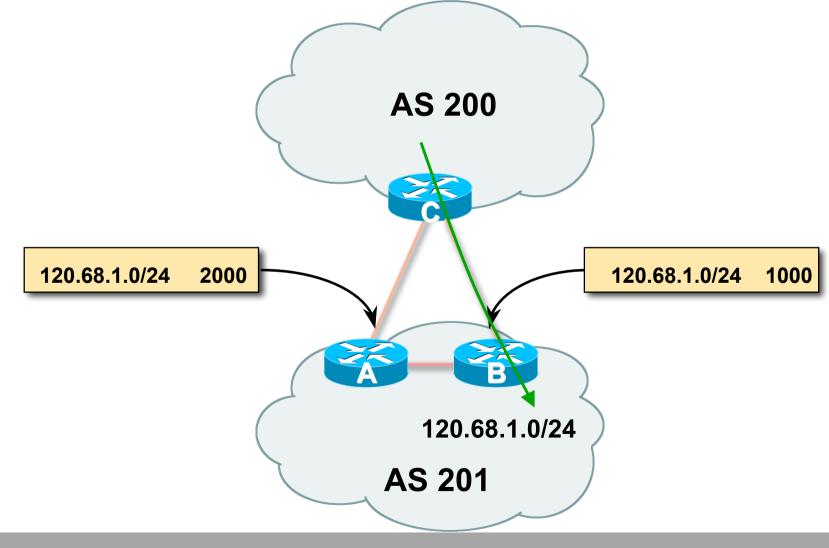
# Local Preference



#### **Local Preference**

- Local to an AS non-transitive
   Default local preference is 100 (IOS)
- Used to influence BGP path selection
   determines best path for *outbound* traffic
- Path with highest local preference wins

# Multi-Exit Discriminator (MED)



# **Multi-Exit Discriminator**

- Inter-AS non-transitive & optional attribute
- Used to convey the relative preference of entry points

determines best path for *inbound* traffic

Comparable if paths are from same AS

bgp always-compare-med allows comparisons of MEDs from different ASes

- Path with lowest MED wins
- Absence of MED attribute implies MED value of zero (RFC4271)

#### Multi-Exit Discriminator "metric confusion"

- MED is non-transitive and optional attribute
  - Some implementations send learned MEDs to iBGP peers by default, others do not
  - Some implementations send MEDs to eBGP peers by default, others do not
- Default metric varies according to vendor implementation

**Original BGP spec (RFC1771) made no recommendation** 

Some implementations said that absence of metric was equivalent to 0

Other implementations said that absence of metric was equivalent to 2<sup>32</sup>-1 (highest possible) or 2<sup>32</sup>-2

Potential for "metric confusion"

### Community

Communities are described in RFC1997

**Transitive and Optional Attribute** 

• 32 bit integer

**Represented as two 16 bit integers (RFC1998)** 

Common format is <*local-ASN*>:xx

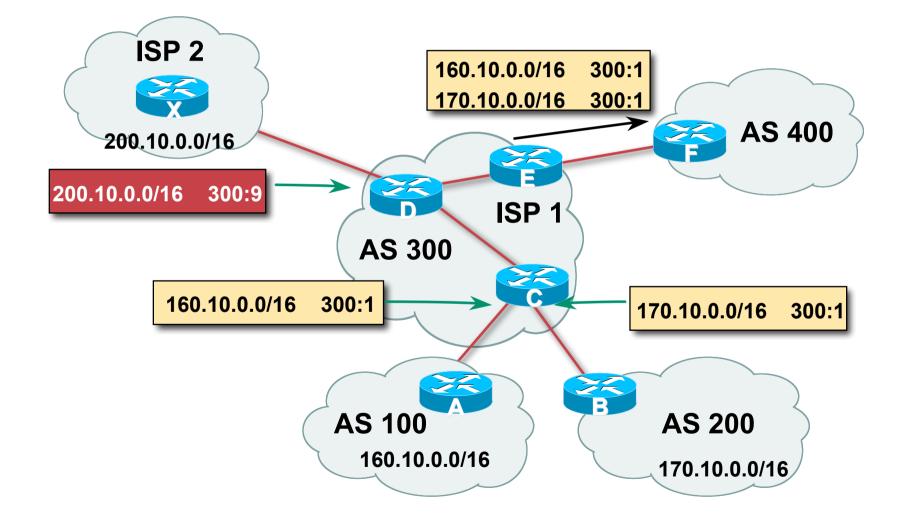
0:0 to 0:65535 and 65535:0 to 65535:65535 are reserved

Used to group destinations

Each destination could be member of multiple communities

 Very useful in applying policies within and between ASes

### Community



# **Well-Known Communities**

#### Several well known communities

www.iana.org/assignments/bgp-well-known-communities

no-export 65535:65281

do not advertise to any eBGP peers

no-advertise
 <u>65535:65282</u>

do not advertise to any BGP peer

no-export-subconfed
 65535:65283

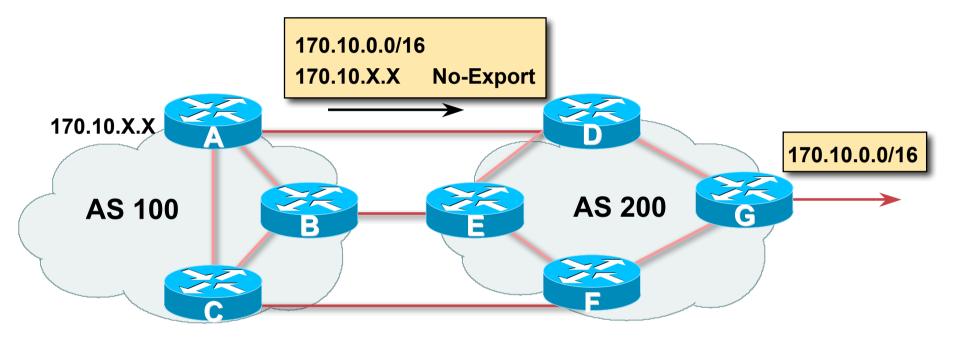
do not advertise outside local AS (only used with confederations)

• no-peer

65535:65284

do not advertise to bi-lateral peers (RFC3765)

# **No-Export Community**

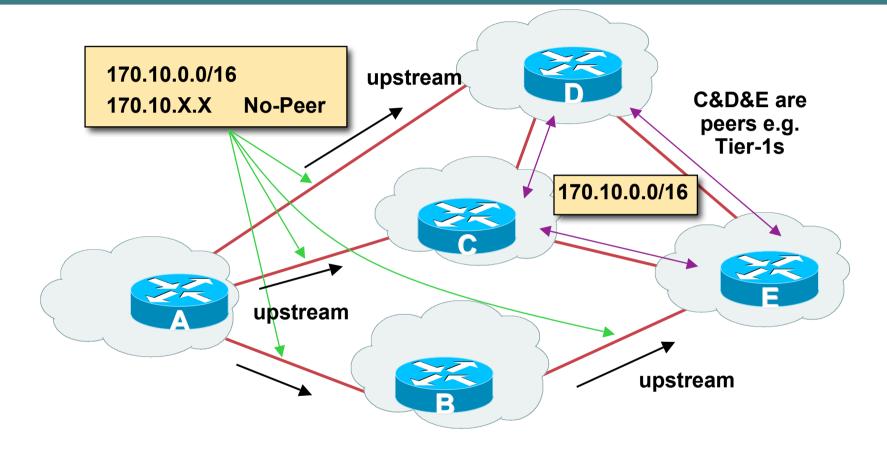


AS100 announces aggregate and subprefixes

aim is to improve loadsharing by leaking subprefixes

- Subprefixes marked with no-export community
- Router G in AS200 does not announce prefixes with no-export community set

## **No-Peer Community**



 Sub-prefixes marked with no-peer community are not sent to bilateral peers

They are only sent to upstream providers

### Community Implementation details

#### Community is an optional attribute

Some implementations send communities to iBGP peers by default, some do not

Some implementations send communities to eBGP peers by default, some do not

#### Being careless can lead to community "confusion"

ISPs need consistent community policy within their own networks

And they need to inform peers, upstreams and customers about their community expectations



# **BGP Path Selection Algorithm**

Why Is This the Best Path?

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# BGP Path Selection Algorithm for IOS Part One

- Do not consider path if no route to next hop
- Do not consider iBGP path if not synchronised (Cisco IOS only)
- Highest weight (local to router)
- Highest local preference (global within AS)
- Prefer locally originated route
- Shortest AS path

# BGP Path Selection Algorithm for IOS Part Two

Lowest origin code

**IGP < EGP < incomplete** 

Lowest Multi-Exit Discriminator (MED)

If bgp deterministic-med, order the paths before comparing

(BGP spec does not specify in which order the paths should be compared. This means best path depends on order in which the paths are compared.)

If bgp always-compare-med, then compare for all paths

otherwise MED only considered if paths are from the same AS (default)

# BGP Path Selection Algorithm for IOS Part Three

- Prefer eBGP path over iBGP path
- Path with lowest IGP metric to next-hop
- Lowest router-id (originator-id for reflected routes)
- Shortest Cluster-List

**Client must be aware of Route Reflector attributes!** 

Lowest neighbour IP address

# **BGP Path Selection Algorithm**

#### • In multi-vendor environments:

Make sure the path selection processes are understood for each brand of equipment

Each vendor has slightly different implementations, extra steps, extra features, etc

Watch out for possible MED confusion



# **Applying Policy with BGP**

#### **Controlling Traffic Flow & Traffic Engineering**

# Applying Policy in BGP: Why?

- Network operators rarely "plug in routers and go"
- External relationships:
  - Control who they peer with
  - Control who they give transit to
  - Control who they get transit from
- Traffic flow control:
  - Efficiently use the scarce infrastructure resources (external link load balancing)
  - **Congestion avoidance**
  - **Terminology: Traffic Engineering**

### Applying Policy in BGP: How?

- Policies are applied by:
  - Setting BGP attributes (local-pref, MED, AS-PATH, community), thereby influencing the path selection process
  - **Advertising or Filtering prefixes**
  - Advertising or Filtering prefixes according to ASN and AS-PATHs
  - Advertising or Filtering prefixes according to Community membership

### Applying Policy with BGP: Tools

- Most implementations have tools to apply policies to BGP:
  - Prefix manipulation/filtering
  - **AS-PATH** manipulation/filtering
  - **Community Attribute setting and matching**
- Implementations also have policy language which can do various match/set constructs on the attributes of chosen BGP routes



**Extending BGP** 

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- Documented in RFC2842
- Capabilities parameters passed in BGP open message
- Unknown or unsupported capabilities will result in NOTIFICATION message
- Codes:

0 to 63 are assigned by IANA by IETF consensus 64 to 127 are assigned by IANA "first come first served" 128 to 255 are vendor specific

#### **Current capabilities are:**

0	Reserved	[RFC3392]
1	Multiprotocol Extensions for BGP-4	[RFC4760]
2	Route Refresh Capability for BGP-4	[RFC2918]
3	Cooperative Route Filtering Capability	[ID]
4	Multiple routes to a destination capability	[RFC3107]
64	Graceful Restart Capability	[RFC4724]
65	Support for 4 octet ASNs	[RFC4893]
66	Deprecated 2003-03-06	
67	Support for Dynamic Capability	[ID]
68	Multisession BGP	[ID]

See www.iana.org/assignments/capability-codes

Multiprotocol extensions

This is a whole different world, allowing BGP to support more than IPv4 unicast routes

**Examples include: v4 multicast, IPv6, v6 multicast, VPNs** 

Another tutorial (or many!)

- Route refresh is a well known scaling technique covered shortly
- 32-bit ASNs have recently arrived
- The other capabilities are still in development or not widely implemented or deployed yet

# **BGP for Internet Service Providers**

- BGP Basics
- Scaling BGP
- Using Communities
- Deploying BGP in an ISP network



# **BGP Scaling Techniques**

# **BGP Scaling Techniques**

• How does a service provider:

Scale the iBGP mesh beyond a few peers?

- Implement new policy without causing flaps and route churning?
- Keep the network stable, scalable, as well as simple?

# **BGP Scaling Techniques**

- Route Refresh
- Route Reflectors
- (Confederations)
- Why Route Flap Damping is considered harmful



# **Dynamic Reconfiguration**

**Route Refresh** 

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## **Route Refresh**

- BGP peer reset required after every policy change Because the router does not store prefixes which are rejected by policy
- Hard BGP peer reset:

**Terminates BGP peering & Consumes CPU** 

Severely disrupts connectivity for all networks

#### Soft BGP peer reset (or Route Refresh):

**BGP** peering remains active

Impacts only those prefixes affected by policy change

# **Route Refresh Capability**

- Facilitates non-disruptive policy changes
- For most implementations, no configuration is needed

Automatically negotiated at peer establishment

- No additional memory is used
- Requires peering routers to support "route refresh capability" – RFC2918

# **Dynamic Reconfiguration**

Use Route Refresh capability if supported

find out from the BGP neighbour status display Non-disruptive, "Good For the Internet"

- If not supported, see if implementation has a workaround
- Only hard-reset a BGP peering as a last resort

## **Consider the impact to be** equivalent to a router reboot



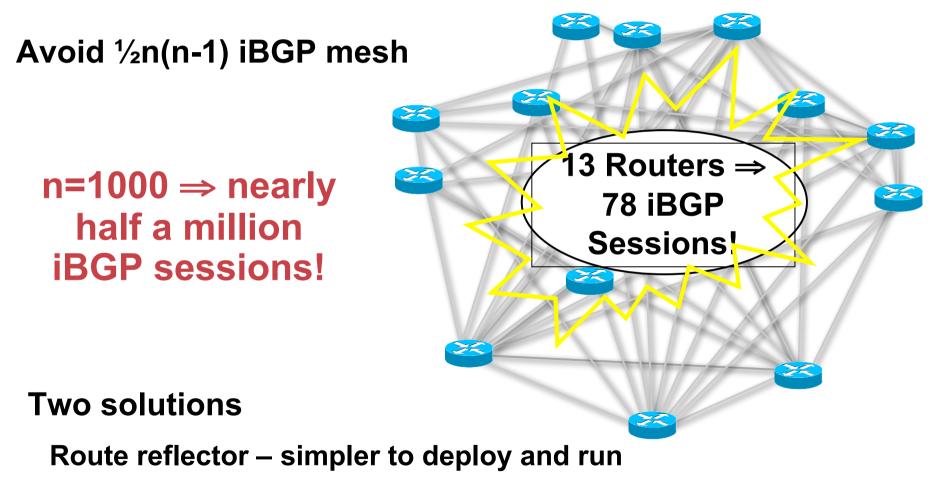
# **Route Reflectors**

Scaling the iBGP mesh

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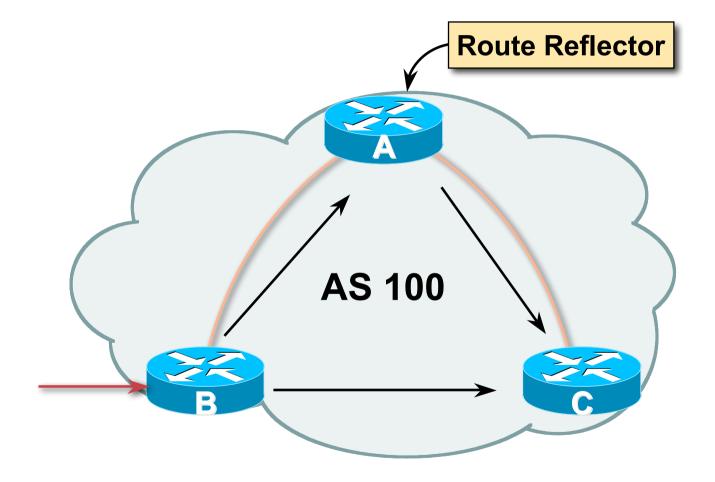
60

# Scaling iBGP mesh



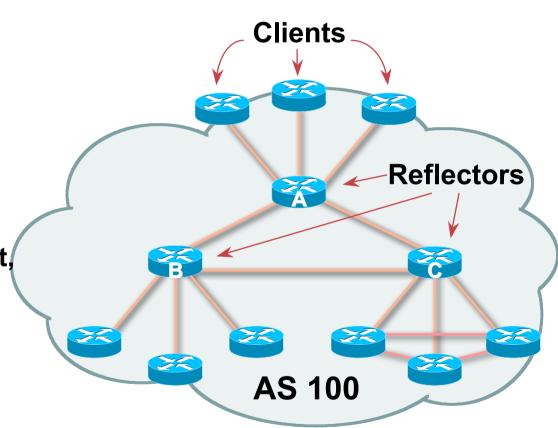
**Confederation – more complex, has corner case advantages** 

# **Route Reflector: Principle**



## **Route Reflector**

- Reflector receives path from clients and non-clients
- Selects best path
- If best path is from client, reflect to other clients and non-clients
- If best path is from non-client, reflect to clients only
- Non-meshed clients
- Described in RFC4456



# **Route Reflector: Topology**

- Divide the backbone into multiple clusters
- At least one route reflector and few clients per cluster
- Route reflectors are fully meshed
- Clients in a cluster could be fully meshed
- Single IGP to carry next hop and local routes

# **Route Reflector: Loop Avoidance**

#### Originator\_ID attribute

Carries the RID of the originator of the route in the local AS (created by the RR)

#### Cluster\_list attribute

The local cluster-id is added when the update is sent by the RR

Best to set cluster-id is from router-id (address of loopback)

(Some ISPs use their own cluster-id assignment strategy – but needs to be well documented!)

## **Route Reflector: Redundancy**

 Multiple RRs can be configured in the same cluster – not advised!

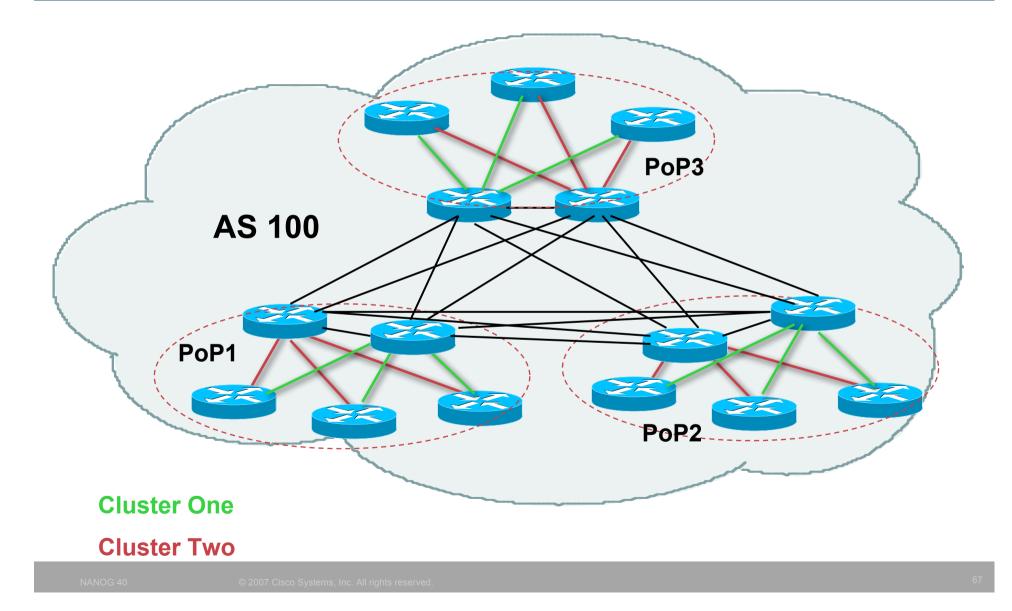
All RRs in the cluster must have the same cluster-id (otherwise it is a different cluster)

#### • A router may be a client of RRs in different clusters

Common today in ISP networks to overlay two clusters – redundancy achieved that way

→ Each client has two RRs = redundancy

# **Route Reflector: Redundancy**



# **Route Reflector: Benefits**

- Solves iBGP mesh problem
- Packet forwarding is not affected
- Normal BGP speakers co-exist
- Multiple reflectors for redundancy
- Easy migration
- Multiple levels of route reflectors

# **Route Reflector: Deployment**

Where to place the route reflectors?

Always follow the physical topology!

This will guarantee that the packet forwarding won't be affected

- Typical ISP network:
  - PoP has two core routers
  - Core routers are RR for the PoP
  - Two overlaid clusters

## **Route Reflector: Migration**

#### • Typical ISP network:

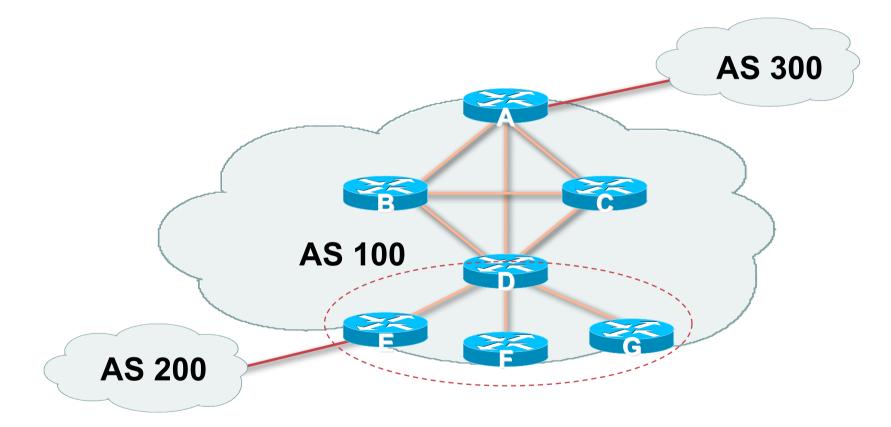
Core routers have fully meshed iBGP Create further hierarchy if core mesh too big Split backbone into regions

### Configure one cluster pair at a time

Eliminate redundant iBGP sessions Place maximum one RR per cluster

Easy migration, multiple levels

# **Route Reflector: Migration**



#### Migrate small parts of the network, one part at a time



# **BGP Confederations**

#### Confederations

Divide the AS into sub-AS

eBGP between sub-AS, but some iBGP information is kept

Preserve NEXT\_HOP across the sub-AS (IGP carries this information)

Preserve LOCAL\_PREF and MED

- Usually a single IGP
- Described in RFC3065

#### **Confederations (Cont.)**

 Visible to outside world as single AS – "Confederation Identifier"

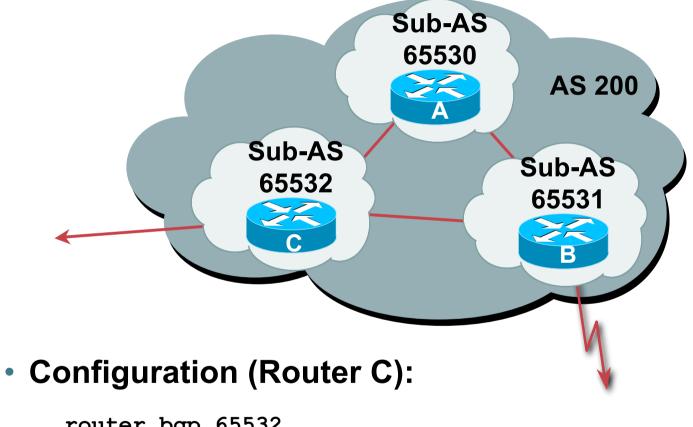
Each sub-AS uses a number from the private AS range (64512-65534)

#### • iBGP speakers in each sub-AS are fully meshed

The total number of neighbours is reduced by limiting the full mesh requirement to only the peers in the sub-AS

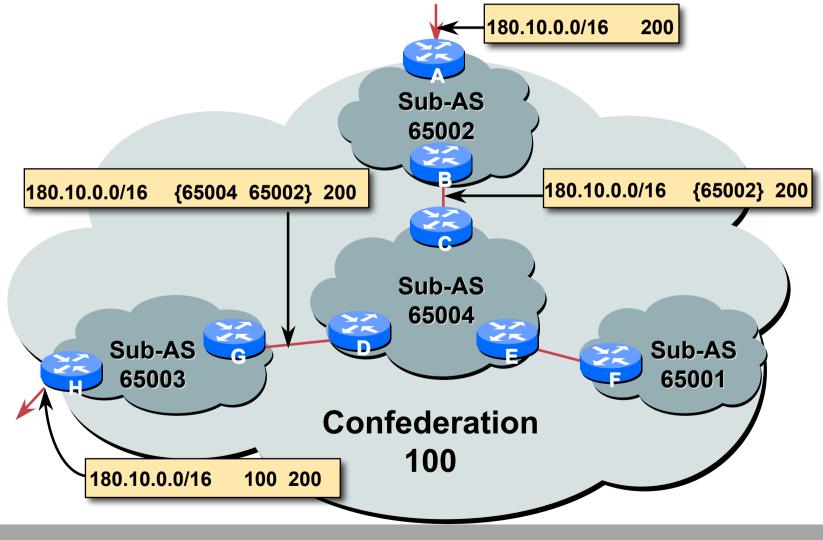
Can also use Route-Reflector within sub-AS

#### Confederations



#### router bgp 65532 bgp confederation identifier 200 bgp confederation peers 65530 65531 neighbor 141.153.12.1 remote-as 65530 neighbor 141.153.17.2 remote-as 65531

#### **Confederations: AS-Sequence**



#### **Route Propagation Decisions**

• Same as with "normal" BGP:

From peer in same sub-AS  $\rightarrow$  only to external peers

From external peers  $\rightarrow$  to all neighbors

"External peers" refers to

Peers outside the confederation

Peers in a different sub-AS

Preserve LOCAL\_PREF, MED and NEXT\_HOP

### **RRs or Confederations**

	Internet Connectivity	Multi-Level Hierarchy	Policy Control	Scalability	Migration Complexity
Confederations	Anywhere in the Network	Yes	Yes	Medium	Medium to High
Route Reflectors	Anywhere in the Network	Yes	Yes	Very High	Very Low

#### Most new service provider networks now deploy Route Reflectors from Day One

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#### More points about Confederations

 Can ease "absorbing" other ISPs into you ISP – e.g., if one ISP buys another

Or can use AS masquerading feature available in some implementations to do a similar thing

 Can use route-reflectors with confederation sub-AS to reduce the sub-AS iBGP mesh



# **Route Flap Damping**

**Network Stability for the 1990s** 

**Network Instability for the 21st Century!** 

#### **Route Flap Damping**

- For many years, Route Flap Damping was a strongly recommended practice
- Now it is strongly discouraged as it appears to cause far greater network instability than it cures
- But first, the theory...

#### **Route Flap Damping**

Route flap

Going up and down of path or change in attribute BGP WITHDRAW followed by UPDATE = 1 flap eBGP neighbour going down/up is NOT a flap Ripples through the entire Internet Wastes CPU

 Damping aims to reduce scope of route flap propagation

## **Route Flap Damping (continued)**

#### Requirements

Fast convergence for normal route changes

History predicts future behaviour

Suppress oscillating routes

Advertise stable routes

#### Implementation described in RFC 2439

#### Operation

Add penalty (1000) for each flap

Change in attribute gets penalty of 500

Exponentially decay penalty

half life determines decay rate

Penalty above suppress-limit

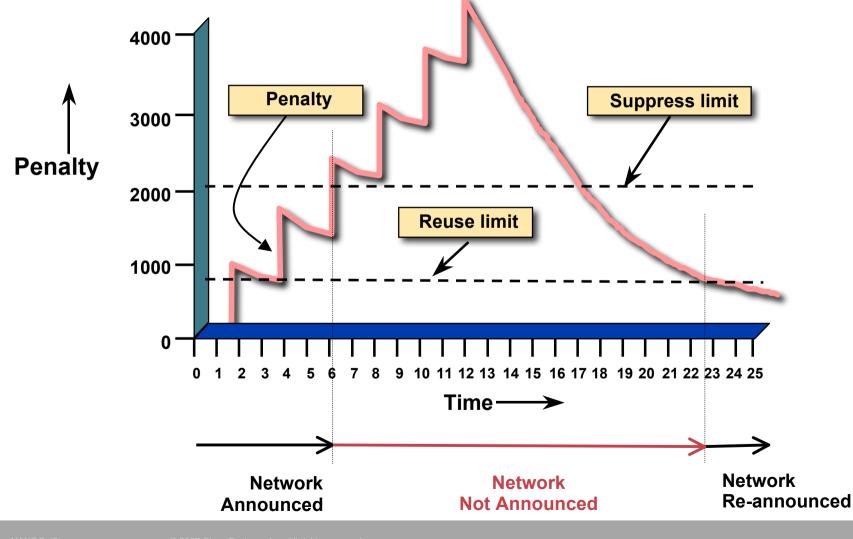
do not advertise route to BGP peers

Penalty decayed below reuse-limit

re-advertise route to BGP peers

penalty reset to zero when it is half of reuse-limit

## **Operation**



#### Operation

- Only applied to inbound announcements from eBGP peers
- Alternate paths still usable
- Controllable by at least:
  - Half-life
  - reuse-limit
  - suppress-limit
  - maximum suppress time

## Configuration

- Implementations allow various policy control with flap damping
  - Fixed damping, same rate applied to all prefixes
  - Variable damping, different rates applied to different ranges of prefixes and prefix lengths

#### **Route Flap Damping History**

- First implementations on the Internet by 1995
- Vendor defaults too severe

RIPE Routing Working Group recommendations in ripe-178, ripe-210, and ripe-229

http://www.ripe.net/ripe/docs

But many ISPs simply switched on the vendors' default values without thinking

#### **Serious Problems:**

- "Route Flap Damping Exacerbates Internet Routing Convergence"
  - Zhuoqing Morley Mao, Ramesh Govindan, George Varghese & Randy H. Katz, August 2002
- "What is the sound of one route flapping?"

Tim Griffin, June 2002

- Various work on routing convergence by Craig Labovitz and Abha Ahuja a few years ago
- "Happy Packets"

Closely related work by Randy Bush et al

#### Problem 1:

• One path flaps:

BGP speakers pick next best path, announce to all peers, flap counter incremented

Those peers see change in best path, flap counter incremented

After a few hops, peers see multiple changes simply caused by a single flap  $\rightarrow$  prefix is suppressed

#### **Problem 2:**

Different BGP implementations have different transit time for prefixes

Some hold onto prefix for some time before advertising

Others advertise immediately

 Race to the finish line causes appearance of flapping, caused by a simple announcement or path change → prefix is suppressed

#### **Solution:**

- Do NOT use Route Flap Damping whatever you do!
- RFD will unnecessarily impair access
  - to your network and
  - to the Internet
- More information contained in RIPE Routing Working Group recommendations:

www.ripe.net/ripe/docs/ripe-378.[pdf,html,txt]

#### **BGP for Internet Service Providers**

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## **Service Provider use of Communities**

Some examples of how ISPs make life easier for themselves

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#### **BGP Communities**

- Another ISP "scaling technique"
- Prefixes are grouped into different "classes" or communities within the ISP network
- Each community means a different thing, has a different result in the ISP network

#### **BGP** Communities

 Communities are generally set at the edge of the ISP network

**Customer edge:** customer prefixes belong to different communities depending on the services they have purchased

Internet edge: transit provider prefixes belong to difference communities, depending on the loadsharing or traffic engineering requirements of the local ISP, or what the demands from its BGP customers might be

Two simple examples follow to explain the concept

- This demonstrates how communities might be used at the customer edge of an ISP network
- ISP has three connections to the Internet:

IXP connection, for local peers

Private peering with a competing ISP in the region

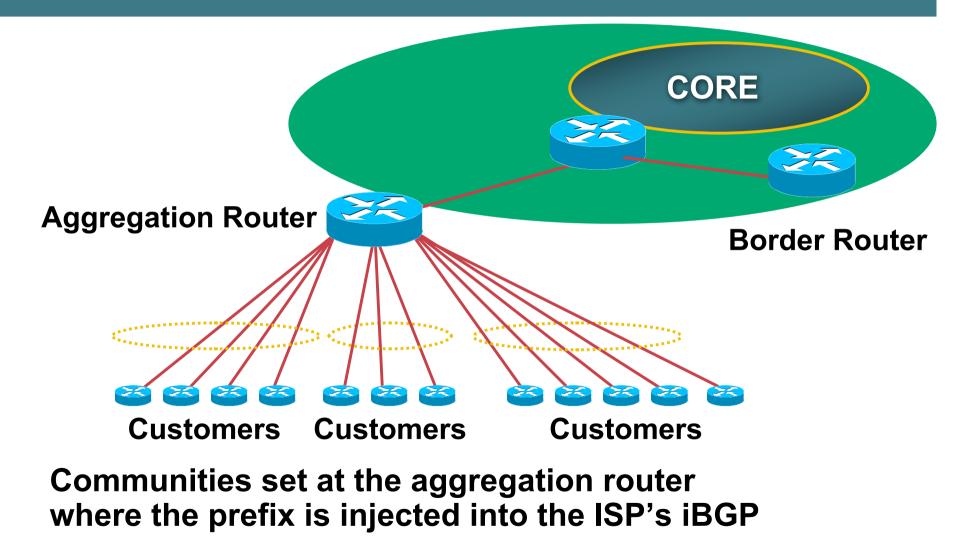
Transit provider, who provides visibility to the entire Internet

 Customers have the option of purchasing combinations of the above connections

Community assignments:

IXP connection:	community 100:2100
Private peer:	community 100:2200

- Customer who buys local connectivity (via IXP) is put in community 100:2100
- Customer who buys peer connectivity is put in community 100:2200
- Customer who wants both IXP and peer connectivity is put in 100:2100 and 100:2200
- Customer who wants "the Internet" has no community set
   We are going to announce his prefix everywhere



- No need to alter filters at the network border when adding a new customer
- New customer simply is added to the appropriate community

Border filters already in place take care of announcements

 $\Rightarrow$  Ease of operation!

#### **Community Example: Internet Edge**

- This demonstrates how communities might be used at the peering edge of an ISP network
- ISP has four types of BGP peers:
  - Customer
  - **IXP** peer
  - **Private peer**
  - Transit provider
- The prefixes received from each can be classified using communities
- Customers can opt to receive any or all of the above

#### **Community Example: Internet Edge**

Community assignments:

Customer prefix:	community 100:3000		
IXP prefix:	community 100:3100		
Private peer prefix:	community 100:3200		

- BGP customer who buys local connectivity gets 100:3000
- BGP customer who buys local and IXP connectivity receives community 100:3000 and 100:3100
- BGP customer who buys full peer connectivity receives community 100:3000, 100:3100, and 100:3200
- Customer who wants "the Internet" gets everything Gets default route originated by aggregation router Or pays money to get all 220k prefixes

### **Community Example: Internet Edge**

 No need to create customised filters when adding customers

**Border router already sets communities** 

Installation engineers pick the appropriate community set when establishing the customer BGP session

 $\Rightarrow$  Ease of operation!

### **Community Example – Summary**

- Two examples of customer edge and internet edge can be combined to form a simple community solution for ISP prefix policy control
- More experienced operators tend to have more sophisticated options available

Advice is to start with the easy examples given, and then proceed onwards as experience is gained

#### Some ISP Examples

- ISPs also create communities to give customers bigger routing policy control
- Public policy is usually listed in the IRR

Following examples are all in the IRR

Examples build on the configuration concepts from the introductory example

 Consider creating communities to give policy control to customers

**Reduces technical support burden** 

Reduces the amount of router reconfiguration, and the chance of mistakes

#### Some ISP Examples: Sprintlink

000		Welcome to Sprint.net		
	.//www.sprint.net/ii	ndex.php?module=policies/bgp_policy	• Q▼ Google	
			Google	
Welcome to Sprint.net	п евау тапоо! п	News (144) ▼ Apple (40) ▼		
				6
	WHAT YOU CAN	CONTROL		
	AS-PATH PREPENDS			
		ers to use AS-path prepending to adjust route prefere notifiying Sprint of your change in announcments.	ence on the network. Such prepending will be received and passed	
		will prepend AS1239 to eBGP sessions with certain aut pported: 1668, 209, 2914, 3300, 3356, 3549, 3561, 4635	conomous systems depending on a received community. Currently, the i, 701, 7018, 702 and 8220.	
	String	Resulting AS Path to ASXXX		
	65000:XXX	Do not advertise to ASXXX		
	65001:XXX	1239 (default)		
	65002:XXX	1239 1239		
	65003:XXX	1239 1239 1239		
	65004:XXX	1239 1239 1239 1239		
	String	Resulting AS Path to ASXXX in Asia		
	65070:XXX	Do not advertise to ASXXX		
	65071:XXX	1239 (default)	More info at	
	65072:XXX	1239 1239	www.sprintlink.net/pol	icy/ban html
	65073:XXX	1239 1239 1239	www.spinitink.neupoi	icy/bgp.num
	65074:XXX	1239 1239 1239 1239		
	String	Resulting AS Path to ASXXX in Europe		
	65050:XXX	Do not advertise to ASXXX		
	65051:XXX	1239 (default)		
	65052:XXX	1239 1239		
	65053:XXX	1239 1239 1239		

# Some ISP Examples AAPT

- Australian ISP
- Run their own Routing Registry
   Whois.connect.com.au
- Offer 6 different communities to customers to aid with their traffic engineering

# Some ISP Examples AAPT

aut-num:	AS2764
as-name:	ASN-CONNECT-NET
descr:	AAPT Limited
admin-c:	CNO2-AP
tech-c:	CNO2-AP
remarks:	Community support definitions
remarks:	
remarks:	Community Definition
remarks:	
remarks:	2764:2 Don't announce outside local POP
remarks:	2764:4 Lower local preference by 15
remarks:	2764:5 Lower local preference by 5
remarks:	2764:6 Announce to customers and all peers
	(incl int'l peers), but not transit
remarks:	2764:7 Announce to customers only
remarks:	2764:14 Announce to AANX
notify:	routing@connect.com.au
mnt-by:	CONNECT-AU
changed:	nobody@connect.com.au 20050225
source:	CCAIR

#### More at http://info.connect.com.au/docs/routing/general/multi-faq.shtml#q13

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#### Some ISP Examples Verizon Business EMEA

- Verizon Business' European operation
- Permits customers to send communities which determine
  - **local preferences within Verizon Business' network**
  - **Reachability of the prefix**
  - How the prefix is announced outside of Verizon Business' network

#### Some ISP Examples Verizon Business Europe

aut-num:	AS702							
descr:	Verizon B	usiness EMEA - Commercial IP service provider in Eur						
remarks:	VzBi uses	Bi uses the following communities with its customers:						
	702:80	Set Local Pref 80 within AS702						
	702:120	Set Local Pref 120 within AS702						
	702:20	Announce only to VzBi AS'es and VzBi customers						
	702:30	Keep within Europe, don't announce to other VzBi AS						
	702:1	Prepend AS702 once at edges of VzBi to Peers						
	702:2	Prepend AS702 twice at edges of VzBi to Peers						
	702:3	:3 Prepend AS702 thrice at edges of VzBi to Peers						
	Advanced of	Advanced communities for customers						
	702:7020	Do not announce to AS702 peers with a scope of						
		National but advertise to Global Peers, European						
		Peers and VzBi customers.						
	702:7001	Prepend AS702 once at edges of VzBi to AS702						
		peers with a scope of National.						
	702:7002	Prepend AS702 twice at edges of VzBi to AS702						
		peers with a scope of National.						
(more)								

#### Some ISP Examples VzBi Europe

(more)							
	702:7003 Prepend AS702 thrice at edges of VzBi to AS702						
	peers with a scope of National.						
	702:8020 Do not announce to AS702 peers with a scope of						
	European but advertise to Global Peers, National						
	Peers and VzBi customers.						
	702:8001 Prepend AS702 once at edges of VzBi to AS702						
	peers with a scope of European.						
	702:8002 Prepend AS702 twice at edges of VzBi to AS702						
	peers with a scope of European.						
	702:8003 Prepend AS702 thrice at edges of VzBi to AS702						
	peers with a scope of European.						
	Additional details of the VzBi communities are located at:						
	http://www.verizonbusiness.com/uk/customer/bgp/						
mnt-by:	WCOM-EMEA-RICE-MNT						
source:	RIPE						

#### Some ISP Examples BT Ignite

 One of the most comprehensive community lists around

Seems to be based on definitions originally used in Tiscali's network

whois -h whois.ripe.net AS5400 reveals all

 Extensive community definitions allow sophisticated traffic engineering by customers

#### Some ISP Examples BT Ignite

aut-num:	AS5400	
descr:	BT Ignite European Backbone	
remarks:		
remarks:	Community to	Community to
remarks:	Not announce To peer:	AS prepend 5400
remarks:		
remarks:	5400:1000 All peers & Transits	5400:2000
remarks:		
remarks:	5400:1500 All Transits	5400:2500
remarks:	5400:1501 Sprint Transit (AS1239)	5400:2501
remarks:	5400:1502 SAVVIS Transit (AS3561)	5400:2502
remarks:	5400:1503 Level 3 Transit (AS3356)	5400:2503
remarks:	5400:1504 AT&T Transit (AS7018)	5400:2504
remarks:	5400:1506 GlobalCrossing Trans(AS3549)	5400:2506
remarks:		
remarks:	5400:1001 Nexica (AS24592)	5400:2001
remarks:	5400:1002 Fujitsu (AS3324)	5400:2002
remarks:	5400:1004 C&W EU (1273)	5400:2004
<snip></snip>		
notify:	notify@eu.bt.net And mar	าง
mnt-by:		
source:	RIPE many mo	

#### Some ISP Examples Carrier1

- European ISP
- Another very comprehensive list of community definitions

whois -h whois.ripe.net AS8918 reveals all

#### Some ISP Examples Carrier1

aut-num:						
descr: <snip></snip>	Carrier1 Autonomous System					
remarks:	Community Definition					
remarks:	*					
remarks:	8918:2000 Do not announce to C1 customers					
remarks:	8918:2010 Do not announce to C1 peers, peers+ and transit					
remarks:	8918:2015 Do not announce to C1 transit providers					
remarks:	*					
remarks:	8918:2020 Do not announce to Teleglobe (AS 6453)					
remarks:	8918:2035 Do not announce to UUNet (AS 702)					
remarks:	8918:2050 Do not announce to T-Systems (AS 3320)					
remarks:	8918:2060 Do not announce to JointTransit (AS 24785/20562)					
remarks:	*					
remarks:	8918:2070 Do not announce to AMS-IX peers					
remarks:	8918:2080 Do not announce to NL-IX peers					
<snip></snip>						
notify:	inoc@carrier1.net					
mnt-by:	CARRIER1-MNT					
source:	RIPE And many					
	many more!					

#### Some ISP Examples Level 3

- Highly detailed AS object held on the RIPE Routing Registry
- Also a very comprehensive list of community definitions

whois -h whois.ripe.net AS3356 reveals all

#### Some ISP Examples Level 3

aut-num:	AS3356				
descr:	Level 3 Communications				
<snip></snip>					
remarks:					
remarks:	customer traffic engineering communities - Suppression				
remarks:					
remarks:	64960:XXX - announce to AS XXX if 65000:0				
remarks:	65000:0 - announce to customers but not to peers				
remarks:	65000:XXX - do not announce at peerings to AS XXX				
remarks:					
remarks:	customer traffic engineering communities - Prepending				
remarks:					
remarks:	65001:0 - prepend once to all peers				
remarks:	65001:XXX - prepend once at peerings to AS XXX				
<snip></snip>					
remarks:	3356:70 - set local preference to 70				
remarks:	3356:80 - set local preference to 80				
remarks:	3356:90 - set local preference to 90				
remarks:	- 3356:9999 - blackhole (discard) traffic				
<snip></snip>					
mnt-by:	LEVEL3-MNT				
source:	RIPE And many				
	many more!				

### **BGP for Internet Service Providers**

- BGP Basics
- Scaling BGP
- Using Communities
- Deploying BGP in an ISP network



# **Deploying BGP in an ISP Network**

Okay, so we've learned all about BGP now; how do we use it on our network??

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### **Deploying BGP**

- The role of IGPs and iBGP
- Aggregation
- Receiving Prefixes
- Configuration Tips



# The role of IGP and iBGP

Ships in the night? Or Good foundations?

### **BGP versus OSPF/ISIS**

- Internal Routing Protocols (IGPs)
  - examples are ISIS and OSPF
  - used for carrying infrastructure addresses
  - **NOT** used for carrying Internet prefixes or customer prefixes
  - design goal is to minimise number of prefixes in IGP to aid scalability and rapid convergence

### **BGP versus OSPF/ISIS**

BGP used internally (iBGP) and externally (eBGP)

#### iBGP used to carry

some/all Internet prefixes across backbone

customer prefixes

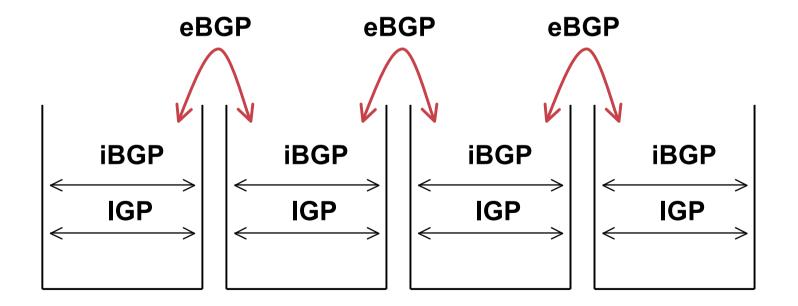
#### eBGP used to

exchange prefixes with other ASes

implement routing policy

### **BGP/IGP model used in ISP networks**

Model representation



### **BGP versus OSPF/ISIS**

#### • DO NOT:

distribute BGP prefixes into an IGP

distribute IGP routes into BGP

use an IGP to carry customer prefixes

• YOUR NETWORK WILL NOT SCALE

## Injecting prefixes into iBGP

- Use iBGP to carry customer prefixes
   Don't ever use IGP
- Point static route to customer interface
- Enter network into BGP process

Ensure that implementation options are used so that the prefix always remains in iBGP, regardless of state of interface

i.e. avoid iBGP flaps caused by interface flaps



# Aggregation

**Quality or Quantity?** 

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### Aggregation

- Aggregation means announcing the address block received from the RIR to the other ASes connected to your network
- Subprefixes of this aggregate *may* be:

Used internally in the ISP network

Announced to other ASes to aid with multihoming

 Unfortunately too many people are still thinking about class Cs, resulting in a proliferation of /24s in the Internet routing table

## Aggregation

- Address block should be announced to the Internet as an aggregate
- Subprefixes of address block should NOT be announced to Internet unless special circumstances (more later)
- Aggregate should be generated internally

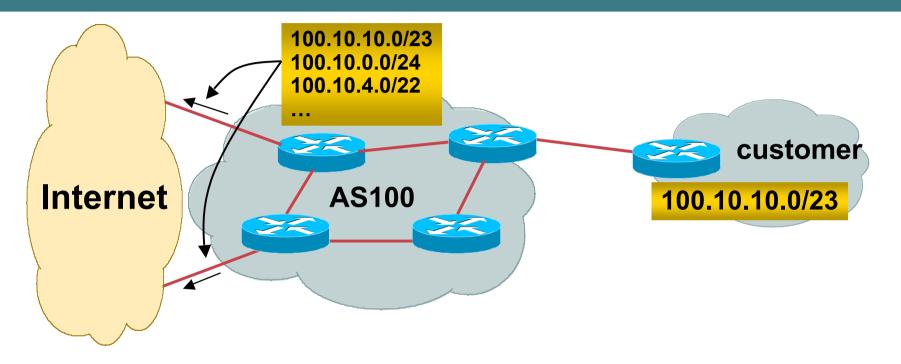
Not on the network borders!

## Announcing an Aggregate

- ISPs who don't and won't aggregate are held in poor regard by community
- Registries publish their minimum allocation size Anything from a /20 to a /22 depending on RIR
- No real reason to see anything longer than a /22 prefix in the Internet

BUT there are currently >117000 /24s!

### **Aggregation – Example**



- Customer has /23 network assigned from AS100's /19 address block
- AS100 announces customers' individual networks to the Internet

#### **Aggregation – Bad Example**

- Customer link goes down
  - Their /23 network becomes unreachable
  - /23 is withdrawn from AS100's iBGP
- Their ISP doesn't aggregate its /19 network block
  - /23 network withdrawal announced to peers
  - starts rippling through the Internet
  - added load on all Internet backbone routers as network is removed from routing table

#### Customer link returns

Their /23 network is now visible to their ISP

Their /23 network is readvertised to peers

Starts rippling through Internet

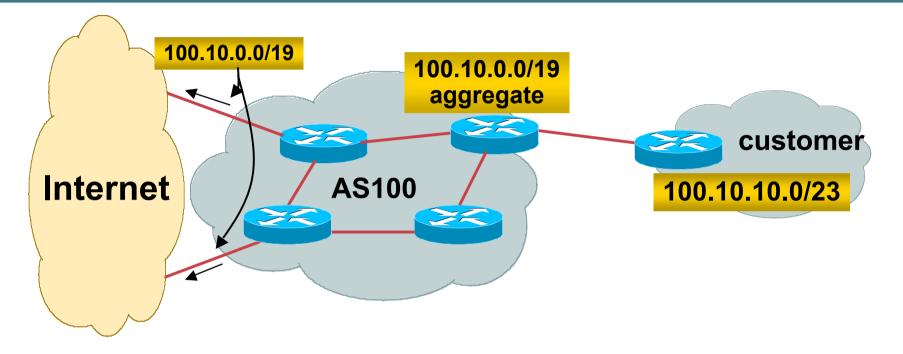
Load on Internet backbone routers as network is reinserted into routing table

Some ISP's suppress the flaps

Internet may take 10-20 min or longer to be visible

Where is the Quality of Service???

### **Aggregation – Example**



- Customer has /23 network assigned from AS100's /19 address block
- AS100 announced /19 aggregate to the Internet

### **Aggregation – Good Example**

- Customer link goes down
  - their /23 network becomes unreachable
  - /23 is withdrawn from AS100's iBGP
- /19 aggregate is still being announced
  - no BGP hold down problems
  - no BGP propagation delays
  - no damping by other ISPs

- Customer link returns
- Their /23 network is visible again

The /23 is re-injected into AS100's iBGP

- The whole Internet becomes visible immediately
- Customer has Quality of Service perception

#### **Aggregation – Summary**

- Good example is what everyone should do!
  - Adds to Internet stability
  - **Reduces size of routing table**
  - **Reduces routing churn**
  - Improves Internet QoS for everyone
- Bad example is what too many still do!
  - Why? Lack of knowledge?
  - Laziness?

## The Internet Today (May 2007)

Current Internet Routing Table Statistics

BGP Routing Table Entries	220947
Prefixes after maximum aggregation	117493
Unique prefixes in Internet	107152
Prefixes smaller than registry alloc	115865
/24s announced	117175
only 5762 /24s are from 192.0.0.0/8	
ASes in use	25241

### "The New Swamp"

Swamp space is name used for areas of poor aggregation

The original swamp was 192.0.0.0/8 from the former class C block

Name given just after the deployment of CIDR

The new swamp is creeping across all parts of the Internet

Not just RIR space, but "legacy" space too

#### "The New Swamp" RIR Space – February 1999

# **RIR blocks contribute 49393 prefixes or 88% of the Internet Routing Table**

Block	Networks	Block	Networks	Block	Networks	Block	Networks
24/8	165	77/8	0	118/8	0	203/8	3622
41/8	0	78/8	0	119/8	0	204/8	3792
58/8	0	79/8	0	120/8	0	205/8	2584
59/8	0	80/8	0	121/8	0	206/8	3127
60/8	0	81/8	0	122/8	0	207/8	2723
61/8	0 3	82/8	0	123/8	0	208/8	2817
62/8	87	83/8	0	124/8	0	209/8	2574
63/8	20	84/8	0	125/8	0	210/8	617
64/8	0	85/8	0	126/8	0	211/8	0
65/8	0	86/8	0	189/8	0	212/8	717
66/8	0	87/8	0	190/8	0	213/8	1
67/8	0	88/8	0	192/8	6275	216/8	943
68/8	0	89/8	0	193/8	2390	217/8	0
69/8	0	90/8	0	194/8	2932	218/8	0
70/8	0	91/8	0	195/8	1338	219/8	0
71/8	0	96/8	0	196/8	513	220/8	0
72/8	0	97/8	0	198/8	4034	221/8	0
73/8	0	98/8	0	199/8	3495	222/8	0
74/8	0	99/8	0	200/8	1348		
75/8	0	116/8	0	201/8	0		
76/8	0	117/8	0	202/8	2276		

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#### "The New Swamp" RIR Space – February 2007

# **RIR blocks contribute 192490 prefixes or 90% of the Internet Routing Table**

Block	Networks	Block	Networks	Block	Networks	Block	Networks
24/8	2930	77/8	1214	118/8	3	203/8	10459
41/8	288	78/8	8	119/8	3	204/8	5569
58/8	1097	79/8	2	120/8	3	205/8	2892
59/8	1152	80/8	2053	121/8	426	206/8	3857
60/8	604	81/8	1695	122/8	698	207/8	4331
61/8	2589	82/8	1564	123/8	534	208/8	4258
62/8	2193	83/8	1172	124/8	1340	209/8	5540
63/8	2967	84/8	1269	125/8	1554	210/8	4759
64/8	5501	85/8	1891	126/8	41	211/8	2733
65/8	3917	86/8	800	189/8	169	212/8	2900
66/8	6575	87/8	1157	190/8	1077	213/8	3052
67/8	2015	88/8	847	192/8	6927	216/8	6930
68/8	2770	89/8	1970	193/8	5704	217/8	2615
69/8	3702	90/8	105	194/8	4652	218/8	1561
70/8	1693	91/8	577	195/8	4279	219/8	1197
71/8	1188	96/8	8	196/8	1600	220/8	1988
72/8	2878	97/8	1	198/8	4748	221/8	894
73/8	273	98/8	3	199/8	4184	222/8	1241
74/8	1483	99/8	0	200/8	7482		
75/8	483	116/8	3 3	201/8	2927		
76/8	194	117/8	3	202/8	10529		

### "The New Swamp" Summary

#### RIR space shows creeping deaggregation

It seems that an RIR /8 block averages around 5000 prefixes once fully allocated

So their existing 81 /8s will eventually cause 405000 prefix announcements

#### • Food for thought:

Remaining 48 unallocated /8s and the 81 RIR /8s combined will cause:

645000 prefixes with 5000 prefixes per /8 density

774000 prefixes with 6000 prefixes per /8 density

Plus 12% due to "non RIR space deaggregation"

→ Routing Table size of 866880 prefixes

### "The New Swamp" Summary

- Rest of address space is showing similar deaggregation too <sup>(2)</sup>
- What are the reasons?
  - Main justification is traffic engineering
- Real reasons are:
  - Lack of knowledge
  - Laziness
  - **Deliberate & knowing actions**

### BGP Report (bgp.potaroo.net)

199336 total announcements in October 2006

#### 129795 prefixes

After aggregating including full AS PATH info

i.e. including each ASN's traffic engineering

35% saving possible

#### 109034 prefixes

After aggregating by Origin AS

i.e. ignoring each ASN's traffic engineering

10% saving possible

### **Deaggregation: The Excuses**

- Traffic engineering causes 10% of the Internet Routing table
- Deliberate deaggregation causes 35% of the Internet Routing table

### **Efforts to Improve Aggregation**

#### • The CIDR Report

Initiated and operated for many years by Tony Bates Now combined with Geoff Huston's routing analysis

www.cidr-report.org

Results e-mailed on a weekly basis to most operations lists around the world

Lists the top 30 service providers who could do better at aggregating

### Efforts to Improve Aggregation The CIDR Report

- Also computes the size of the routing table assuming ISPs performed optimal aggregation
- Website allows searches and computations of aggregation to be made on a per AS basis

Flexible and powerful tool to aid ISPs

Intended to show how greater efficiency in terms of BGP table size can be obtained without loss of routing and policy information

Shows what forms of origin AS aggregation could be performed and the potential benefit of such actions to the total table size

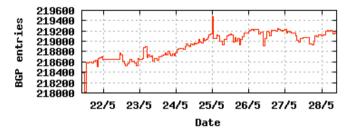
Very effectively challenges the traffic engineering excuse



#### **Status Summary**

#### **Table History**

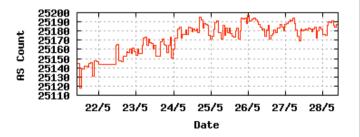
Date	Prefixes	CIDR Aggregated
21-05-07	218385	140025
22-05-07	218650	139831
23-05-07	218653	139850
24-05-07	218776	139698
25-05-07	219469	139898
26-05-07	219203	139943
27-05-07	219232	139870
28-05-07	219115	140020



Plot: BGP Table Size

#### AS Summary

- 25190 Number of ASes in routing system
- 10666 Number of ASes announcing only one prefix
- 1483 Largest number of prefixes announced by an AS AS7018: ATT-INTERNET4 - AT&T WorldNet Services
- <sup>89890048</sup> Largest address span announced by an AS (/32s) AS721: DISA-ASNBLK - DoD Network Information Center



Plot: AS count

Plot: Average announcements per origin AS

Report: ASes ordered by originating address span

Report: ASes ordered by transit address span

Report: Autonomous System number-to-name mapping (from Registry WHOIS data)

000	CIDR R	eport
<ul> <li>+ Shttp:</li> </ul>	//www.cidr-report.org/	<ul> <li>Q→ Google</li> </ul>
Apple .Mac Amazon	eBay Yahoo! News (135) ▼ Apple (40) ▼	
CIDR Report		

#### **Aggregation Summary**

The algorithm used in this report proposes aggregation only when there is a precise match using AS path so as to preserve traffic transit policies. Aggregation is also proposed across non-advertised address space ('holes').

#### --- 28May07 ---

20110	y07			-	
ASnum	NetsNow	NetsAggr	NetGain	9/ Gain	<sup>6</sup> Description
Table	219180	140003	79177	36.1%	All ASes
AS18566	1010	31	979	96.9%	COVAD - Covad Communications Co.
AS4755	1221	271	950	77.8%	VSNL-AS Videsh Sanchar Nigam Ltd. Autonomous System
AS4134	1257	316	941	74.9%	CHINANET-BACKBONE No.31, Jin-rong Street
AS4323	1266	369	897	70.9%	TWTC - Time Warner Telecom, Inc.
AS9498	982	98	884	90.0%	BBIL-AP BHARTI BT INTERNET LTD.
AS8151	1306	442	864	66.2%	Uninet S.A. de C.V.
AS6478	1100	237	863	78.5%	ATT-INTERNET3 - AT&T WorldNet Services
AS11492	1074	365	709	66.0%	CABLEONE - CABLE ONE
AS22773	720	56	664	92.2%	CCINET-2 - Cox Communications Inc.
AS19262	777	210	567	73.0%	VZGNI-TRANSIT - Verizon Internet Services Inc.
AS6197	1040	516	524	50.4%	BATI-ATL - BellSouth Network Solutions, Inc
AS17488	694	182	512	73.8%	HATHWAY-NET-AP Hathway IP Over Cable Internet
AS18101	546	36	510	93.4%	RIL-IDC Reliance Infocom Ltd Internet Data Centre,
AS7018	1483	980	503	33.9%	ATT-INTERNET4 - AT&T WorldNet Services
AS19916	568	75	493	86.8%	ASTRUM-0001 - OLM LLC
AS7545	629	164	465	73.9%	TPG-INTERNET-AP TPG Internet Pty Ltd
AS15270	535	83	452	84.5%	AS-PAETEC-NET - PaeTec.net -a division of PaeTecCommunications, Inc.
AS17676	504	65	439	87.1%	JPNIC-JP-ASN-BLOCK Japan Network Information Center
AS4766	740	317	423		KIXS-AS-KR Korea Telecom
AS2386	1151	759	392	34.1%	INS-AS - AT&T Data Communications Services
AS9443	464	79	385	83.0%	INTERNETPRIMUS-AS-AP Primus Telecommunications
AS4812	455	76	379	83.3%	CHINANET-SH-AP China Telecom (Group)
AS7029	587	239	348	59.3%	WINDSTREAM - Windstream Communications Inc

\*

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Apple .Mac	Amazon eBa	y Yahoo! News (135) ▼ Apple (40) ▼		
CIDR Report				
Top 20 A	dded Rou	tes this week per Originating AS		
Prefixes	ASnum	AS Description		
65	AS8151	Uninet S.A. de C.V.		
62	AS577	BACOM - Bell Canada		
45	AS5089	NTL NTL Group Limited		
		RAPIDUS - COGECO Cable Canada Inc.		
		INDONET-AS-AP INDO Internet, PT		
		CHINATELECOM-HLJ-AS-AP asn for Heilongjiang Provincial Net of CT		
		SWISSCOM Swisscom Solutions Ltd		
		CCH-AS7 - Comcast Cable Communications Holdings, Inc		
		TPG-INTERNET-AP TPG Internet Pty Ltd		
		DDN-ASNBLK1 - DoD Network Information Center		
		FasoNet-AS		
		HATHWAY-NET-AP Hathway IP Over Cable Internet		
		ASN-QWEST - Qwest		
		CABLEONE - CABLE ONE		
		NET Servicos de Comunicao S.A.		
		HUNT-BROTHERS-OF-LOUISIANA-LLC - Hunt Brothers		
		CAJUNNET-ASN - CAJUNNET LLC		
		Suporte Tecnologia e Instalações Ltda.		
		SKKUNET-AS SungKyunKwan University (SKKU)		
14	AS38548	INFRATEL-AS-ID-AP PT. Info Sarana Telekomunikasi		
Top 20 V	Vithdrawn	Routes this week per Originating AS		

- -143 AS20858 EGYNET-AS
- -39 AS3602 AS3602-RTI Rogers Telecom Inc.
- -28 AS33490 DNEO-OSP5 Comcast Cable Communications, Inc.
- -24 AS721 DISA-ASNBLK DoD Network Information Center
- -24 AS22909 DNEO-OSP1 Comcast Cable Communications, Inc.
- -23 AS21455 PLANETSKY Planetsky Ltd. AS object
- -23 AS23966 DANCOM-AS-AP Dancom Online Services
- -21 AS9121 TTNET TTnet Autonomous System
- -20 AS2706 HKSUPER-HK-AP Pacific Internet (Hong Kong) Limited
- -19 AS9584 GENESIS-AP Divixian.com Limited

CIDR Report	
C + Shttp://www.cidr-report.org/	<ul> <li>Q- Google</li> </ul>
□ Apple .Mac Amazon eBay Yahoo! News (135) ▼ Apple (40) ▼	
CIDR Report	

#### **More Specifics**

A list of route advertisements that appear to be more specific than the original Class-based prefix mask, or more specific than the registry allocation size.

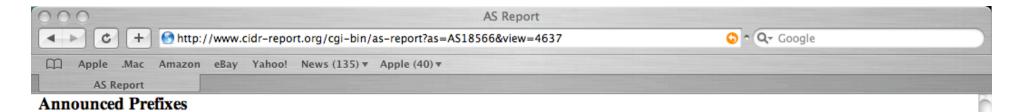
#### Top 20 ASes advertising more specific prefixes

More Specifics	Total Prefixes	ASnum	AS Description
1299	1306	AS8151	Uninet S.A. de C.V.
1218	1483	AS7018	ATT-INTERNET4 - AT&T WorldNet Services
1205	1221	AS4755	VSNL-AS Videsh Sanchar Nigam Ltd. Autonomous System
1110	1110	AS9583	SIFY-AS-IN Sify Limited
1100	1100	AS6478	ATT-INTERNET3 - AT&T WorldNet Services
1074	1266	AS4323	TWTC - Time Warner Telecom, Inc.
1069	1074	AS11492	CABLEONE - CABLE ONE
1058	1151	AS2386	INS-AS - AT&T Data Communications Services
1019	1040	AS6197	BATI-ATL - BellSouth Network Solutions, Inc
1001	1010	AS18566	COVAD - Covad Communications Co.
983	1257	AS4134	CHINANET-BACKBONE No.31, Jin-rong Street
963	982	AS9498	BBIL-AP BHARTI BT INTERNET LTD.
791	807	AS20115	CHARTER-NET-HKY-NC - Charter Communications
790	799	AS7011	FRONTIER-AND-CITIZENS - Frontier Communications of America, Inc.
738			VZGNI-TRANSIT - Verizon Internet Services Inc.
729	946	AS174	COGENT Cogent/PSI
702	740	AS4766	KIXS-AS-KR Korea Telecom
694	694	AS17488	HATHWAY-NET-AP Hathway IP Over Cable Internet
692			CCINET-2 - Cox Communications Inc.
635	954	AS701	UUNET - MCI Communications Services, Inc. d/b/a Verizon Business
Bonorty A	Coo ordoro	d by purch	vor of more creditic profives

Report: ASes ordered by number of more specific prefixes Report: More Specific prefix list (by AS) Report: More Specific prefix list (ordered by prefix)

OOO AS Report	
Image: Antip://www.cidr-report.org/cgi-bin/as-report?as=AS4134&view=4637	
☐ Apple .Mac Amazon eBay Yahoo! News (135) ▼ Apple (40) ▼	
AS Report	
Announced Prefixes	
Rank AS Type Originate Addr Space (pfx) Transit Addr space (pfx) Description 4 AS4134 ORG+TRN Originate: 56476672/6.25 Transit: 30243328/7.15 CHINANET-BACKBONE No.31,Jin-rong Stre	et O
Aggregation Suggestions	
This report does not take into account conditions local to each origin AS in terms of policy or traffic engineering requirements, so this is an approxim guideline as to aggregation possibilities.	iate

AS 4134: CHINANET	-BACKBONE No.31, Jin-rong Street
Prefix (AS Pat)	h) Aggregation Action
58.30.0.0/15	4608 1221 4637 4134
58.32.0.0/13	4608 1221 4637 4134
58.40.0.0/15	4608 1221 4637 4134
58.42.0.0/15	4608 1221 4637 4134 + Announce - aggregate of 58.42.0.0/16 (4608 1221 4637 4134) and 58.43.0.0/1
58.42.0.0/17	4608 1221 4637 4134 - Withdrawn - aggregated with 58.42.128.0/17 (4608 1221 4637 4134)
58.42.128.0/17	4608 1221 4637 4134 - Withdrawn - aggregated with 58.42.0.0/17 (4608 1221 4637 4134)
58.43.0.0/16	4608 1221 4637 4134 - Withdrawn - aggregated with 58.42.0.0/16 (4608 1221 4637 4134)
58.44.0.0/14	4608 1221 4637 4134
58.48.0.0/13	4608 1221 4637 4134
58.48.0.0/14	4608 1221 4637 4134 - Withdrawn - matching aggregate 58.48.0.0/13 4608 1221 4637 4134
58.52.0.0/14	4608 1221 4637 4134  — Withdrawn — matching aggregate 58.48.0.0/13 4608 1221 4637 4134
58.56.0.0/15	4608 1221 4637 4134
58.58.0.0/15	4608 1221 4637 4134 + Announce - aggregate of 58.58.0.0/16 (4608 1221 4637 4134) and 58.59.0.0/1
58.58.0.0/16	4608 1221 4637 4134 - Withdrawn - aggregated with 58.59.0.0/16 (4608 1221 4637 4134)
58.59.0.0/17	4608 1221 4637 4134 - Withdrawn - aggregated with 58.59.128.0/17 (4608 1221 4637 4134)
58.59.128.0/17	4608 1221 4637 4134 - Withdrawn - aggregated with 58.59.0.0/17 (4608 1221 4637 4134)
58.59.128.0/19	4608 1221 4637 4134  — Withdrawn — matching aggregate 58.59.128.0/17 4608 1221 4637 4134
58.59.160.0/19	4608 1221 4637 4134 - Withdrawn - matching aggregate 58.59.128.0/17 4608 1221 4637 4134
58.59.192.0/19	4608 1221 4637 4134  — Withdrawn — matching aggregate 58.59.128.0/17 4608 1221 4637 4134
58.59.224.0/19	4608 1221 4637 4134 - Withdrawn - matching aggregate 58.59.128.0/17 4608 1221 4637 4134
58.60.0.0/14	4608 1221 4637 4134
58.60.0.0/15	4608 1221 4637 4134 - Withdrawn - matching aggregate 58.60.0.0/14 4608 1221 4637 4134
58.62.0.0/15	4608 1221 4637 4134 - Withdrawn - matching aggregate 58.60.0.0/14 4608 1221 4637 4134
58.66.0.0/17	4608 1221 4637 4134
58.66.128.0/18	4608 1221 4637 4134
58.67.0.0/17	4608 1221 4637 4134
58.82.0.0/17	4608 1221 4637 4134



Rank AS Type Originate Addr Space (pfx) Transit Addr space (pfx) Description 144 AS18566 ORIGIN Originate: 2268160 /10.89 Transit: 0 /0.00 COVAD - Covad Communications Co.

#### Aggregation Suggestions

64.105.58.0/23

64.105.60.0/23

64.105.62.0/23

64.105.64.0/23 64.105.66.0/23

64.105.68.0/23

This report does not take into account conditions local to each origin AS in terms of policy or traffic engineering requirements, so this is an approximate guideline as to aggregation possibilities.

Rank AS AS Nam 2 <u>AS18566</u> COVAD	Current Wthdw Aggte Annce Redctn % Covad Communications Co. 1010 979 0 31 979 96.93%	
AS18566: COVAD - Cov	d Communications Co.	
Prefix (AS Path)	Aggregation Action	
64.105.0.0/16	4608 1221 4637 3356 18566	
64.105.0.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.4.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.6.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.8.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.10.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.14.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.16.0/24	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.17.0/24	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.18.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.20.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.22.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.24.0/21	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.32.0/21	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.40.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.42.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.44.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.46.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.48.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.50.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.52.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.54.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6
64.105.56.0/23	4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 335	6

4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 3356 18566

4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 3356 18566

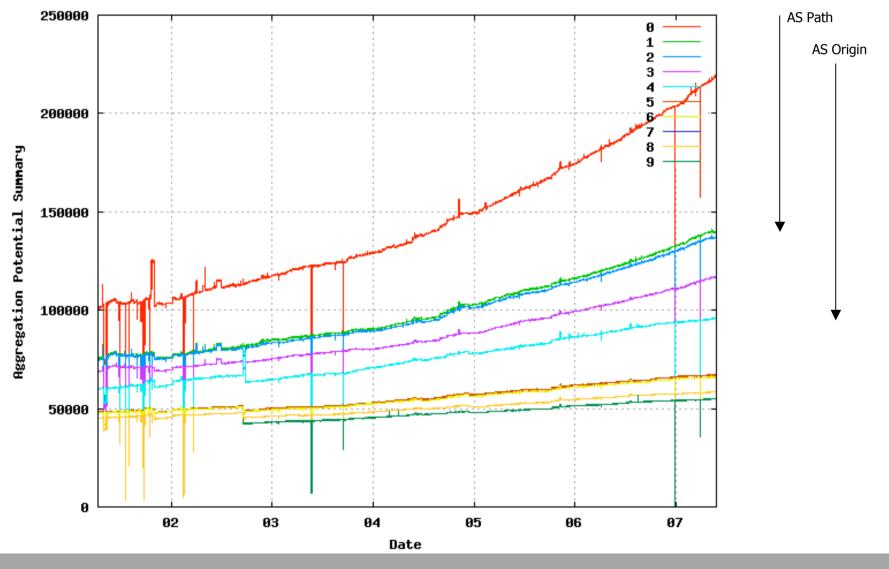
4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 3356 18566

4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 3356 18566

4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 3356 18566

4608 1221 4637 3356 18566 - Withdrawn - matching aggregate 64.105.0.0/16 4608 1221 4637 3356 18566

## Aggregation Potential (source: bgp.potaroo.net/as4637/)



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### Aggregation Summary

Aggregation on the Internet could be MUCH better

35% saving on Internet routing table size is quite feasible

Tools are available

Commands on the routers are not hard

**CIDR-Report webpage** 

#### RIPE Routing WG aggregation recommendation

RIPE-399 — http://www.ripe.net/ripe/docs/ripe-399.html



# **Receiving Prefixes**

## **Receiving Prefixes**

- There are three scenarios for receiving prefixes from other ASNs
  - Customer talking BGP
  - Peer talking BGP
  - **Upstream/Transit talking BGP**
- Each has different filtering requirements and need to be considered separately

### **Receiving Prefixes: From Customers**

- ISPs should only accept prefixes which have been assigned or allocated to their downstream customer
- If ISP has assigned address space to its customer, then the customer IS entitled to announce it back to his ISP
- If the ISP has NOT assigned address space to its customer, then:

Check the five RIR databases to see if this address space really has been assigned to the customer

The tool: whois

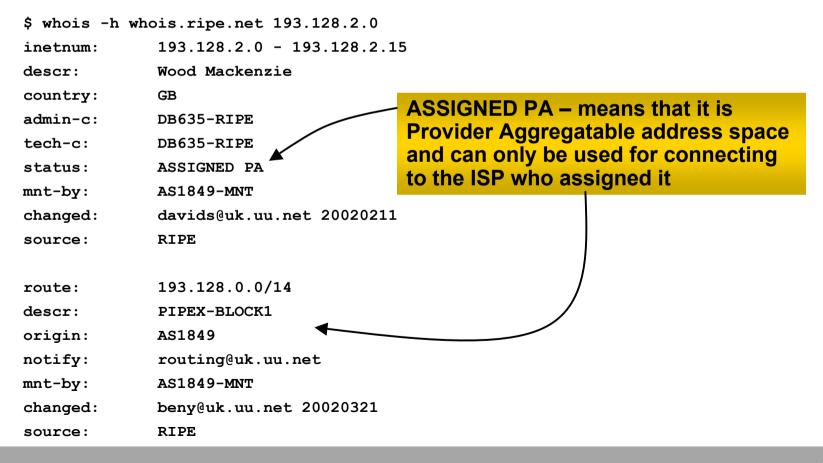
#### **Receiving Prefixes: From Customers**

• Example use of whois to check if customer is entitled to announce address space:

pfs-pc\$ whois	-h whois.apnic.ne	et 202.12.29.0
inetnum:	202.12.29.0 - 202	2.12.29.255
netname:	APNIC-AP-AU-BNE	
descr:	APNIC Pty Ltd - H	Brisbane Offices + Servers
descr:	Level 1, 33 Park	Rd
descr:	PO Box 2131, Milt	con
descr:	Brisbane, QLD.	
country:	AU	Portable – means its an assignment
admin-c:	HM20-AP	to the customer, the customer can
tech-c:	NO4-AP	announce it to you
tech-c: mnt-by:	NO4-AP APNIC-HM	announce it to you
	-	
mnt-by:	APNIC-HM	net 20030108

#### **Receiving Prefixes: From Customers**

 Example use of whois to check if customer is entitled to announce address space:



#### **Receiving Prefixes: From Peers**

• A peer is an ISP with whom you agree to exchange prefixes you originate into the Internet routing table

Prefixes you accept from a peer are only those they have indicated they will announce

Prefixes you announce to your peer are only those you have indicated you will announce

### **Receiving Prefixes: From Peers**

Agreeing what each will announce to the other:

Exchange of e-mail documentation as part of the peering agreement, and then ongoing updates

OR

Use of the Internet Routing Registry and configuration tools such as the IRRToolSet

www.isc.org/sw/IRRToolSet/

### **Receiving Prefixes: From Upstream/Transit Provider**

- Upstream/Transit Provider is an ISP who you pay to give you transit to the WHOLE Internet
- Receiving prefixes from them is not desirable unless really necessary

special circumstances – see later

• Ask upstream/transit provider to either:

originate a default-route

OR

announce one prefix you can use as default

#### **Receiving Prefixes: From Upstream/Transit Provider**

- If necessary to receive prefixes from any provider, care is required
  - don't accept RFC1918 etc prefixes
    - ftp://ftp.rfc-editor.org/in-notes/rfc3330.txt
  - don't accept your own prefixes
  - don't accept default (unless you need it)
  - don't accept prefixes longer than /24
- Check Rob Thomas' list of "bogons"

http://www.cymru.com/Documents/bogon-list.html

## **Receiving Prefixes**

 Paying attention to prefixes received from customers, peers and transit providers assists with:

The integrity of the local network

The integrity of the Internet

 Responsibility of all ISPs to be good Internet citizens



# **Preparing the network**

Before we begin...

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## **Preparing the Network**

- We will deploy BGP across the network before we try and multihome
- BGP will be used therefore an ASN is required
- If multihoming to different ISPs, public ASN needed:

Either go to upstream ISP who is a registry member, or

Apply to the RIR yourself for a one off assignment, or

Ask an ISP who is a registry member, or

Join the RIR and get your own IP address allocation too

(this option strongly recommended)!

### Preparing the Network Initial Assumptions

- The network is not running any BGP at the moment single statically routed connection to upstream ISP
- The network is not running any IGP at all

Static default and routes through the network to do "routing"

### Preparing the Network First Step: IGP

- Decide on an IGP: OSPF or ISIS ③
- Assign loopback interfaces and /32 address to each router which will run the IGP

Loopback is used for OSPF and BGP router id anchor

Used for iBGP and route origination

Deploy IGP (e.g. OSPF)

IGP can be deployed with NO IMPACT on the existing static routing

e.g. OSPF distance might be 110m static distance is 1

**Smallest distance wins** 

## Preparing the Network IGP (cont)

- Be prudent deploying IGP keep the Link State Database Lean!
  - **Router loopbacks go in IGP**
  - WAN point to point links go in IGP
  - (In fact, any link where IGP dynamic routing will be run should go into IGP)
  - Summarise on area/level boundaries (if possible) i.e. think about your IGP address plan

## Preparing the Network IGP (cont)

Routes which don't go into the IGP include:

**Dynamic assignment pools (DSL/Cable/Dial)** 

**Customer point to point link addressing** 

(using next-hop-self in iBGP ensures that these do NOT need to be in IGP)

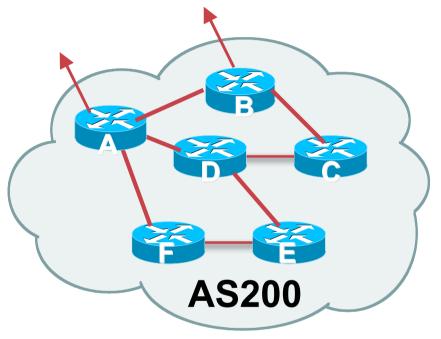
**Static/Hosting LANs** 

**Customer assigned address space** 

Anything else not listed in the previous slide

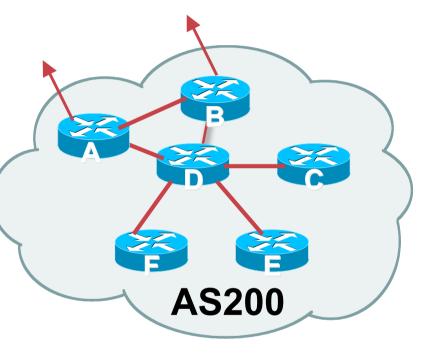
#### Preparing the Network Second Step: iBGP

- Second step is to configure the local network to use iBGP
- iBGP can run on
  - all routers, or
  - a subset of routers, or
  - just on the upstream edge
- *iBGP must run on all routers which are in the transit path between external connections*



#### Preparing the Network Second Step: iBGP (Transit Path)

- iBGP must run on all routers which are in the transit path between external connections
- Routers C, E and F are not in the transit path
  - Static routes or IGP will suffice
- Router D is in the transit path
  - Will need to be in iBGP mesh, otherwise routing loops will result



## Preparing the Network Layers

• Typical SP networks have three layers:

Core – the backbone, usually the transit path Distribution – the middle, PoP aggregation layer Aggregation – the edge, the devices connecting customers

## Preparing the Network Aggregation Layer

#### iBGP is optional

Many ISPs run iBGP here, either partial routing (more common) or full routing (less common)

Full routing is not needed unless customers want full table

Partial routing is cheaper/easier, might usually consist of internal prefixes and, optionally, external prefixes to aid external load balancing

Communities and peer-groups make this administratively easy

Many aggregation devices can't run iBGP

Static routes from distribution devices for address pools IGP for best exit

### Preparing the Network Distribution Layer

#### Usually runs iBGP

Partial or full routing (as with aggregation layer)

#### But does not have to run iBGP

IGP is then used to carry customer prefixes (does not scale)

**IGP** is used to determine nearest exit

#### Networks which plan to grow large should deploy iBGP from day one

Migration at a later date is extra work

No extra overhead in deploying iBGP, indeed IGP benefits

## Preparing the Network Core Layer

- Core of network is usually the transit path
- iBGP necessary between core devices

Full routes or partial routes:

**Transit ISPs carry full routes in core** 

Edge ISPs carry partial routes only

• Core layer includes AS border routers

Decide on:

Best iBGP policy

Will it be full routes everywhere, or partial, or some mix?

iBGP scaling technique

Community policy?

**Route-reflectors?** 

Techniques such as peer groups and peer templates?

#### • Then deploy iBGP:

Step 1: Introduce iBGP mesh on chosen routers

make sure that iBGP distance is greater than IGP distance (it usually is)

Step 2: Install "customer" prefixes into iBGP

**Check!** Does the network still work?

Step 3: Carefully remove the static routing for the prefixes now in IGP and iBGP

**Check!** Does the network still work?

**Step 4: Deployment of eBGP follows** 

#### Install "customer" prefixes into iBGP?

Customer assigned address space

**Network statement/static route combination** 

Use unique community to identify customer assignments

#### Customer facing point-to-point links

Redistribute connected through filters which only permit point-topoint link addresses to enter iBGP

Use a unique community to identify point-to-point link addresses (these are only required for your monitoring system)

#### Dynamic assignment pools & local LANs

Simple network statement will do this

Use unique community to identify these networks

Carefully remove static routes?

- Work on one router at a time:
  - Check that static route for a particular destination is also learned by the iBGP

If so, remove it

If not, establish why and fix the problem

(Remember to look in the RIB, not the FIB!)

- Then the next router, until the whole PoP is done
- Then the next PoP, and so on until the network is now dependent on the IGP and iBGP you have deployed

### Preparing the Network Completion

#### • Previous steps are NOT flag day steps

Each can be carried out during different maintenance periods, for example:

Step One on Week One

Step Two on Week Two

**Step Three on Week Three** 

And so on

And with proper planning will have NO customer visible impact at all

#### Preparing the Network Configuration Summary

- IGP essential networks are in IGP
- Customer networks are now in iBGP iBGP deployed over the backbone
   Full or Partial or Upstream Edge only
- BGP distance is greater than any IGP
- Now ready to deploy eBGP



# **Configuration Tips**

Of passwords, tricks and templates

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#### iBGP and IGPs Reminder!

- Make sure loopback is configured on router iBGP between loopbacks, NOT real interfaces
- Make sure IGP carries loopback /32 address
- Consider the DMZ nets:

Use unnumbered interfaces?

- Use next-hop-self on iBGP neighbours
- Or carry the DMZ /30s in the iBGP
- Basically keep the DMZ nets out of the IGP!

## iBGP: Next-hop-self

- BGP speaker announces external network to iBGP peers using router's local address (loopback) as next-hop
- Used by many ISPs on edge routers

Preferable to carrying DMZ /30 addresses in the IGP

**Reduces size of IGP to just core infrastructure** 

Alternative to using unnumbered interfaces

Helps scale network

Many ISPs consider this "best practice"

# Limiting AS Path Length

 Some BGP implementations have problems with long AS\_PATHS

Memory corruption

Memory fragmentation

 Even using AS\_PATH prepends, it is not normal to see more than 20 ASes in a typical AS\_PATH in the Internet today

The Internet is around 5 ASes deep on average

Largest AS\_PATH is usually 16-20 ASNs

## Limiting AS Path Length

 Some announcements have ridiculous lengths of AS-paths:

\*> 3FFE:1600::/24 22 11537 145 12199 10318 10566
13193 1930 2200 3425 293 5609 5430 13285 6939 14277 1849 33
15589 25336 6830 8002 2042 7610 i

#### This example is an error in one IPv6 implementation

\*> 194.146.180.0/22 2497 3257 29686 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 1

This example shows 20 prepends (for no obvious reason)

 If your implementation supports it, consider limiting the maximum AS-path length you will accept

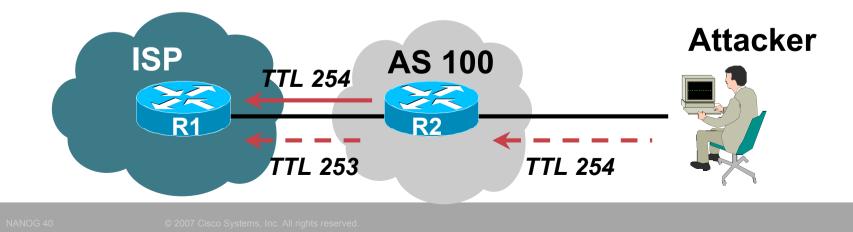
## **BGP TTL "hack"**

#### Implement RFC3682 on BGP peerings

Neighbour sets TTL to 255

Local router expects TTL of incoming BGP packets to be 254

No one apart from directly attached devices can send BGP packets which arrive with TTL of 254, so any possible attack by a remote miscreant is dropped due to TTL mismatch



## **BGP TTL "hack"**

#### • TTL Hack:

Both neighbours must agree to use the feature TTL check is much easier to perform than MD5 (Called BTSH – BGP TTL Security Hack)

#### Provides "security" for BGP sessions

In addition to packet filters of course

MD5 should still be used for messages which slip through the TTL hack

See www.nanog.org/mtg-0302/hack.html for more details

## **Templates**

Good practice to configure templates for everything

Vendor defaults tend not to be optimal or even very useful for ISPs

ISPs create their own defaults by using configuration templates

eBGP and iBGP examples follow

Also see Project Cymru's BGP templates

www.cymru.com/Documents

#### **iBGP Template** Example

- iBGP between loopbacks!
- Next-hop-self

Keep DMZ and external point-to-point out of IGP

#### Always send communities in iBGP

Otherwise accidents will happen

#### Hardwire BGP to version 4

Yes, this is being paranoid!

#### **iBGP Template** Example continued

#### Use passwords on iBGP session

Not being paranoid, VERY necessary

It's a secret shared between you and your peer

If arriving packets don't have the correct MD5 hash, they are ignored

Helps defeat miscreants who wish to attack BGP sessions

 Powerful preventative tool, especially when combined with filters and the TTL "hack"

## eBGP Template Example

#### BGP damping

Do NOT use it unless you understand the impact Do NOT use the vendor defaults without thinking

# Remove private ASes from announcements

**Common omission today** 

Use extensive filters, with "backup"

Use as-path filters to backup prefix filters

Keep policy language for implementing policy, rather than basic filtering

 Use password agreed between you and peer on eBGP session

#### **eBGP Template** Example continued

Use maximum-prefix tracking

Router will warn you if there are sudden increases in BGP table size, bringing down eBGP if desired

- Limit maximum as-path length inbound
- Log changes of neighbour state

...and monitor those logs!

Make BGP admin distance higher than that of any IGP

Otherwise prefixes heard from outside your network could override your IGP!!

#### Summary

- Use configuration templates
- Standardise the configuration
- Be aware of standard "tricks" to avoid compromise of the BGP session
- Anything to make your life easier, network less prone to errors, network more likely to scale
- It's all about scaling if your network won't scale, then it won't be successful



# BGP Techniques for Internet Service Providers

Philip Smith <pfs@cisco.com>NANOG 403-6 June 2007Bellevue, Washington