# BGP for Internet Service Providers Philip Smith <pfs@cisco.com> NANOG 22, Scottsdale, Arizona

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## BGP for Internet Service Providers

- BGP Basics (quick recap)
- Scaling BGP
- Deploying BGP in an ISP network
- Trouble & Troubleshooting
- Multihoming Examples
- Using Communities

## **BGP Basics**

## What is this BGP thing?



#### **Border Gateway Protocol**

 Routing Protocol used to exchange routing information between networks

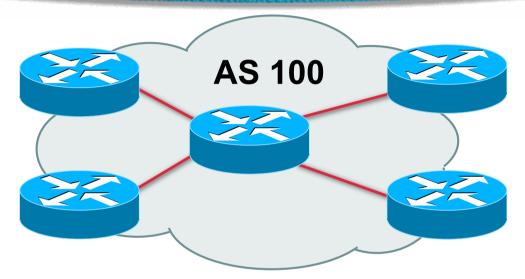
exterior gateway protocol

• RFC1771

work in progress to update

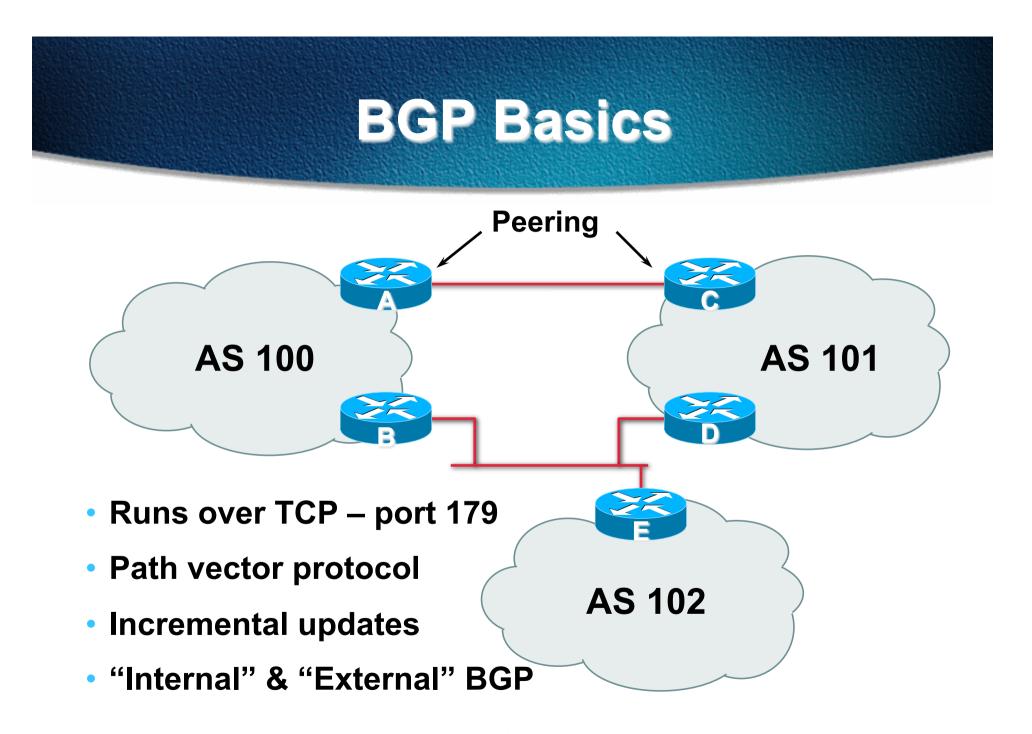
draft-ietf-idr-bgp4-12.txt

#### Autonomous System (AS)

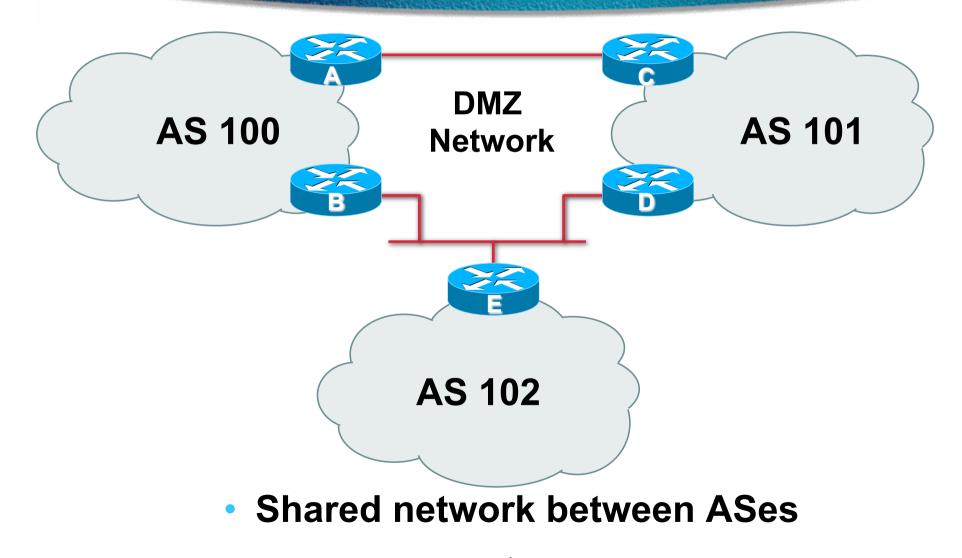


- Collection of networks with same routing policy
- Single routing protocol
- Usually under single ownership, trust and administrative control

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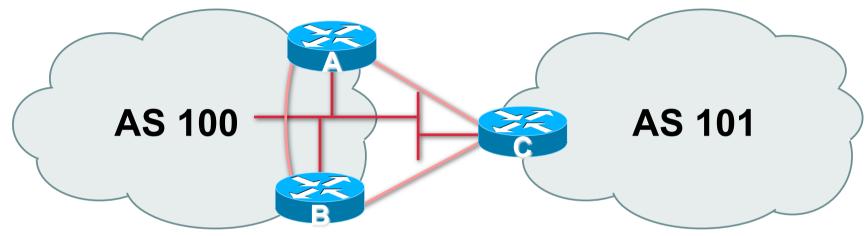
#### **Demarcation Zone (DMZ)**



#### **BGP General Operation**

- Learns multiple paths via internal and external BGP speakers
- Picks the best path and installs in the forwarding table
- Best path is sent to external BGP neighbours
- Policies applied by influencing the best path selection





- Between BGP speakers in different AS
- Should be directly connected
- Never run an IGP between eBGP peers

#### **Configuring External BGP**

#### **Router A in AS100**

interface ethernet 5/0
ip address 222.222.10.2 255.255.255.240
router bgp 100
network 220.220.8.0 mask 255.255.252.0
neighbor 222.222.10.1 remote-as 101
neighbor 222.222.10.1 prefix-list RouterC in
neighbor 222.222.10.1 prefix-list RouterC out

#### Router C in AS101

interface ethernet 1/0/0
ip address 222.222.10.1 255.255.255.240
router bgp 101
network 220.220.16.0 mask 255.255.240.0
neighbor 222.222.10.2 remote-as 100
neighbor 222.222.10.2 prefix-list RouterA in
neighbor 222.222.10.2 prefix-list RouterA out
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## Internal BGP (iBGP)

- BGP peer within the same AS
- Not required to be directly connected
- iBGP speakers need to be fully meshed they originate connected networks
   they do not pass on prefixes learned from other iBGP speakers

## Internal BGP Peering (iBGP)

**AS 100** 

- Topology independent
- Each iBGP speaker must peer with every other iBGP speaker in the AS

#### Peering to Loop-Back Address

**AS 100** 

Peer with loop-back address

Loop-back interface does not go down - ever!

- iBGP session is not dependent on state of a single interface
- iBGP session is not dependent on physical topology

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#### **Configuring Internal BGP**

#### **Router A**

```
interface loopback 0
ip address 215.10.7.1 255.255.255.255
router bgp 100
network 220.220.1.0
neighbor 215.10.7.2 remote-as 100
neighbor 215.10.7.2 update-source loopback0
neighbor 215.10.7.3 remote-as 100
neighbor 215.10.7.3 update-source loopback0
```

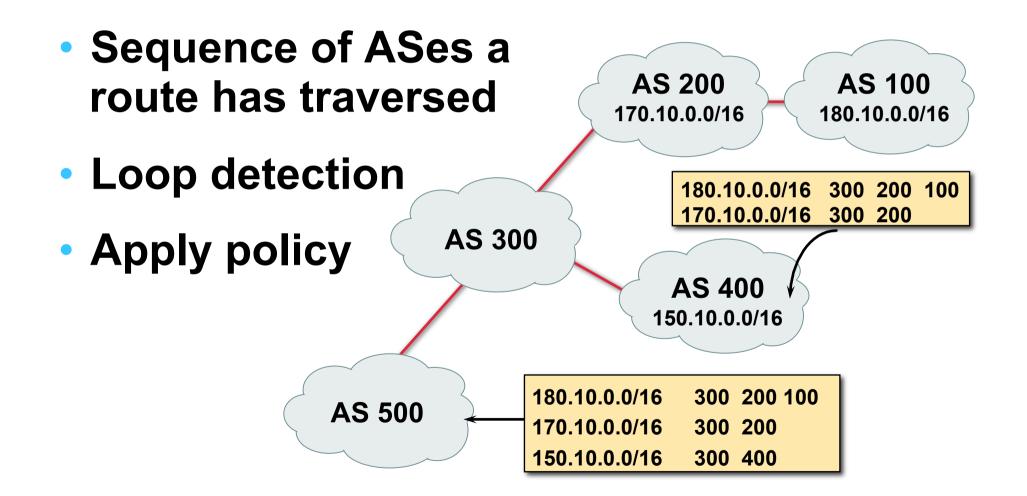
#### **Router B**

```
interface loopback 0
ip address 215.10.7.2 255.255.255.255
router bgp 100
network 220.220.5.0
neighbor 215.10.7.1 remote-as 100
neighbor 215.10.7.1 update-source loopback0
neighbor 215.10.7.3 remote-as 100
neighbor 215.10.7.3 update-source loopback0
```

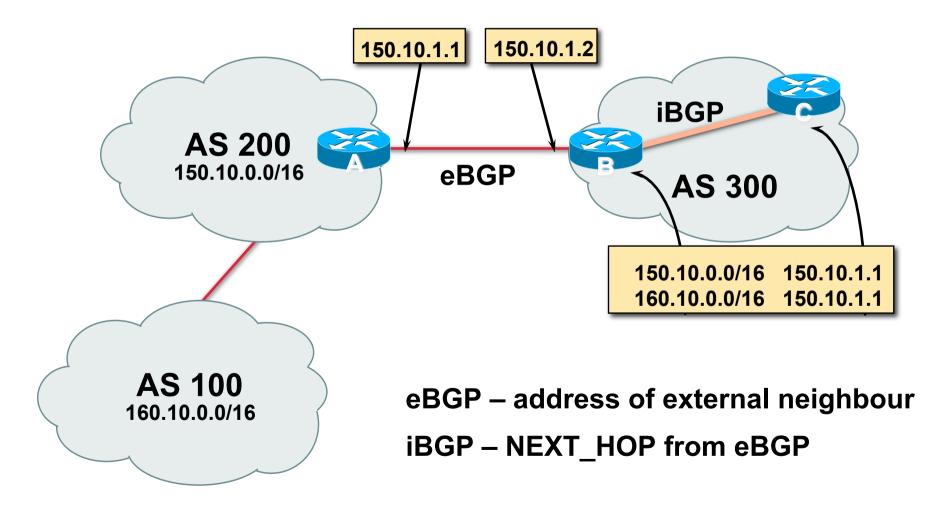
# **BGP Attributes**

Recap

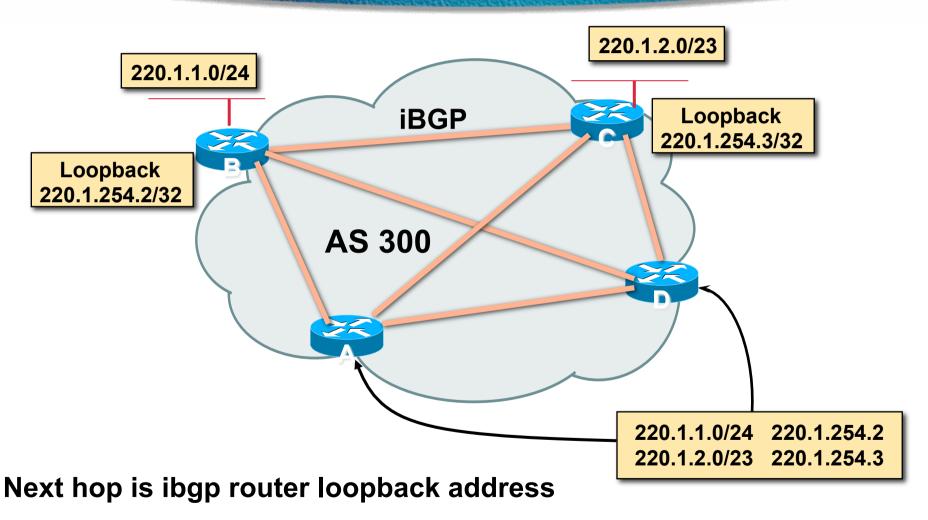
#### **AS-Path**



#### **Next Hop**



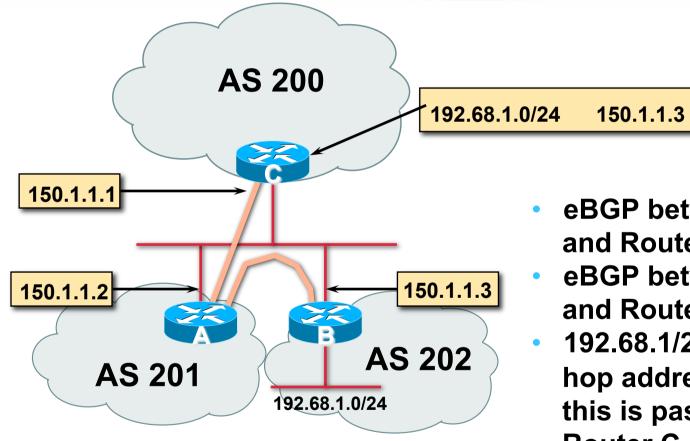
#### **iBGP Next Hop**



#### **Recursive route look-up**

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#### **Third Party Next Hop**



- eBGP between Router A and Router C
- eBGP between Router A and Router B
- 192.68.1/24 prefix has next hop address of 150.1.1.3 – this is passed on to Router C instead of 150.1.1.2

## Next Hop (summary)

- IGP should carry route to next hops
- Recursive route look-up
- Unlinks BGP from actual physical topology
- Allows IGP to make intelligent forwarding decision

## Origin

- Conveys the origin of the prefix
- "Historical" attribute
- Influences best path selection
- Three values: IGP, EGP, incomplete
  - **IGP generated by BGP network statement**

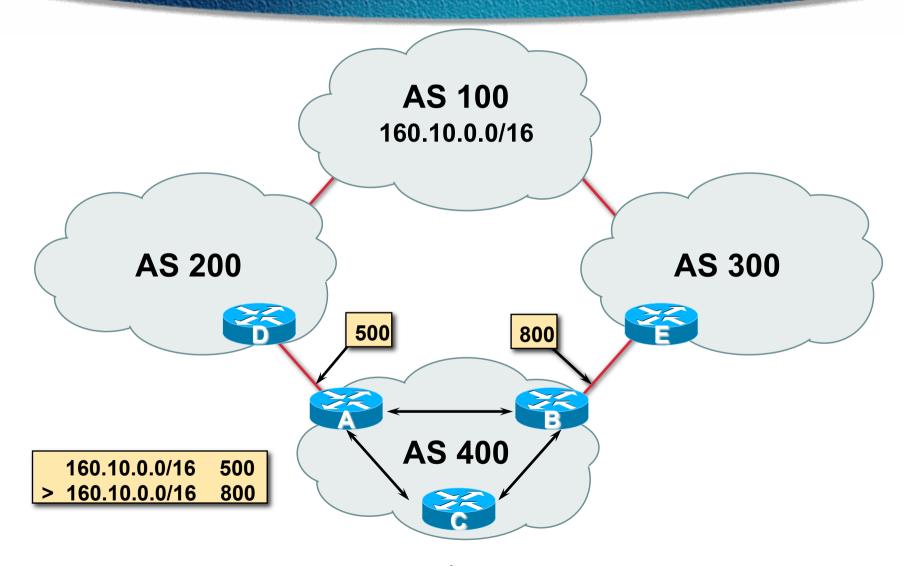
EGP – generated by EGP

incomplete – redistributed from another routing protocol

## Aggregator

- Conveys the IP address of the router/BGP speaker generating the aggregate route
- Useful for debugging purposes
- Does not influence best path selection

#### **Local Preference**



#### Local Preference

- Local to an AS non-transitive
   Default local preference is 100
- Used to influence BGP path selection

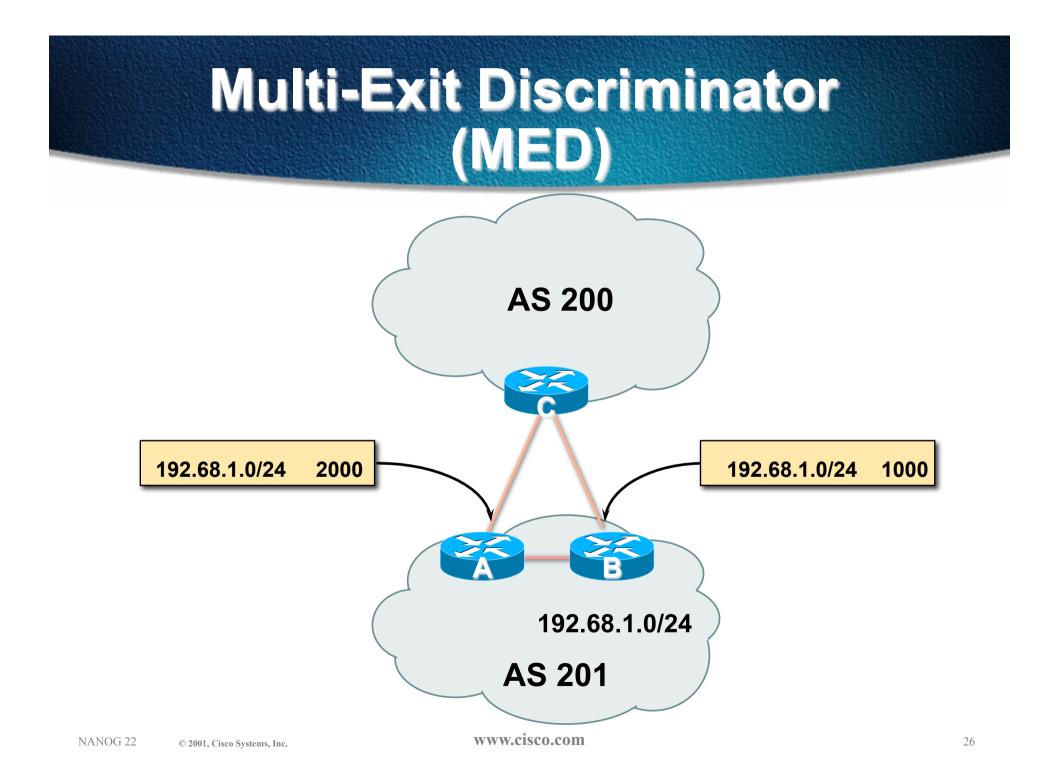
determines best path for *outbound* traffic

 Path with highest local preference wins

#### Local Preference

#### Configuration of Router B:

```
router bgp 400
neighbor 220.5.1.1 remote-as 300
neighbor 220.5.1.1 route-map local-pref in
!
route-map local-pref permit 10
match ip address prefix-list MATCH
set local-preference 800
!
ip prefix-list MATCH permit 160.10.0.0/16
```



#### **Multi-Exit Discriminator**

- Inter-AS non-transitive
- Used to convey the relative preference of entry points

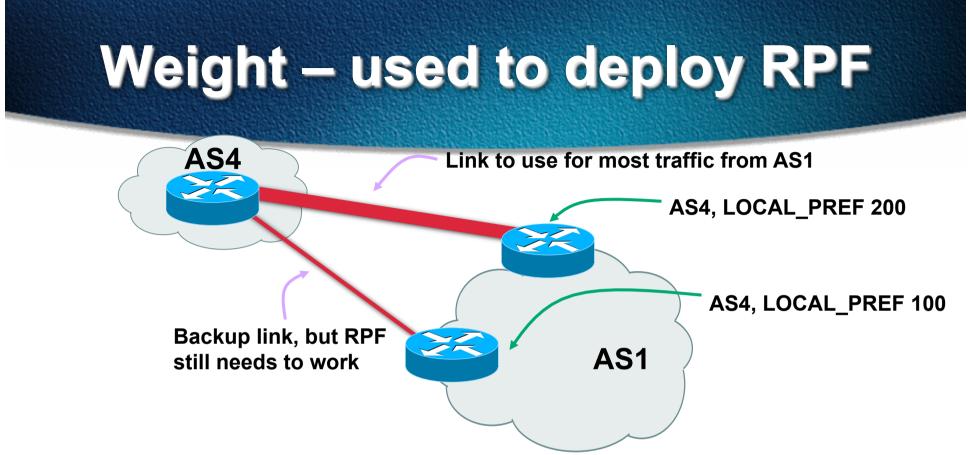
determines best path for *inbound* traffic

- Comparable if paths are from same AS
- IGP metric can be conveyed as MED set metric-type internal in route-map

#### **Multi-Exit Discriminator**

#### Configuration of Router B:

```
router bgp 400
neighbor 220.5.1.1 remote-as 200
neighbor 220.5.1.1 route-map set-med out
!
route-map set-med permit 10
match ip address prefix-list MATCH
set metric 1000
!
ip prefix-list MATCH permit 192.68.1.0/24
```



- Local to router on which it's configured Not really an attribute
- route-map: set weight
- Highest weight wins over all valid paths
- Weight customer eBGP on edge routers to allow RPF to work correctly

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#### Community

- BGP attribute
- Described in RFC1997
- 32 bit integer

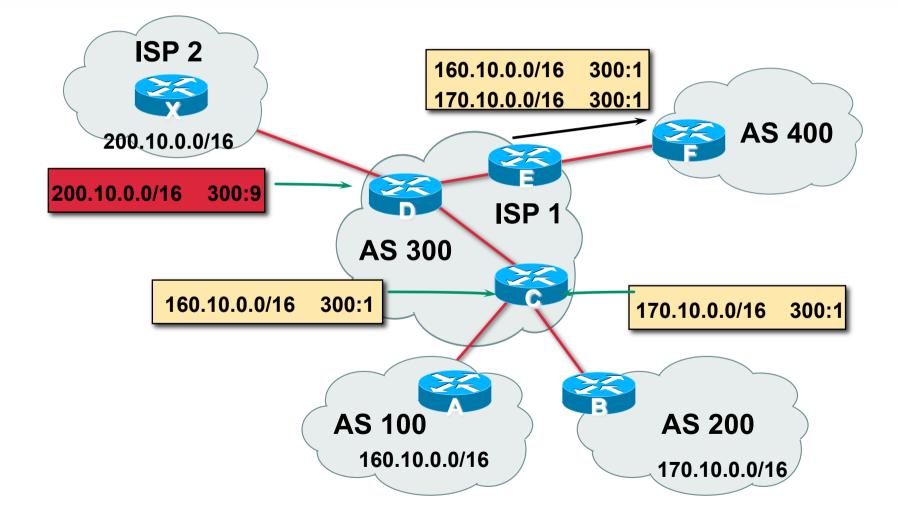
**Represented as two 16bit integers** 

Used to group destinations

Each destination could be member of multiple communities

- Community attribute carried across AS's
- Very useful in applying policies

#### Community



#### **Well-Known Communities**

#### no-export

do not advertise to eBGP peers

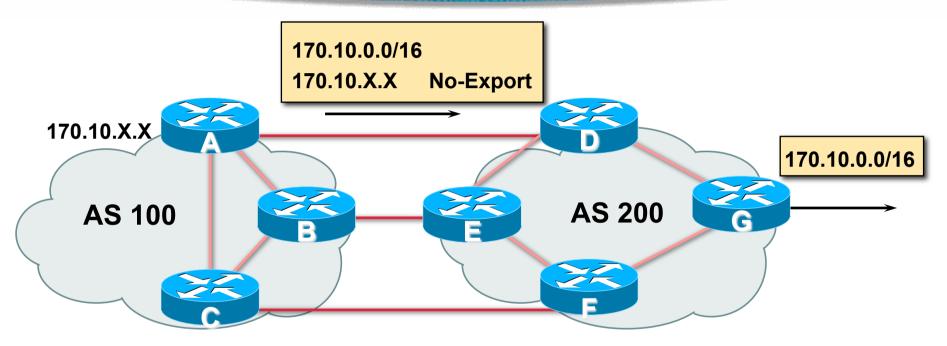
#### no-advertise

do not advertise to any peer

#### local-AS

do not advertise outside local AS (only used with confederations)

### **No-Export Community**



- AS100 announces aggregate and subprefixes aim is to improve loadsharing by leaking subprefixes
- Subprefixes marked with no-export community
- Router G in AS200 strips out all prefixes with no-export community set

# BGP Path Selection Algorithm

Why is this the best path?

## BGP Path Selection Algorithm

- Do not consider path if no route to next hop
- Do not consider iBGP path if not synchronised (Cisco IOS)
- Highest weight (local to router)
- Highest local preference (global within AS)
- Prefer locally originated route
- Shortest AS path

## BGP Path Selection Algorithm (continued)

- Lowest origin code
   IGP < EGP < incomplete</li>
- Lowest Multi-Exit Discriminator (MED)

If **bgp deterministic-med**, order the paths before comparing

If bgp always-compare-med, then compare for all paths

otherwise MED only considered if paths are from the same AS (default)

## BGP Path Selection Algorithm (continued)

- Prefer eBGP path over iBGP path
- Path with lowest IGP metric to next-hop
- Lowest router-id (originator-id for reflected routes)
- Shortest Cluster-List

Client must be aware of Route Reflector attributes!

Lowest neighbour IP address

## Applying Policy with BGP

**Control!** 

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## **Applying Policy with BGP**

Applying Policy

Decisions based on AS path, community or the prefix

**Rejecting/accepting selected routes** 

Set attributes to influence path selection

• Tools:

**Prefix-list (filter prefixes)** 

**Filter-list (filter ASes)** 

**Route-maps and communities** 

### Policy Control Prefix List

Filter routes based on prefix

### Inbound and Outbound

```
router bgp 200
neighbor 220.200.1.1 remote-as 210
neighbor 220.200.1.1 prefix-list PEER-IN in
neighbor 220.200.1.1 prefix-list PEER-OUT out
!
ip prefix-list PEER-IN deny 218.10.0.0/16
ip prefix-list PEER-IN permit 0.0.0.0/0 le 32
ip prefix-list PEER-OUT permit 215.7.0.0/16
```

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## Policy Control Filter List

### Filter routes based on AS path

### Inbound and Outbound

```
router bgp 100
neighbor 220.200.1.1 remote-as 210
neighbor 220.200.1.1 filter-list 5 out
neighbor 220.200.1.1 filter-list 6 in
!
ip as-path access-list 5 permit ^200$
ip as-path access-list 6 permit ^150$
```

## Policy Control Regular Expressions

### • Like Unix regular expressions

- Match one character
- \* Match any number of preceding expression
- + Match at least one of preceding expression
- A Beginning of line
- \$ End of line
  - Beginning, end, white-space, brace
  - Or
  - brackets to contain expression

## Policy Control Regular Expressions

#### Simple Examples

- Match anything
- .+ Match at least one character
- **^\$** Match routes local to this AS
  - Originated by 1800
  - Received from 1800
    - Via 1800
    - Passing through 1800 then 790
    - Match at least one of 1800 in sequence
    - Via 65350 (confederation AS)

1800\$

^1800

1800

\_790\_1800\_

\_(1800\_)+

\_\(65350\)\_

\*

- A route-map is like a "programme" for IOS
- Has "line" numbers, like programmes
- Each line is a separate condition/action
- Concept is basically:

if *match* then do *expression* and *exit* else

if *match* then do *expression* and *exit* else *etc* 

#### Example using prefix-lists

```
router bgp 100
neighbor 1.1.1.1 route-map infilter in
route-map infilter permit 10
match ip address prefix-list HIGH-PREF
set local-preference 120
route-map infilter permit 20
match ip address prefix-list LOW-PREF
set local-preference 80
route-map infilter permit 30
ip prefix-list HIGH-PREF permit 10.0.0/8
ip prefix-list LOW-PREF permit 20.0.0/8
```

Example using filter lists

```
router bgp 100
neighbor 220.200.1.2 route-map filter-on-as-path in
route-map filter-on-as-path permit 10
match as-path 1
 set local-preference 80
route-map filter-on-as-path permit 20
match as-path 2
 set local-preference 200
route-map filter-on-as-path permit 30
ip as-path access-list 1 permit 150$
ip as-path access-list 2 permit 210
```

#### Example configuration of AS-PATH prepend

router bgp 300
network 215.7.0.0
neighbor 2.2.2.2 remote-as 100
neighbor 2.2.2.2 route-map SETPATH out
!
route-map SETPATH permit 10
set as-path prepend 300 300

 Use your own AS number when prepending Otherwise BGP loop detection may cause disconnects

## Policy Control Setting Communities

#### Example Configuration

```
router bgp 100
 neighbor 220.200.1.1 remote-as 200
 neighbor 220.200.1.1 send-community
neighbor 220.200.1.1 route-map set-community out
1
route-map set-community permit 10
match ip address prefix-list NO-ANNOUNCE
 set community no-export
route-map set-community permit 20
ip prefix-list NO-ANNOUNCE permit 172.168.0.0/16 ge 17
```

## Policy Control Matching Communities

#### Example Configuration

```
router bgp 100
neighbor 220.200.1.2 remote-as 200
neighbor 220.200.1.2 route-map filter-on-community in
ļ
route-map filter-on-community permit 10
match community 1
set local-preference 50
route-map filter-on-community permit 20
match community 2 exact-match
set local-preference 200
ip community-list 1 permit 150:3 200:5
ip community-list 2 permit 88:6
```

## **BGP Capabilities**

### **Extending BGP**

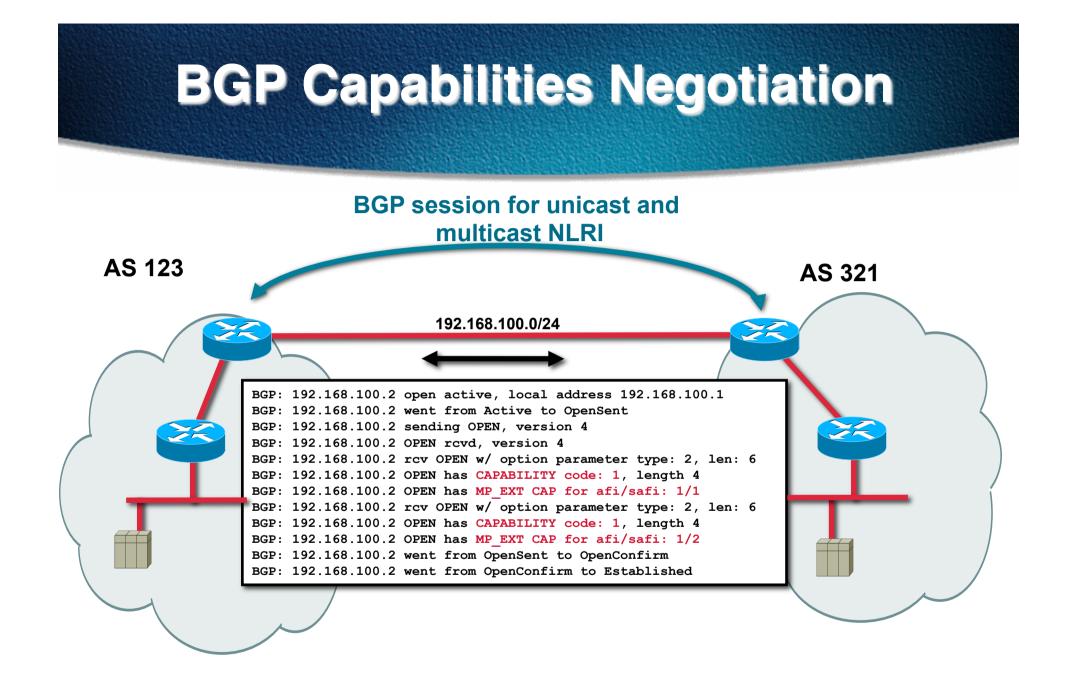
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### **BGP Capabilities**

- Documented in RFC2842
- Capabilities parameters passed in BGP open message
- Unknown or unsupported capabilities will result in NOTIFICATION message

### Current capabilities are:

0	Reserved	[RFC2842]
1	Multiprotocol Extensions for BGP-4	[RFC2858]
2	Route Refresh Capability for BGP-4	[RFC2918]
3	Cooperative Route Filtering Capability	[]
4	Multiple routes to a destination capability	[RFC3107]
64	Graceful Restart Capability	[]



### BGP for Internet Service Providers

- BGP Basics (quick recap)
- Scaling BGP
- Deploying BGP in an ISP network
- Trouble & Troubleshooting
- Multihoming Examples
- Using Communities

## **BGP Scaling Techniques**



### **BGP Scaling Techniques**

- How to scale iBGP mesh beyond a few peers?
- How to implement new policy without causing flaps and route churning?
- How to reduce the overhead on the routers?
- How to keep the network stable, scalable, as well as simple?

### **BGP Scaling Techniques**

- Dynamic Reconfiguration
- Peer groups
- Route flap damping
- Route Reflectors & Confederations

## Dynamic Reconfiguration Soft Reconfiguration and Route Refresh

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### **Soft Reconfiguration**

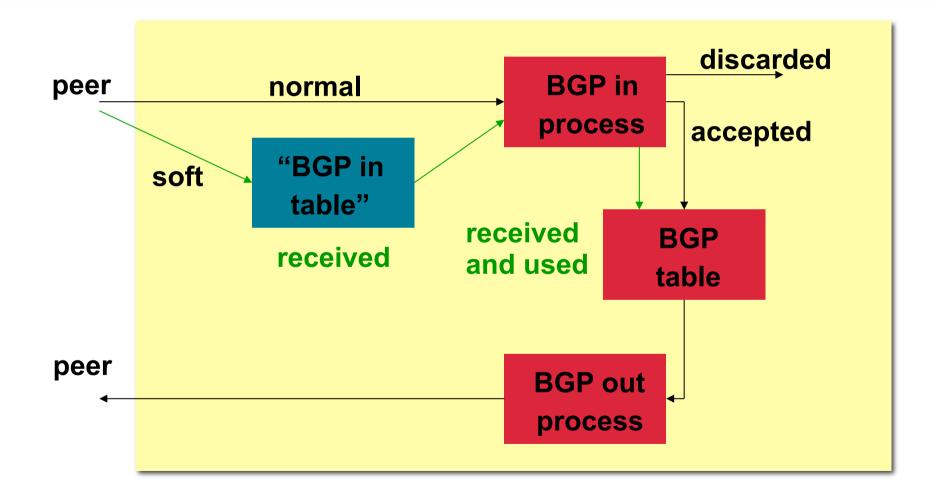
#### **Problem:**

- Hard BGP peer clear required after every policy change because the router does not store prefixes that are denied by a filter
- Hard BGP peer clearing consumes CPU and affects connectivity for all networks

### Solution:

Soft-reconfiguration

### **Soft Reconfiguration**



### **Soft Reconfiguration**

- New policy is activated without tearing down and restarting the peering session
- Per-neighbour basis
- Use more memory to keep prefixes whose attributes have been changed or have not been accepted

# Configuring Soft reconfiguration

```
router bgp 100
```

```
neighbor 1.1.1.1 remote-as 101
```

```
neighbor 1.1.1.1 route-map infilter in
```

```
neighbor 1.1.1.1 soft-reconfiguration inbound
```

! Outbound does not need to be configured !

Then when we change the policy, we issue an exec command

clear ip bgp 1.1.1.1 soft [in | out]

### **Route Refresh Capability**

- Facilitates non-disruptive policy changes
- No configuration is needed
- No additional memory is used
- Requires peering routers to support "route refresh capability" – RFC2918
- clear ip bgp x.x.x.x in tells peer to resend full BGP announcement

## Soft Reconfiguration vs Route Refresh

- Use Route Refresh capability if supported find out from "show ip bgp neighbor" uses much less memory
- Otherwise use Soft Reconfiguration
- Only hard-reset a BGP peering as a last resort

## **Peer Groups**

### **Peer Groups**

Without peer groups

- iBGP neighbours receive same update
- Large iBGP mesh slow to build
- Router CPU wasted on repeat calculations
   Solution peer groups!
- Group peers with same outbound policy
- Updates are generated once per group

### Peer Groups – Advantages

- Makes configuration easier
- Makes configuration less prone to error
- Makes configuration more readable
- Lower router CPU load
- iBGP mesh builds more quickly
- Members can have different inbound policy
- Can be used for eBGP neighbours too!

### **Configuring Peer Group**

```
router bgp 100
neighbor ibgp-peer peer-group
neighbor ibgp-peer remote-as 100
neighbor ibgp-peer update-source loopback 0
neighbor ibgp-peer send-community
neighbor ibgp-peer route-map outfilter out
neighbor 1.1.1.1 peer-group ibgp-peer
neighbor 2.2.2.2 peer-group ibgp-peer
neighbor 2.2.2.2 route-map infilter in
neighbor 3.3.3.3 peer-group ibgp-peer
```

! note how 2.2.2.2 has different inbound filter from peer-group !

### **Configuring Peer Group**

```
router bgp 109
neighbor external-peer peer-group
neighbor external-peer send-community
neighbor external-peer route-map set-metric out
neighbor 160.89.1.2 remote-as 200
neighbor 160.89.1.2 peer-group external-peer
neighbor 160.89.1.4 remote-as 300
neighbor 160.89.1.4 peer-group external-peer
neighbor 160.89.1.6 remote-as 400
neighbor 160.89.1.6 peer-group external-peer
neighbor 160.89.1.6 filter-list infilter in
```

## **Route Flap Damping**

### **Stabilising the Network**

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### **Route Flap Damping**

- Route flap
  - Going up and down of path or change in attribute
    - **BGP WITHDRAW** followed by UPDATE = 1 flap
    - eBGP neighbour going down/up is NOT a flap
  - **Ripples through the entire Internet**
  - Wastes CPU
- Damping aims to reduce scope of route flap propagation

## Route Flap Damping (Continued)

### Requirements

Fast convergence for normal route changes

**History predicts future behaviour** 

**Suppress oscillating routes** 

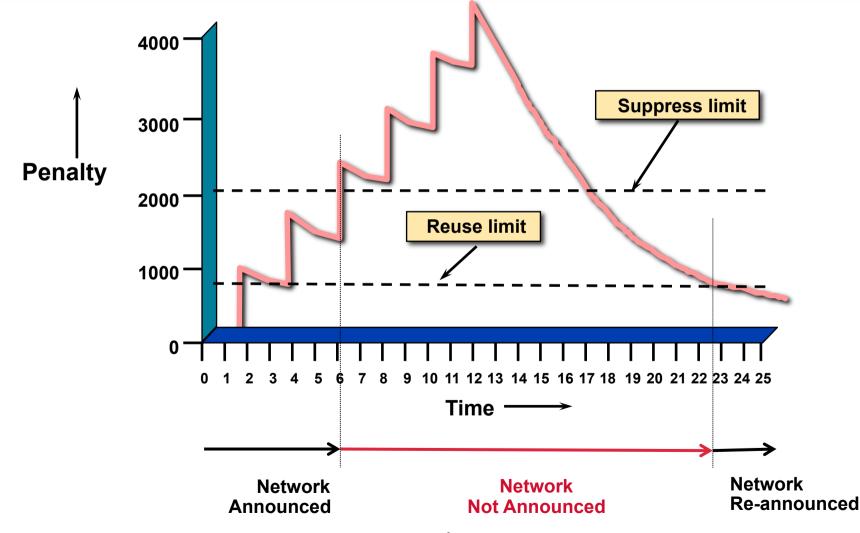
**Advertise stable routes** 

Documented in RFC2439

### Operation

- Add penalty (1000) for each flap
   Change in attribute gets penalty of 500
- Exponentially decay penalty half life determines decay rate
- Penalty above suppress-limit do not advertise route to BGP peers
- Penalty decayed below reuse-limit re-advertise route to BGP peers penalty reset to zero when it is half of reuse-limit

# Operation



# Operation

- Only applied to inbound announcements from eBGP peers
- Alternate paths still usable
- Controlled by:

Half-life (default 15 minutes)

reuse-limit (default 750)

suppress-limit (default 2000)

maximum suppress time (default 60 minutes)

# Configuration

#### **Fixed damping**

router bgp 100

bgp dampening [<half-life> <reuse-value> <suppresspenalty> <maximum suppress time>]

#### Selective and variable damping

bgp dampening [route-map <name>]

#### Variable damping

recommendations for ISPs

http://www.ripe.net/docs/ripe-210.html

# Operation

- Care required when setting parameters
- Penalty must be less than reuse-limit at the maximum suppress time
- Maximum suppress time and half life must allow penalty to be larger than suppress limit

# Configuration

#### Examples - ×

bgp dampening 30 750 3000 60

reuse-limit of 750 means maximum possible penalty is 3000 – no prefixes suppressed as penalty cannot exceed suppress-limit

Examples - ✓

bgp dampening 30 2000 3000 60

reuse-limit of 2000 means maximum possible penalty is 8000 – suppress limit is easily reached



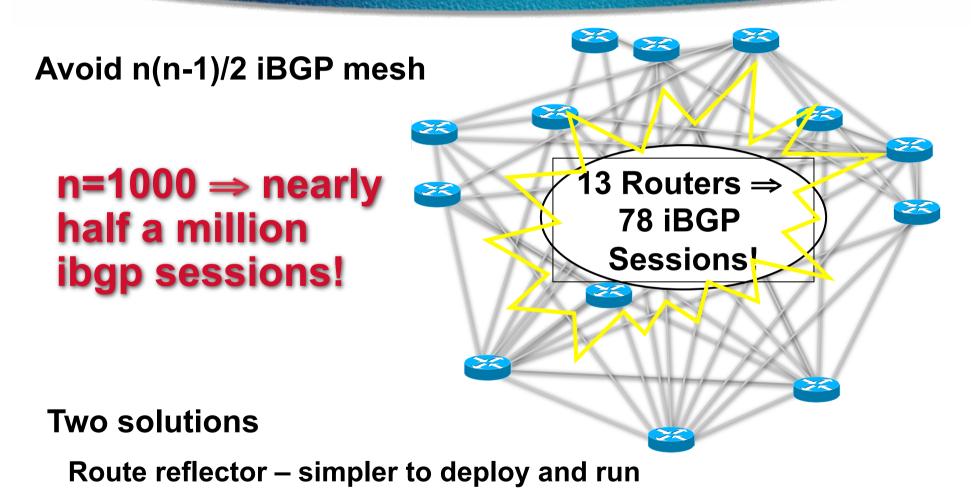
## Maximum value of penalty is

max-penalty = reuse-limit x 2  $\left(\frac{\frac{max-suppress-time}{half-life}}{half-life}\right)$ 

 Always make sure that suppress-limit is LESS than max-penalty otherwise there will be no flap damping

# Route Reflectors and Confederations

# Scaling iBGP mesh

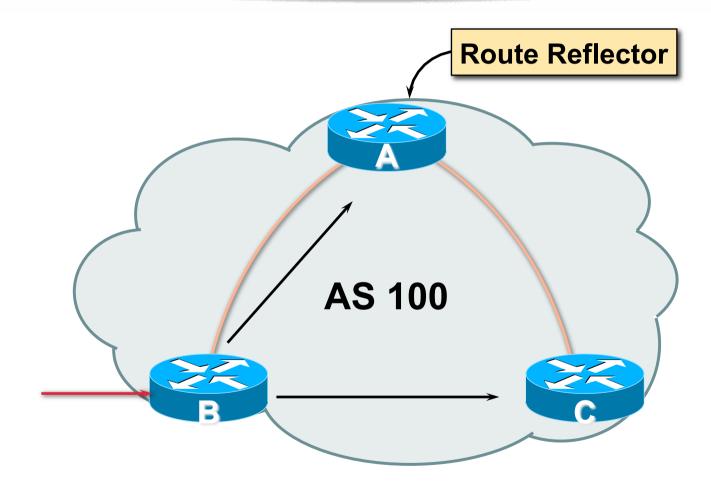


**Confederation – more complex, corner case benefits** 

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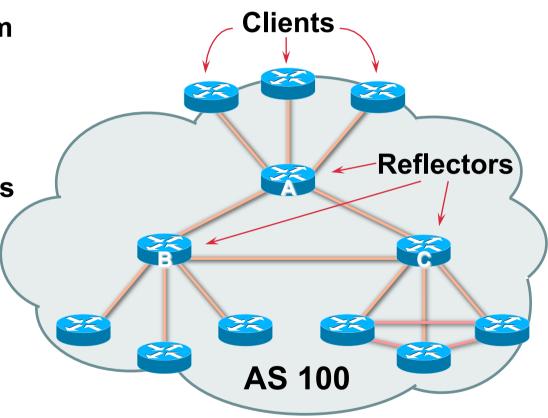
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# **Route Reflector: Principle**



# **Route Reflector**

- Reflector receives path from clients and non-clients
- Selects best path
- If best path is from client, reflect to other clients and non-clients
- If best path is from non-client, reflect to clients only
- Non-meshed clients
- Described in RFC2796



# **Route Reflector Topology**

- Divide the backbone into multiple clusters
- At least one route reflector and few clients per cluster
- Route reflectors are fully meshed
- Clients in a cluster could be fully meshed
- Single IGP to carry next hop and local routes

# Route Reflectors: Loop Avoidance

#### Originator\_ID attribute

Carries the RID of the originator of the route in the local AS (created by the RR)

#### Cluster\_list attribute

The local cluster-id is added when the update is sent by the RR

Cluster-id is automatically set from routerid (address of loopback)

Do NOT use bgp cluster-id x.x.x.x

# Route Reflectors: Redundancy

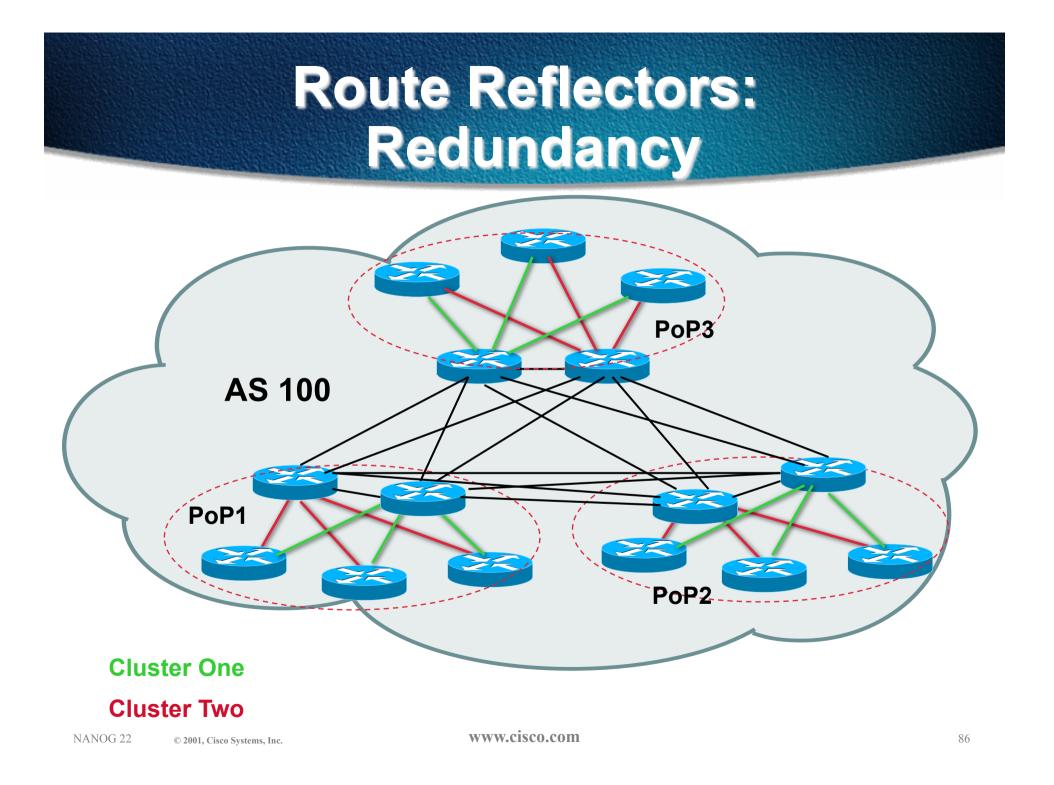
 Multiple RRs can be configured in the same cluster – not advised!

All RRs in the cluster must have the same cluster-id (otherwise it is a different cluster)

 A router may be a client of RRs in different clusters

Common today in ISP networks to overlay two clusters – redundancy achieved that way

→ Each client has two RRs = redundancy



# **Route Reflectors: Migration**

Where to place the route reflectors?
 Always follow the physical topology!

This will guarantee that the packet forwarding won't be affected

• Typical ISP network:

PoP has two core routers

**Core routers are RR for the PoP** 

Two overlaid clusters

# **Route Reflectors: Migration**

Typical ISP network:

Core routers have fully meshed iBGP

Create further hierarchy if core mesh too big

Split backbone into regions

 Configure one cluster pair at a time Eliminate redundant iBGP sessions
 Place maximum one RR per cluster
 Easy migration, multiple levels

# **Route Reflector: Migration AS 300 AS 100 AS 200**

#### Migrate small parts of the network, one part at a time.

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# Configuring a Route Reflector

router bgp 100 neighbor 1.1.1.1 remote-as 100 neighbor 1.1.1.1 route-reflector-client neighbor 2.2.2.2 remote-as 100 neighbor 2.2.2.2 route-reflector-client neighbor 3.3.3.3 remote-as 100

# Confederations

Divide the AS into sub-AS

eBGP between sub-AS, but some iBGP information is kept

Preserve NEXT\_HOP across the sub-AS (IGP carries this information)

Preserve LOCAL\_PREF and MED

- Usually a single IGP
- Described in RFC3065

# **Confederations (Cont.)**

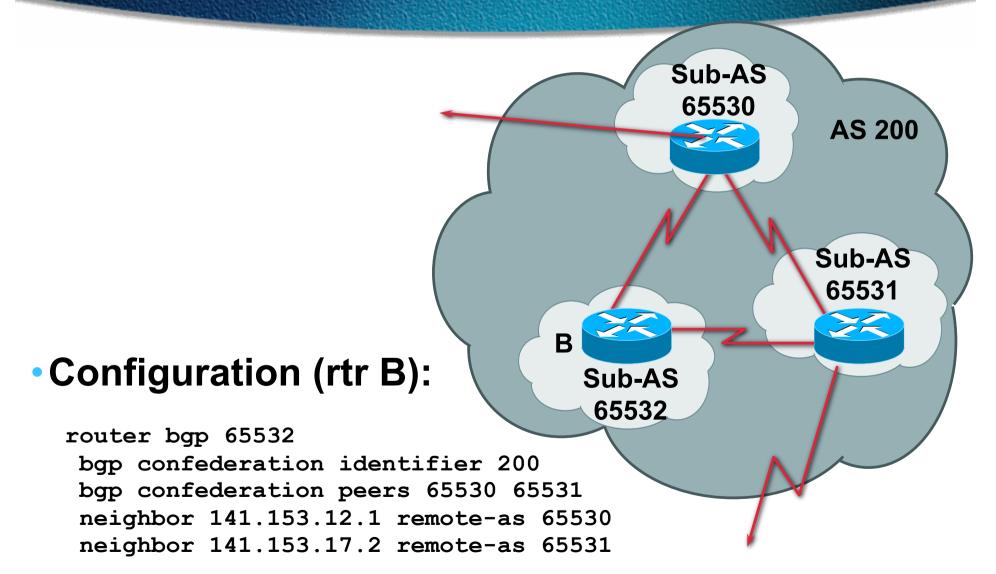
 Visible to outside world as single AS – "Confederation Identifier"

Each sub-AS uses a number from the private space (64512-65534)

 iBGP speakers in sub-AS are fully meshed

The total number of neighbors is reduced by limiting the full mesh requirement to only the peers in the sub-AS





# **Route Propagation Decisions**

- Same as with "normal" BGP:
  - From peer in same sub-AS → only to external peers
  - From external peers  $\rightarrow$  to all neighbors
- "External peers" refers to
  - Peers outside the confederation
  - Peers in a different sub-AS
    - Preserve LOCAL\_PREF, MED and NEXT\_HOP

# **Confederations (cont.)**

#### • Example (cont.):

BGP table version is 78, local router ID is 141.153.17.1 Status codes: s suppressed, d damped, h history, \* valid, > best, i - internal Origin codes: i - IGP, e - EGP, ? - incomplete Network Next Hop Metric LocPrf Weight Path \*> 10.0.0.0 141.153.14.3 0 100 (65531) 1 0 i \*> 141.153.0.0 141.153.30.2 100 (65530) i 0 0 \*> 144.10.0.0 141.153.12.1 100 (65530) i 0 0 \*> 199.10.10.0 141.153.29.2 0 100 0 (65530) 1 i

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# **RRs or Confederations**

	Internet Connectivity	Multi-Level Hierarchy	Policy Control	Scalability	Migration Complexity
Confederations	Anywhere in the Network	Yes	Yes	Medium	Medium to High
Route Reflectors	Anywhere in the Network	Yes	Yes	Very High	Very Low

#### Most new service provider networks now deploy Route Reflectors from Day One

# More points about confederations

- Can ease "absorbing" other ISPs into you ISP – e.g., if one ISP buys another (can use local-as feature to do a similar thing)
- You can use route-reflectors with confederation sub-AS to reduce the sub-AS iBGP mesh

# **BGP Scaling Techniques**

- These 4 techniques should be core requirements in all ISP networks
  - Soft reconfiguration/Route Refresh
  - Peer groups
  - **Route flap damping**
  - **Route reflectors**

# BGP for Internet Service Providers

- BGP Basics (quick recap)
- Scaling BGP
- Deploying BGP in an ISP network
- Trouble & Troubleshooting
- Multihoming Examples
- Using Communities

# Deploying BGP in an ISP Network

## **Current Practices**



# **BGP versus OSPF/ISIS**

 Internal Routing Protocols (IGPs) examples are ISIS and OSPF

used for carrying infrastructure addresses

**NOT** used for carrying Internet prefixes or customer prefixes

design goal is to minimise number of prefixes in IGP to aid scalability and rapid convergence

# **BGP versus OSPF/ISIS**

- BGP used internally (iBGP) and externally (eBGP)
- iBGP used to carry

some/all Internet prefixes across backbone

customer prefixes

eBGP used to

exchange prefixes with other ASes implement routing policy

# **BGP versus OSPF/ISIS**

## • DO NOT:

# distribute BGP prefixes into an IGP distribute IGP routes into BGP use an IGP to carry customer prefixes YOUR NETWORK WILL NOT SCALE

# Aggregation

# **Quality or Quantity?**

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# Aggregation

- ISPs receive address block from Regional Registry or upstream provider
- Aggregation means announcing the address block only, not subprefixes

Subprefixes should only be announced in special cases – see later.

 Aggregate should be generated internally Not on the network borders!

# Configuring Aggregation – Method One

- ISP has 221.10.0.0/19 address block
- To put into BGP as an aggregate:

router bgp 100

network 221.10.0.0 mask 255.255.224.0

ip route 221.10.0.0 255.255.224.0 null0

#### • The static route is a "pull up" route

more specific prefixes within this address block ensure connectivity to ISP's customers

"longest match lookup"

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# Configuring Aggregation – Method Two

#### Configuration Example

```
router bgp 109
network 221.10.0.0 mask 255.255.252.0
aggregate-address 221.10.0.0 255.255.224.0 [summary-
only]
```

Requires more specific prefix in routing table before aggregate is announced

#### {summary-only} keyword

ensures that only the summary is announced if a more specific prefix exists in the routing table

Sets "aggregator" attribute

#### Useful for debugging

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# Announcing Aggregate – Cisco IOS

#### Configuration Example

```
router bgp 100
network 221.10.0.0 mask 255.255.224.0
neighbor 222.222.10.1 remote-as 101
neighbor 222.222.10.1 prefix-list out-filter out
!
ip route 221.10.0.0 255.255.224.0 null0
!
ip prefix-list out-filter permit 221.10.0.0/19
```

### **Announcing an Aggregate**

- ISPs who don't and won't aggregate are held in poor regard by community
- Registries' minimum allocation size is now a /20

no real reason to see subprefixes of allocated blocks in the Internet

BUT there are currently >60000 /24s!

# **Receiving Prefixes**

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- ISPs should only accept prefixes which have been assigned or allocated to their downstream peer
- For example

downstream has 220.50.0.0/20 block should only announce this to peers peers should only accept this from them

## Receiving Prefixes – Cisco IOS

### Configuration Example on upstream

router bgp 100
neighbor 222.222.10.1 remote-as 101
neighbor 222.222.10.1 prefix-list customer in
!
ip prefix-list customer permit 220.50.0.0/20

- Not desirable unless really necessary special circumstances – see later
- Ask upstream to either: originate a default-route announce one prefix you can use as default

#### Downstream Router Configuration

router bgp 100
network 221.10.0.0 mask 255.255.224.0
neighbor 221.5.7.1 remote-as 101
neighbor 221.5.7.1 prefix-list infilter in
neighbor 221.5.7.1 prefix-list outfilter out

!
ip prefix-list infilter permit 0.0.0.0/0
!
ip prefix-list outfilter permit 221.10.0.0/19

### Upstream Router Configuration

router bgp 101 neighbor 221.5.7.2 remote-as 100 neighbor 221.5.7.2 default-originate neighbor 221.5.7.2 prefix-list cust-in in neighbor 221.5.7.2 prefix-list cust-out out ! ip prefix-list cust-in permit 221.10.0.0/19 ! ip prefix-list cust-out permit 0.0.0.0/0

 If necessary to receive prefixes from upstream provider, care is required

### don't accept RFC1918 etc prefixes

http://www.ietf.org/internet-drafts/draft-manning-dsua-06.txt

### don' t accept your own prefix don' t accept default (unless you need it) don' t accept prefixes longer than /24

### **Receiving Prefixes**

```
router bop 100
network 221.10.0.0 mask 255.255.224.0
neighbor 221.5.7.1 remote-as 101
neighbor 221.5.7.1 prefix-list in-filter in
I
ip prefix-list in-filter deny 0.0.0.0/0
                                                 ! Block default
ip prefix-list in-filter deny 0.0.0.0/8 le 32
ip prefix-list in-filter deny 10.0.0.0/8 le 32
ip prefix-list in-filter deny 127.0.0.0/8 le 32
ip prefix-list in-filter deny 169.254.0.0/16 le 32
ip prefix-list in-filter deny 172.16.0.0/12 le 32
ip prefix-list in-filter deny 192.0.2.0/24 le 32
ip prefix-list in-filter deny 192.168.0.0/16 le 32
ip prefix-list in-filter deny 221.10.0.0/19 le 32 ! Block local prefix
ip prefix-list in-filter deny 224.0.0.0/3 le 32 ! Block multicast
ip prefix-list in-filter deny 0.0.0.0/0 ge 25
                                                 ! Block prefixes >/24
ip prefix-list in-filter permit 0.0.0.0/0 le 32
```

# **Prefixes into iBGP**

### Injecting prefixes into iBGP

- Use iBGP to carry customer prefixes don't ever use IGP
- Point static route to customer interface
- Use BGP network statement
- As long as static route exists (interface active), prefix will be in BGP

### Router Configuration network statement

#### • Example:

```
interface loopback 0
 ip address 215.17.3.1 255.255.255.255
interface Serial 5/0
 ip unnumbered loopback 0
 ip verify unicast reverse-path
ip route 215.34.10.0 255.255.252.0 Serial 5/0
router bgp 100
network 215.34.10.0 mask 255.255.252.0
```

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### Injecting prefixes into iBGP

 interface flap will result in prefix withdraw and re-announce

use "ip route...permanent"

Static route always exists, even if interface is down  $\rightarrow$  prefix announced in iBGP

 many ISPs use redistribute static rather than network statement

only use this if you understand why

# Inserting prefixes into BGP – redistribute static

• Care required with redistribute!

redistribute <routing-protocol> means everything in the <routing-protocol> will be transferred into the current routing protocol

Does not scale if uncontrolled

Best avoided if at all possible

redistribute normally used with "route-maps" and under tight administrative control

# Router Configuration redistribute static

#### • Example:

```
ip route 215.34.10.0 255.255.252.0 Serial 5/0
router bgp 100
 redistribute static route-map static-to-bgp
<snip>
route-map static-to-bgp permit 10
match ip address prefix-list ISP-block
 set origin igp
<snip>
ip prefix-list ISP-block permit 215.34.10.0/22 le 30
```

### Injecting prefixes into iBGP

Route-map ISP-block can be used for many things:

setting communities and other attributes setting origin code to IGP, etc

 Be careful with prefix-lists and route-maps absence of either/both could mean all statically routed prefixes go into iBGP

# **Configuration Tips**

### **iBGP and IGPs**

Make sure loopback is configured on router

iBGP between loopbacks, NOT real interfaces

- Make sure IGP carries loopback /32 address
- Make sure IGP carries DMZ nets
   Or use next-hop-self on iBGP neighbours
   neighbor x.x.x next-hop-self

### Next-hop-self

Used by many ISPs on edge routers

Preferable to carrying DMZ /30 addresses in the IGP

- Reduces size of IGP to just core infrastructure
- Alternative to using ip unnumbered

Helps scale network

BGP speaker announces external network using local address (loopback) as next-hop

### **BGP Template – iBGP peers**

Z

**iBGP** Peer Group

**AS100** router bgp 100 neighbor internal peer-group neighbor internal description ibgp peers neighbor internal remote-as 100 neighbor internal update-source Loopback0 neighbor internal next-hop-self neighbor internal send-community neighbor internal version 4 neighbor internal password 7 03085A09 neighbor 1.0.0.1 peer-group internal neighbor 1.0.0.2 peer-group internal

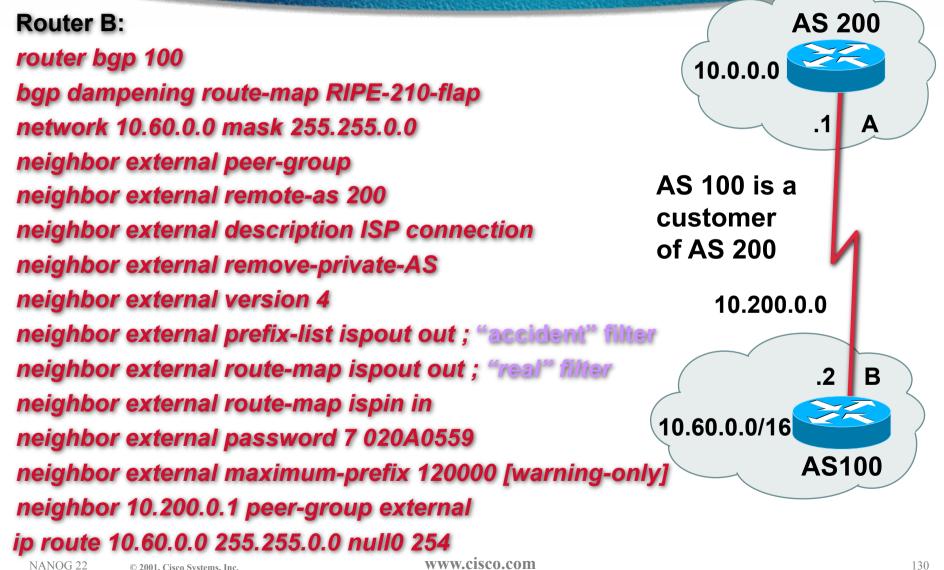
### **BGP Template – iBGP peers**

- Use peer-groups
- iBGP between loopbacks!
- Next-hop-self
  - Keep DMZ and point-to-point out of IGP
- Always send communities in iBGP
   Otherwise accidents will happen
- Hardwire BGP to version 4

Yes, this is being paranoid!

Use passwords on iBGP session

### **BGP Template – eBGP peers**



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### **BGP Template – eBGP peers**

- BGP damping use RIPE-210 parameters
- Remove private ASes from announcements Common omission today
- Use extensive filters, with "backup"
- Use password agreed between you and peer on eBGP session
- Use maximum-prefix tracking

Router will warn you if there are sudden changes in BGP table size, bringing down eBGP if necessary

### More BGP "defaults"

- Log neighbour changes
   bgp log-neighbor-changes
- Enable deterministic MED

bgp deterministic-med

Otherwise bestpath could be different every time BGP session is reset

 Make BGP admin distance higher than any IGP

### distance bgp 200 200 200

### **Customer Aggregation**

BGP customers

Offer max 3 types of feeds (easier than custom configuration per peer)

**Use communities** 

Static customers

**Use communities** 

 Differentiate between different types of prefixes

### Makes eBGP filtering easy

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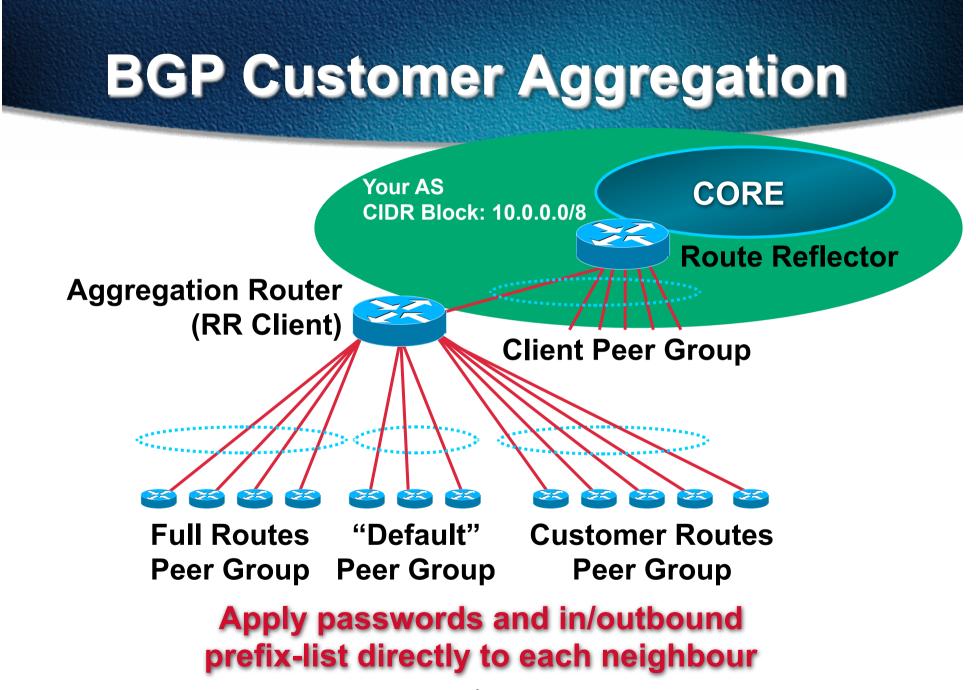
### BGP Customer Aggregation Guidelines

- Define at least three peer groups: cust-default—send default route only cust-cust—send customer routes only cust-full —send full Internet routes
- Identify routes via communities e.g.
   100:4100=customers; 100:4500=peers
- Apply passwords per neighbour
- Apply inbound & outbound prefix-list per neighbour

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### Static Customer Aggregation Guidelines

- Identify routes via communities, e.g. 100:4000=my address blocks 100:4200=customers from my block

  - **100:4300=customers outside my block**
  - Helps with aggregation, iBGP, filtering
- BGP network statements on aggregation routers set correct community

### Sample core configuration

eBGP peers and upstreams

Send communities 100:4000, 100:4100 and 100:4300, receive everything

iBGP full routes

Send everything (only network core)

iBGP partial routes

Send communities 100:4000, 100:4100, 100:4200, 100:4300 and 100:4500 (edge routers, peering routers, IXP routers)

Simple configuration with peer-groups and route-maps

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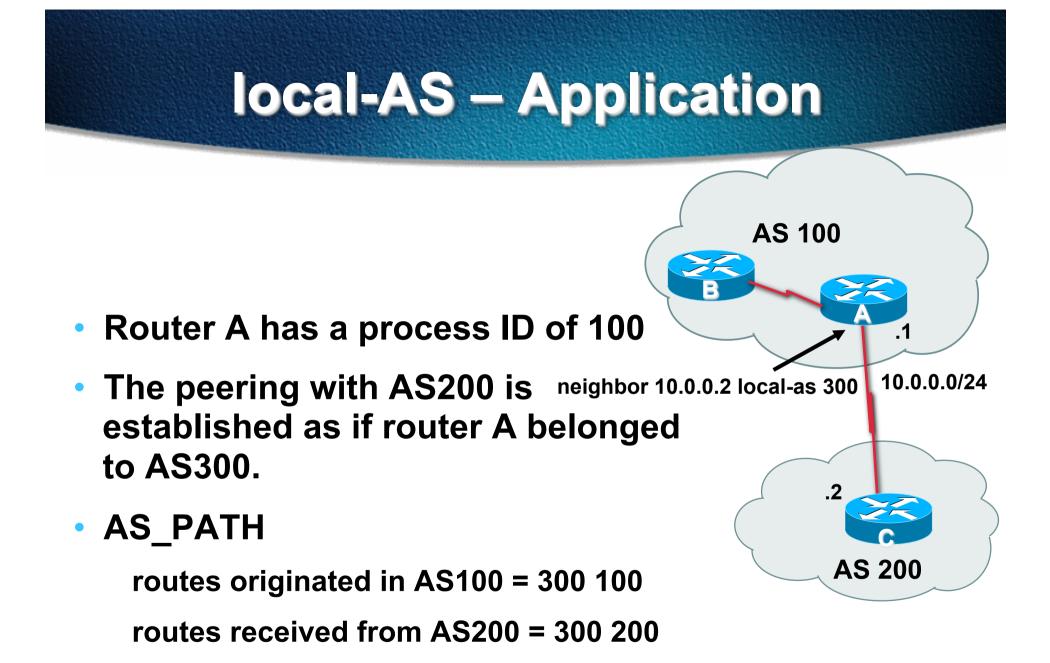
### **Acquisitions!**

- Your ISP has just bought another ISP How to merge networks?
- Options:

use confederations – make their AS a sub-AS (only useful if you are using confederations already)

use the BGP local-as feature to implement a gradual transition – overrides BGP process ID

neighbor x.x.x.x local-as as-number



### BGP for Internet Service Providers

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# Troubleshooting

### **Staying out of Trouble**



## Potential Caveats and Operational Problems

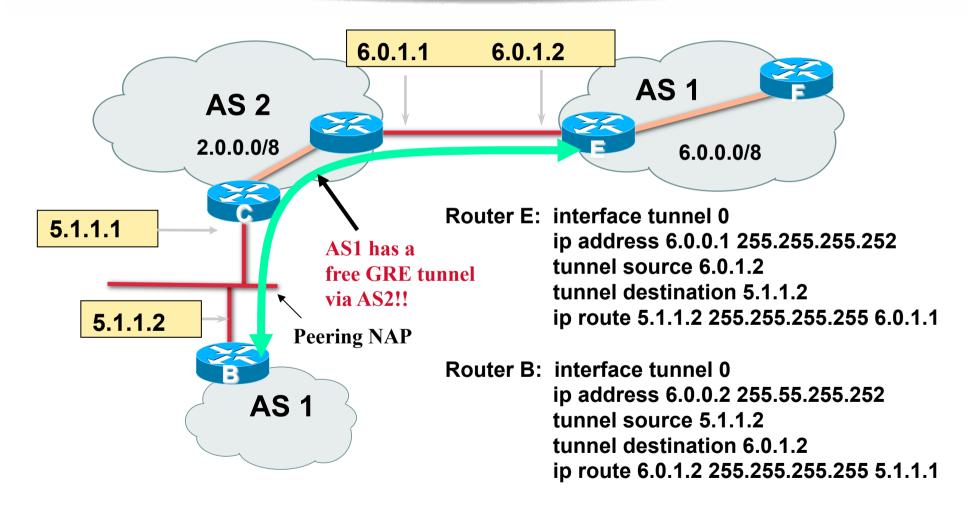
- GRE Tunnels & IXPs
- Auto-summarisation & synchronisation
- Route Reflectors

Follow the topology

Common Problems

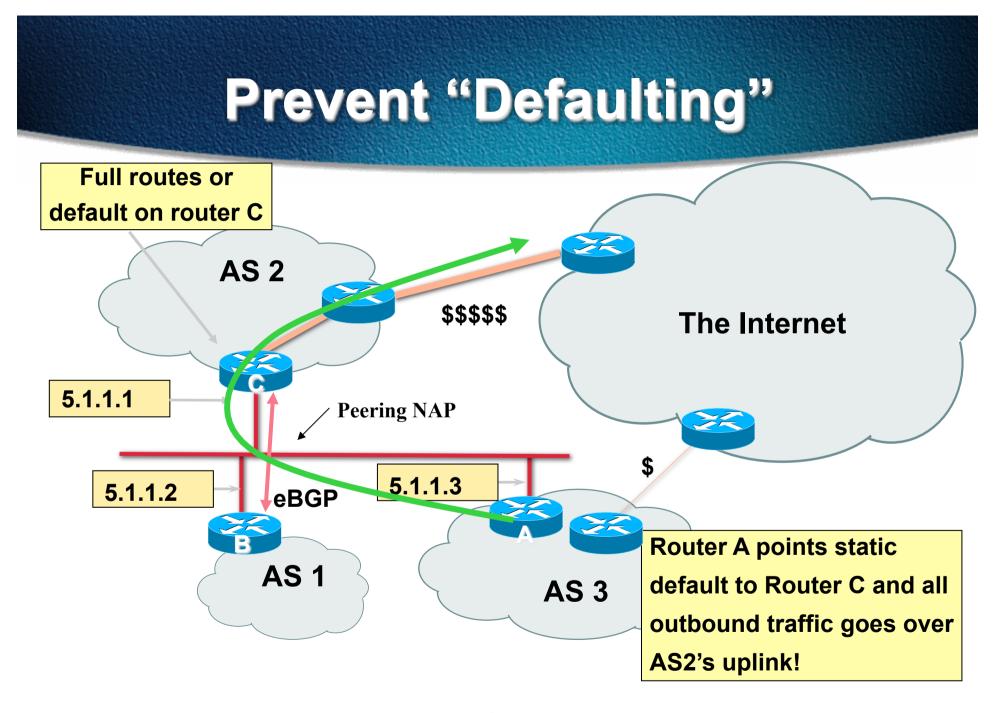
...and the solutions!

### **Prevent GRE VPNs**



Don't carry IXP net in your IGP – use next-hop-self!

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### Watch out at IXPs/NAPs

- IXP router should not carry full routes or have a default
- ISP should not carry IXP/NAP network prefix internally

**Use BGP next-hop-self** 

- or -
- Use RPF check for non-peers
- Use good filters for peers

### Auto Summarisation – Cisco IOS

- Historical feature
- Automatically summarises subprefixes to the classful network for prefixes redistributed into BGP

**Example:** 

61.10.8.0/22 --> 61.0.0/8

• Must be turned off for any Internet connected site using BGP.

router bgp 109

```
no auto-summary
```

### Synchronisation – Cisco IOS

- Historical feature
- BGP will not advertise a route before all routers in the AS have learned it via an IGP
- Disable synchronisation if:

AS doesn't pass traffic from one AS to another, or

All transit routers in AS run BGP, or

iBGP is used across backbone

router bgp 109

no synchronization

# Troubleshooting

# Common Problems and their Solutions

## Troubleshooting – Examples

- Missing routes
- Route Oscillation
- Routing Loops
- Troubleshooting hints

### **Route Origination**

#### Network statement with mask

R1# show run | begin bgp

network 200.200.0.0 mask 255.255.252.0

#### BGP is not originating the route???

R1# show ip bgp | include 200.200.0.0 R1#

• Do we have the exact route?

R1# show ip route 200.200.0.0 255.255.252.0

% Network not in table

### **Route Origination**

Nail down routes you want to originate
 R1#ip route 200.200.0.0 255.255.252.0 Null 0

#### Check the RIB

R1# show ip route 200.200.0.0 255.255.252.0

200.200.0.0/22 is subnetted, 1 subnets

S 200.200.0.0 [1/0] via Null 0

• BGP originates the route!!

R1# show ip bgp | include 200.200.0.0\*> 200.200.0.0/220.0.0.0032768

200

### **Route Oscillation**

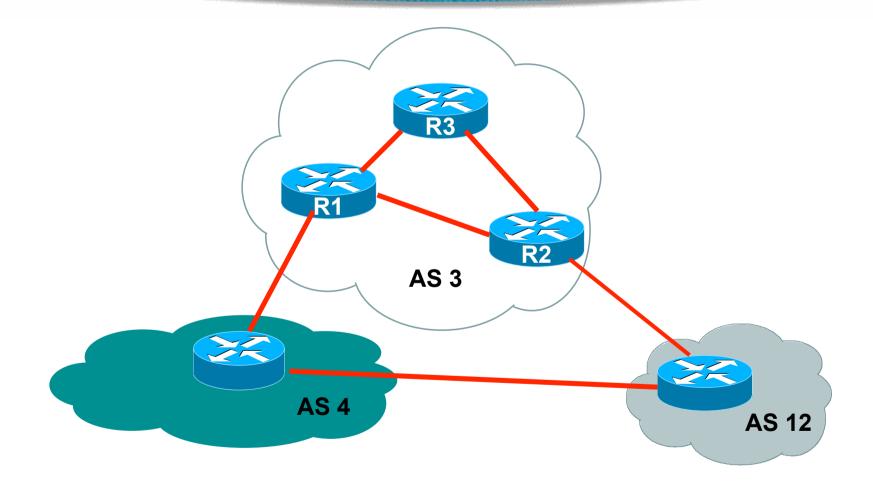
### One of the most common problems!

Every minute routes flap in the routing table from one next hop to another

With large routing table the most obvious symptom is high CPU in the "BGP-Router" process

Can be frustrating to track down unless you have seen it before!

### **Route Oscillation – Diagram**



### **Route Oscillation – Symptom**

R3#show ip bgp summary
BGP router identifier 3.3.3.3, local AS number 3
BGP table version is 502, main routing table version 502
<b>267 network entries and 272 paths</b> using 34623 bytes of memory
R3#sh ip route summary   begin bgp
bgp 3 4 6 520 1400
External: 0 Internal: 10 Local: 0
internal 5 5800
Total         10         263         13936         43320

• Watch for:

table version number incrementing rapidly

number of networks/paths or external/internal routes changing.

# Pick up a bgp route from the RIB that is less than a minute old and watch what happens with the routing/bgp table ...

```
R3#show ip route 156.1.0.0
Routing entry for 156.1.0.0/16
  Known via "bgp 3", distance 200, metric 0
Routing Descriptor Blocks:
  * 1.1.1.1, from 1.1.1.1, 00:00:53 ago
      Route metric is 0, traffic share count is 1
      AS Hops 2, BGP network version 474
R3#show ip bgp 156.1.0.0
BGP routing table entry for 156.1.0.0/16, version 474
Paths: (2 available, best #1)
  Advertised to non peer-group peers:
    2.2.2.2
  4 12
    1.1.1.1 from 1.1.1.1 (1.1.1.1)
      Origin IGP, localpref 100, valid, internal, best
  12
    142.108.10.2 (inaccessible) from 2.2.2.2 (2.2.2.2)
      Origin IGP, metric 0, localpref 100, valid, internal
```

...and after bgp\_scanner runs (by default once a minute):

```
R3#sh ip route 156.1.0.0
Routing entry for 156.1.0.0/16
  Known via "bgp 3", distance 200, metric 0
    Routing Descriptor Blocks:
  * 142.108.10.2, from 2.2.2.2, 00:00:27 ago
      Route metric is 0, traffic share count is 1
      AS Hops 1, BGP network version 478
R3#sh ip bgp 156.1.0.0
BGP routing table entry for 156.1.0.0/16, version 478
Paths: (2 available, best #2)
  Advertised to non peer-group peers:
    1.1.1.1
  4 12
    1.1.1.1 from 1.1.1.1 (1.1.1.1)
      Origin IGP, localpref 100, valid, internal
  12
    142.108.10.2 from 2.2.2.2 (2.2.2.2)
      Origin IGP, metric 0, localpref 100, valid, internal, best
```

#### Let's take a look at the next hop at this point!

```
R3#show ip route 142.108.10.2
Routing entry for 142.108.0.0/16
  Known via "bgp 3", distance 200, metric 0
 Routing Descriptor Blocks:
  * 142.108.10.2, from 2.2.2.2, 00:00:50 ago
      Route metric is 0, traffic share count is 1
      AS Hops 1, BGP network version 476
R3#show ip bqp 142.108.10.2
BGP routing table entry for 142.108.0.0/16, version 476
Paths: (2 available, best #2)
  Advertised to non peer-group peers:
    1.1.1.1
  4 12
    1.1.1.1 from 1.1.1.1 (1.1.1.1)
      Origin IGP, localpref 100, valid, internal
  12
    142.108.10.2 from 2.2.2.2 (2.2.2.2)
      Origin IGP, metric 0, localpref 100, valid, internal, best
```

Next-hop is recursive !!! This will be detected next time the scanner runs and the other path will be installed in the RIB instead

```
R3#sh debug
BGP events debugging is on
BGP updates debugging is on
IP routing debugging is on
R3#
BGP: scanning routing tables
BGP: nettable_walker 142.108.0.0/16 calling revise_route
RT: del 142.108.0.0 via 142.108.10.2, bgp metric [200/0]
BGP: revise route installing 142.108.0.0/16 -> 1.1.1.1
RT: add 142.108.0.0/16 via 1.1.1.1, bgp metric [200/0]
RT: del 156.1.0.0 via 142.108.10.2, bgp metric [200/0]
BGP: revise route installing 156.1.0.0/16 -> 1.1.1.1
RT: add 156.1.0.0/16 via 1.1.1.1, bgp metric [200/0]
```

The route to the next-hop is now valid and at the next bgp scan we will change to the shorter as-path path, and so on ...

#### R3#

#### BGP: scanning routing tables

BGP: ip nettable\_walker 142.108.0.0/16 calling revise\_route
RT: del 142.108.0.0 via 1.1.1.1, bgp metric [200/0]
BGP: revise route installing 142.108.0.0/16 -> 142.108.10.2
RT: add 142.108.0.0/16 via 142.108.10.2, bgp metric [200/0]
BGP: nettable\_walker 156.1.0.0/16 calling revise\_route
RT: del 156.1.0.0 via 1.1.1.1, bgp metric [200/0]
BGP: revise route installing 156.1.0.0/16 -> 142.108.10.2
RT: add 156.1.0.0/16 via 142.108.10.2, bgp metric [200/0]

### **Route Oscillation – Summary**

- iBGP preserves the next-hop information from eBGP
- To avoid problems
  - use "next-hop-self" for iBGP peering

-or-

make sure you advertise the next-hop prefix via the IGP

### **Inconsistent Route Selection**

 Two common problems with route selection Inconsistency

**Appearance of an Incorrect decision** 

- RFC 1771 defines the decision algorithm
- Every vendor has tweaked the algorithm http://www.cisco.com/warp/public/459/25.shtml
- Route Selection problems can result from oversights in RFC1771

### **Inconsistent Route Selection**

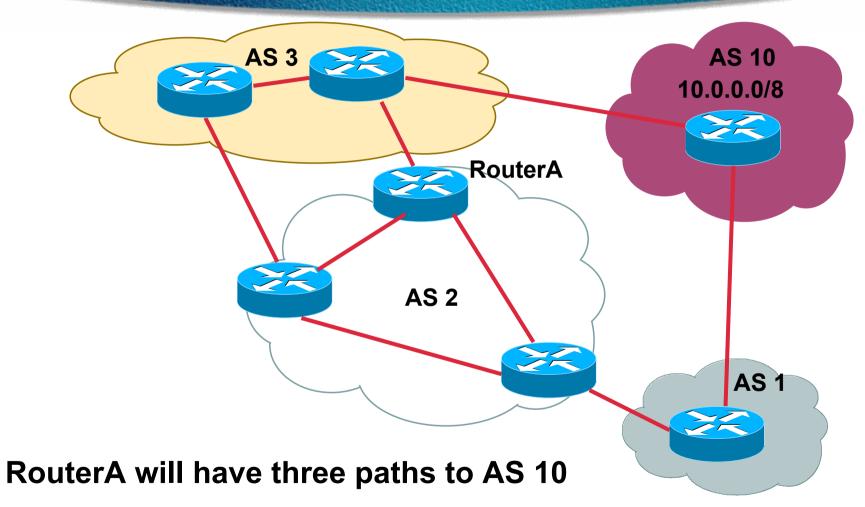
- RFC says that MED is not always compared
- As a result, the ordering of the paths can affect the decision process
- By default, the prefixes are compared in order of arrival (most recent to oldest)

use **bgp deterministic-med** to order paths consistently

the bestpath is recalculated as soon as the command is entered

enable in all the routers in the AS

### Symptom – Diagram



MEDs from AS 3 will not be compared with MEDs from AS 1

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### **Inconsistent Route Selection**

```
RouterA#sh ip bgp 10.0.0.0
BGP routing table entry for 10.0.0.0/8, version 40
Paths: (3 available, best #3, advertised over IBGP, EBGP)
3 10
2.2.2.2 from 2.2.2.2
Origin IGP, metric 20, localpref 100, valid, internal
3 10
3.3.3.3 from 3.3.3.3
Origin IGP, metric 30, valid, external
1 10
1.1.1.1 from 1.1.1.1
Origin IGP, metric 0, localpref 100, valid, internal, best
```

### Initial State

Path 1 beats Path 2 – Lower MED

#### Path 3 beats Path 1 – Lower Router-ID

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### **Inconsistent Route Selection**

```
RouterA#sh ip bgp 10.0.0.0
BGP routing table entry for 10.0.0.0/8, version 40
Paths: (3 available, best #3, advertised over IBGP, EBGP)
1 10
1.1.1.1 from 1.1.1.1
Origin IGP, metric 0, localpref 100, valid, internal
3 10
2.2.2.2 from 2.2.2.2
Origin IGP, metric 20, localpref 100, valid, internal
3 10
3.3.3.3 from 3.3.3.3
Origin IGP, metric 30, valid, external, best
```

 1.1.1.1 bounced so the paths are re-ordered Path 1 beats Path 2 – Lower Router-ID Path 3 beats Path 1 – External vs Internal

## Deterministic MED – Operation

- The paths are ordered by Neighbour AS
- The bestpath for each Neighbour AS group is selected
- The overall bestpath results from comparing the winners from each group
- The bestpath will be consistent because paths will be placed in a deterministic order

### **Deterministic MED – Result**

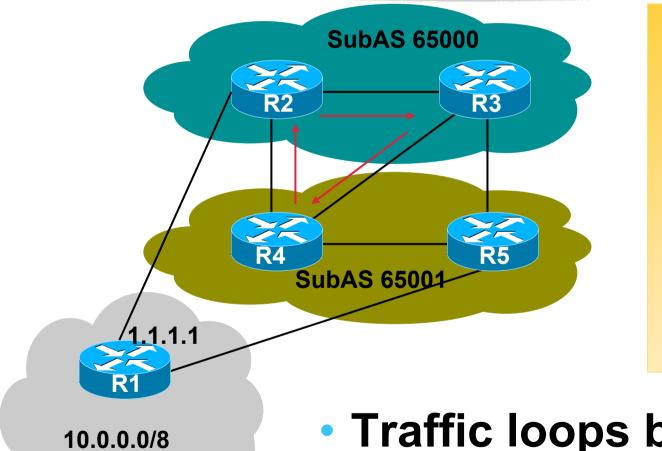
```
RouterA#sh ip bgp 10.0.0.0
BGP routing table entry for 10.0.0.0/8, version 40
Paths: (3 available, best #1, advertised over IBGP, EBGP)
1 10
1.1.1.1 from 1.1.1.1
Origin IGP, metric 0, localpref 100, valid, internal, best
3 10
2.2.2.2 from 2.2.2.2
Origin IGP, metric 20, localpref 100, valid, internal
3 10
3.3.3.3 from 3.3.3.3
Origin IGP, metric 30, valid, external
```

```
Path 1 is best for AS 1
Path 2 beats Path 3 for AS 3 – Lower MED
Path 1 beats Path 2 – Lower Router-ID
```

## Deterministic MED – Summary

- If multihoming with multiple ISPs and peering with one ISP at multiple points:
  - use "bgp deterministic-med" enable it on all routers in the AS
- Always use "bgp deterministic-med"

### **Routing Loop – Problem**



traceroute 10.1.1.1

1 30.100.1.1 2 20.20.20.4 - R3 3 30.1.1.26 - R4 4 30.1.1.17 - R2 5 20.20.20.4 - R3 6 30.1.1.26 - R4 7 30.1.1.17 - R2 8 20.20.20.4 9 30.1.1.26 10 30.1.1.17

#### 10.0.0.0/8 SubAS 65002

### Traffic loops between R3, R4, and R2

- First grab a "show ip route" from the three problem routers
- R3 is forwarding traffic to 1.1.1.1 (R1)

```
R3# show ip route 10.1.1.1
Routing entry for 10.0.0.0/8
Known via "bgp 65000", distance 200, metric 0
Routing Descriptor Blocks:
1.1.1.1, from 5.5.5.5, 01:46:43 ago
Route metric is 0, traffic share count is 1
AS Hops 0, BGP network version 0
* 1.1.1.1, from 4.4.4.4, 01:46:43 ago
Route metric is 0, traffic share count is 1
AS Hops 0, BGP network version 0
```

#### R4 is also forwarding to 1.1.1.1 (R1)

```
R4# show ip route 10.1.1.1
Routing entry for 10.0.0.0/8
Known via "bgp 65001", distance 200, metric 0
Routing Descriptor Blocks:
* 1.1.1.1, from 5.5.5.5, 01:47:02 ago
Route metric is 0, traffic share count is 1
AS Hops 0
```

### • R2 is forwarding to 3.3.3.3? (R3)

```
R2# show ip route 10.1.1.1
```

```
Routing entry for 10.0.0/8
```

Known via "bgp 65000", distance 200, metric 0

**Routing Descriptor Blocks:** 

```
* 3.3.3.3, from 3.3.3.3, 01:47:00 ago
```

Route metric is 0, traffic share count is 1

AS Hops 0, BGP network version 3

#### Very odd that the NEXT\_HOP is in the middle of the network

#### Verify BGP paths on R2

```
R2#show ip bgp 10.0.0.0
BGP routing table entry for 10.0.0.0/8, version 3
Paths: (4 available, best #1)
Advertised to non peer-group peers:
    1.1.1.1 5.5.5.5 4.4.4.4
  (65001 65002)
    3.3.3 (metric 11) from 3.3.3.3 (3.3.3.3)
    Origin IGP, metric 0, localpref 100, valid, confed-
internal, best
  (65002)
    1.1.1.1 (metric 5010) from 1.1.1.1 (1.1.1.1)
    Origin IGP, metric 0, localpref 100, valid, confed-
external
```

- R3 path is better than R1 path because of IGP cost to NEXT\_HOP
- R3 is advertising the path to us with a NEXT\_HOP of 3.3.3.3 ???

#### What is R3 advertising?

```
R3# show ip bgp 10.0.0.0
BGP routing table entry for 10.0.0.0/8, version 3
Paths: (2 available, best #1, table Default-IP-Routing-Table)
Advertised to non peer-group peers:
   5.5.5 2.2.2.2
  (65001 65002)
   1.1.1.1 (metric 5031) from 4.4.4.4 (4.4.4.4)
        Origin IGP, metric 0, localpref 100, valid, confed-
external, best, multipath
  (65001 65002)
   1.1.1.1 (metric 5031) from 5.5.5.5 (5.5.5.5)
        Origin IGP, metric 0, localpref 100, valid, confed-
external, multipath
```

#### Hmmm, R3 is using multipath to load-balance

R3#show run | include maximum

#### maximum-paths 6

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### **Routing Loop – Solution**

 "maximum-paths" tells the router to reset the NEXT\_HOP to himself

R3 sets NEXT\_HOP to 3.3.3.3

- Forces traffic to come to him so he can loadbalance
- Is typically used for multiple eBGP sessions to an AS

Be careful when using in Confederations!!

 Need to make R2 prefer the path from R1 to prevent the routing loop

Make IGP metric to 1.1.1.1 better than IGP metric to 4.4.4.4

- High CPU in "Router BGP" is normally a sign of a convergence problem
- Find a prefix that changes every minute show ip route | include , 00:00
- Troubleshoot/debug that one prefix

BGP routing loop?

First, check for IGP routing loops to BGP NEXT\_HOPs

BGP loops are normally caused by

Not following physical topology in RR environment

Multipath within confederations

Lack of a full iBGP mesh

 Get the following from each router in the loop path

show ip route x.x.x.x

show ip bgp x.x.x.x

show ip route NEXT\_HOP

• "show ip bgp neighbor x.x.x.x advertised-routes"

Lets you see a list of NLRI that you sent a peer

Note: The attribute values shown are taken from the BGP table. Attribute modifications by outbound routemaps will not be shown.

"show ip bgp neighbor x.x.x.x routes"

Displays routes x.x.x.x sent to us that made it through our inbound filters

"show ip bgp neighbor x.x.x.x received-routes"
 Can only use if "soft-reconfig inbound" is configured
 Displays all routes received from a peer, even those that were denied

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• "clear ip bgp x.x.x.x in"

Ask x.x.x.x to resend his UPDATEs to us

• "clear ip bgp x.x.x.x out"

Tells BGP to resend UPDATEs to x.x.x.x

"debug ip bgp update"

**Always** use an ACL to limit output

Great for troubleshooting "Automatic Denies"

"debug ip bgp x.x.x.x update"
 Allows you to debug updates to/from a specific peer
 Handy if multiple peers are sending you the same prefix

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### Summary/Tips

- Isolate the problem!!
- Use ACLs when enabling debug commands
- Enable bgp log-neighbor-changes
- IP reachability must exist for sessions to be established

learned from IGP

make sure the source and destination addresses match the configuration

## BGP for Internet Service Providers

- BGP Basics (quick recap)
- Scaling BGP
- Deploying BGP in an ISP network
- Trouble & Troubleshooting
- Multihoming Examples
- Using Communities

## Multihoming



## Multihoming Definition

 More than one link external to the local network

two or more links to the same ISP two or more links to different ISPs

 Usually two external facing routers one router gives link and provider redundancy only

### **AS Numbers**

- An Autonomous System Number is required by BGP
- Obtained from upstream ISP or Regional Registry
- Necessary when you have links to more than one ISP or exchange point

## **Configuring Policy**

Three BASIC Principles
 prefix-lists to filter prefixes
 filter-lists to filter ASNs
 route-maps to apply policy
 Avoids confusion!

## **Originating Prefixes**

#### Basic Assumptions

MUST announce assigned address block to Internet

MAY also announce subprefixes – reachability is not guaranteed

**RIR minimum allocation is /20** 

several ISPs filter RIR blocks on this boundary called "Net Police" by some

## Part of the "Net Police" prefix list

#### !! APNIC

ip prefix-list FILTER permit 61.0.0.0/8 ge 9 le 20 ip prefix-list FILTER permit 202.0.0.0/7 ge 9 le 20 ip prefix-list FILTER permit 210.0.0.0/7 ge 9 le 20 ip prefix-list FILTER permit 218.0.0.0/8 ge 9 le 20 **!!** ARIN ip prefix-list FILTER permit 63.0.0.0/8 ge 9 le 20 ip prefix-list FILTER permit 64.0.0.0/7 ge 9 le 20 ip prefix-list FILTER permit 66.0.0.0/8 ge 9 le 20 ip prefix-list FILTER permit 199.0.0.0/8 ge 9 le 20 ip prefix-list FILTER permit 200.0.0.0/8 ge 9 le 20 ip prefix-list FILTER permit 204.0.0.0/6 ge 9 le 20 ip prefix-list FILTER permit 208.0.0.0/7 ge 9 le 20 ip prefix-list FILTER permit 216.0.0.0/8 ge 9 le 20 **!!** RIPE NCC ip prefix-list FILTER permit 62.0.0.0/8 ge 9 le 20 ip prefix-list FILTER permit 80.0.0.0/7 ge 9 le 20 ip prefix-list FILTER permit 193.0.0.0/8 ge 9 le 20 ip prefix-list FILTER permit 194.0.0.0/7 ge 9 le 20 ip prefix-list FILTER permit 212.0.0.0/7 ge 9 le 20

## "Net Police" prefix list issues

- meant to "punish" ISPs who won't and don't aggregate
- impacts legitimate multihoming
- impacts regions where domestic backbone is unavailable or costs \$\$\$ compared with international bandwidth
- hard to maintain requires updating when RIRs start allocating from new address blocks
- don't do it unless consequences understood and you are prepared to keep it current

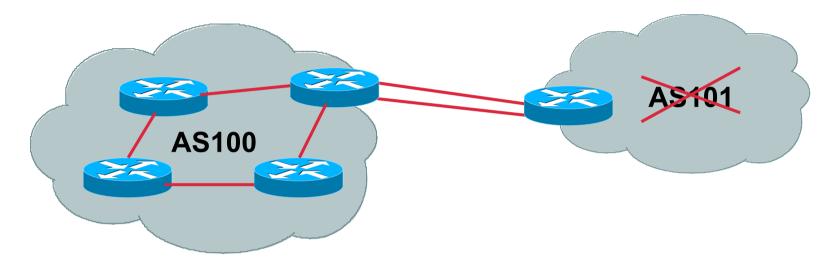
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## **Multihoming Options**

## **Multihoming Scenarios**

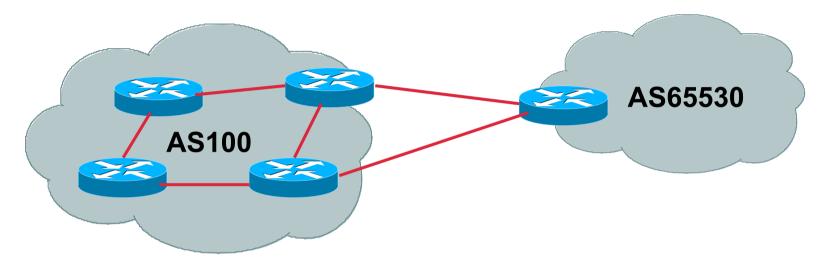
- Stub network
- Multi-homed stub network
- Multi-homed network
- Configuration Options

### **Stub Network**



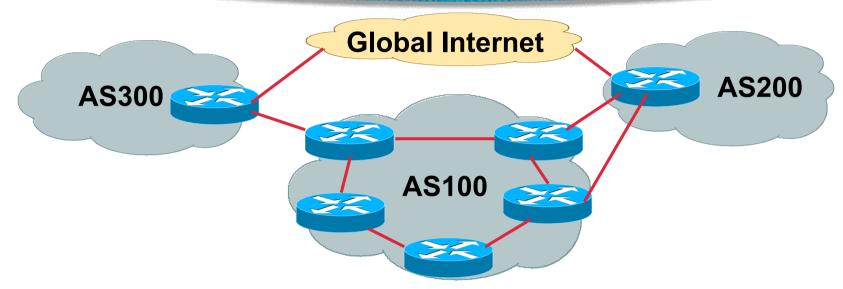
- No need for BGP
- Point static default to upstream ISP
- Upstream ISP advertises stub network
- Policy confined within upstream ISP's policy

## **Multi-homed Stub Network**



- Use BGP (not IGP or static) to loadshare
- Use private AS (ASN > 64511)
- Upstream ISP advertises stub network
- Policy confined within upstream ISP's policy

## **Multi-Homed Network**



Many situations possible

multiple sessions to same ISP

secondary for backup only

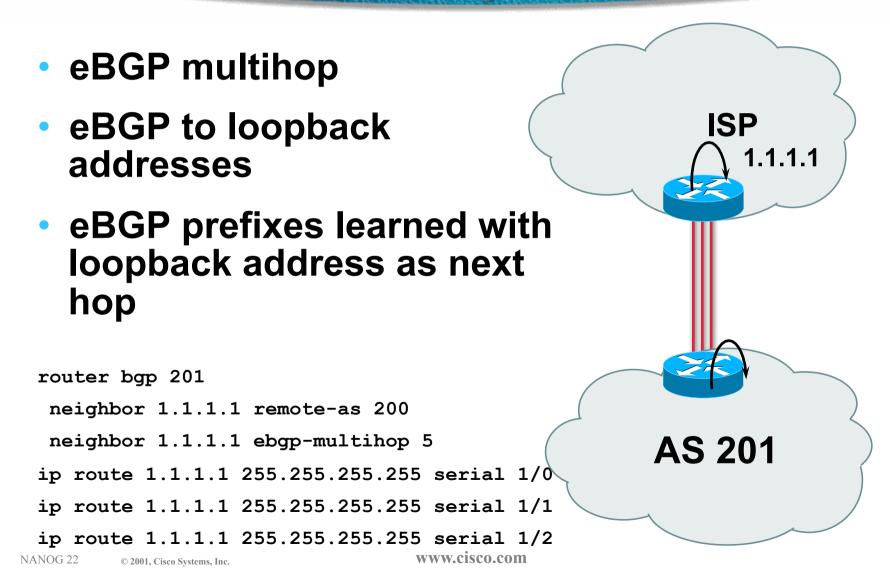
load-share between primary and secondary

selectively use different ISPs

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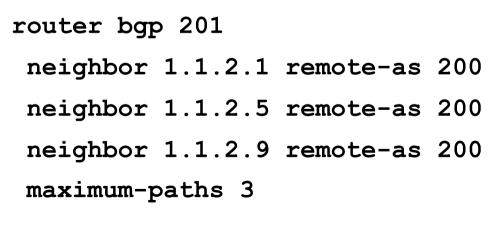
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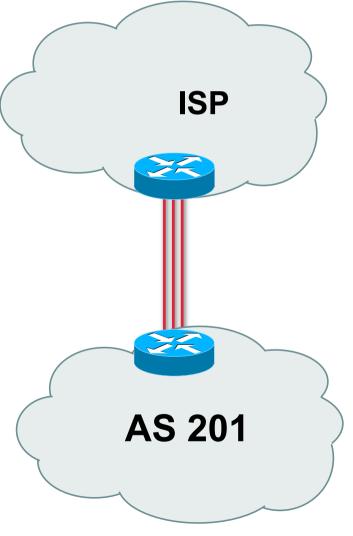
## Multiple Sessions to an ISP – Example One



## Multiple Sessions to an ISP – Example Two

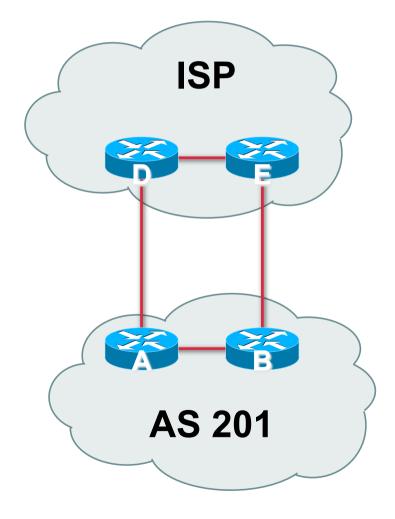
- BGP multi-path
- Three BGP sessions required
- limit of 6 parallel paths





## **Multiple Sessions to an ISP**

- Simplest scheme is to use defaults
- Learn/advertise prefixes for better control
- Planning and some work required to achieve loadsharing
- No magic solution

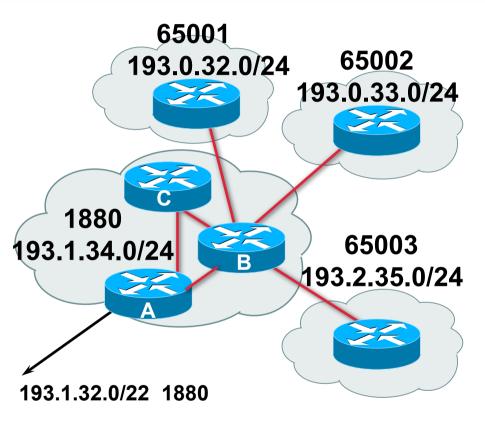


## **Private-AS – Application**

#### Applications

ISP with singlehomed customers (RFC2270)

corporate network with several regions and connections to the Internet only in the core



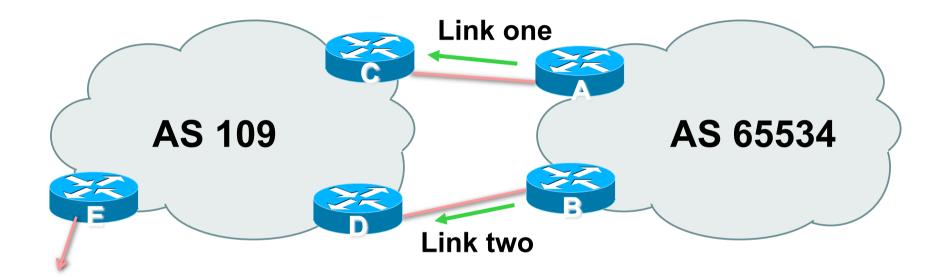
## **Private-AS Removal**

- neighbor x.x.x.x remove-private-AS
- Rules:
  - available for eBGP neighbors only
  - if the update has AS\_PATH made up of private-AS numbers, the private-AS will be dropped
  - if the AS\_PATH includes private and public AS numbers, private AS number will not be removed...it is a configuration error!
  - if AS\_PATH contains the AS number of the eBGP neighbor, the private-AS numbers will not be removed
  - if used with confederations, it will work as long as the private AS numbers are after the confederation portion of the AS\_PATH

# Two links to the same ISP

With Redundancy and Loadsharing

# Two links to the same ISP (with redundancy)



 AS109 removes private AS and any customer subprefixes from Internet announcement

## Loadsharing to the same ISP

- Announce /19 aggregate on each link
- Split /19 and announce as two /20s, one on each link

basic inbound loadsharing

assumes equal circuit capacity and even spread of traffic across address block

- Vary the split until "perfect" loadsharing achieved
- Accept the default from upstream

basic outbound loadsharing by nearest exit okay in first approx as most ISP and end-site traffic is inbound

## Two links to the same ISP

#### Router A Configuration

```
router bgp 65534
network 221.10.0.0 mask 255.255.224.0
network 221.10.0.0 mask 255.255.240.0
neighbor 222.222.10.2 remote-as 109
neighbor 222.222.10.2 prefix-list routerC out
neighbor 222.222.10.2 prefix-list default in
I
ip prefix-list default permit 0.0.0.0/0
ip prefix-list routerC permit 221.10.0.0/20
ip prefix-list routerC permit 221.10.0.0/19
ip route 221.10.0.0 255.255.240.0 null0
ip route 221.10.0.0 255.255.224.0 null0
```

#### Router B configuration is similar but with the other /20

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## Two links to the same ISP

Router C Configuration

```
router bgp 109
```

neighbor 222.222.10.1 remote-as 65534

neighbor 222.222.10.1 default-originate

neighbor 222.222.10.1 prefix-list Customer in

neighbor 222.222.10.1 prefix-list default out

```
ip prefix-list Customer permit 221.10.0.0/19 le 20
ip prefix-list default permit 0.0.0.0/0
```

- Router C only allows in /19 and /20 prefixes from customer block
- Router D configuration is identical

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## Loadsharing to the same ISP

- Loadsharing configuration is only on customer router
- Upstream ISP has to

remove customer subprefixes from external announcements

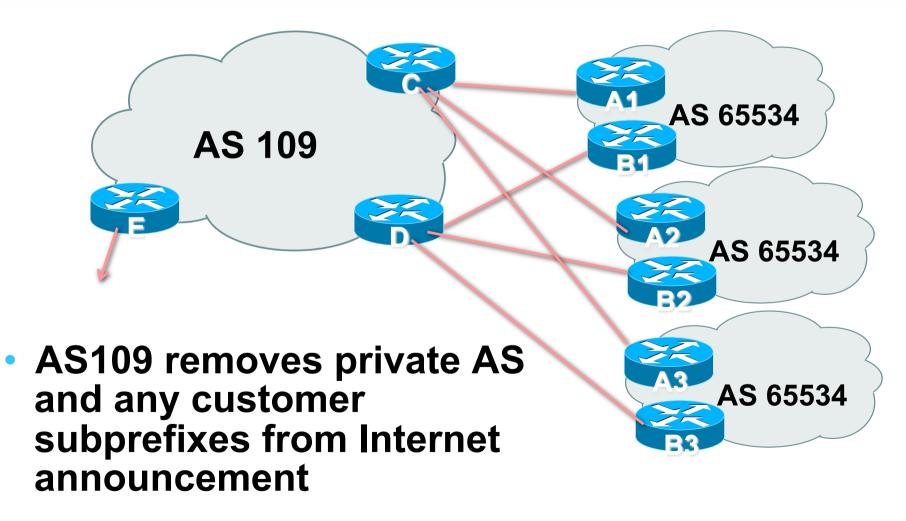
remove private AS from external announcements

Could also use BGP communities

# Two links to the same ISP

Multiple Dualhomed Customers (RFC2270)

## Multiple Dualhomed Customers (RFC2270)



- Customer announcements as per previous example
- Use the same private AS for each customer documented in RFC2270 address space is not overlapping

each customer hears default only

 Router An and Bn configuration same as Router A and B previously

## Two links to the same ISP

#### Router A1 Configuration

```
router bqp 65534
network 221.10.0.0 mask 255.255.224.0
network 221.10.0.0 mask 255.255.240.0
neighbor 222.222.10.2 remote-as 109
neighbor 222.222.10.2 prefix-list routerC out
neighbor 222.222.10.2 prefix-list default in
ļ
ip prefix-list default permit 0.0.0.0/0
ip prefix-list routerC permit 221.10.0.0/20
ip prefix-list routerC permit 221.10.0.0/19
ip route 221.10.0.0 255.255.240.0 null0
ip route 221.10.0.0 255.255.224.0 null0
```

#### Router B1 configuration is similar but for the other /20

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#### Router C Configuration

router bgp 109

neighbor bgp-customers peer-group

neighbor bgp-customers remote-as 65534

neighbor bgp-customers default-originate

neighbor bgp-customers prefix-list default out

neighbor 222.222.10.1 peer-group bgp-customers

neighbor 222.222.10.1 description Customer One

neighbor 222.222.10.1 prefix-list Customer1 in

neighbor 222.222.10.9 peer-group bgp-customers

neighbor 222.222.10.9 description Customer Two

neighbor 222.222.10.9 prefix-list Customer2 in

neighbor 222.222.10.17 peer-group bgp-customers
neighbor 222.222.10.17 description Customer Three
neighbor 222.222.10.17 prefix-list Customer3 in

- ip prefix-list Customer1 permit 221.10.0.0/19 le 20 ip prefix-list Customer2 permit 221.16.64.0/19 le 20 ip prefix-list Customer3 permit 221.14.192.0/19 le 20 ip prefix-list default permit 0.0.0.0/0
- Router C only allows in /19 and /20 prefixes from customer block
- Router D configuration is almost identical

Router E Configuration

assumes customer address space is not part of upstream's address block

```
router bgp 109
neighbor 222.222.10.17 remote-as 110
neighbor 222.222.10.17 remove-private-AS
neighbor 222.222.10.17 prefix-list Customers out
!
ip prefix-list Customers permit 221.10.0.0/19
ip prefix-list Customers permit 221.16.64.0/19
ip prefix-list Customers permit 221.14.192.0/19
```

Private AS still visible inside AS109

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 If customers' prefixes come from ISP's address block

do NOT announce them to the Internet

announce ISP aggregate only

### Router E configuration:

```
router bgp 109
neighbor 222.222.10.17 remote-as 110
neighbor 222.222.10.17 prefix-list my-aggregate out
!
ip prefix-list my-aggregate permit 221.8.0.0/13
```

## Two links to different ISPs With Redundancy

www.cisco.com

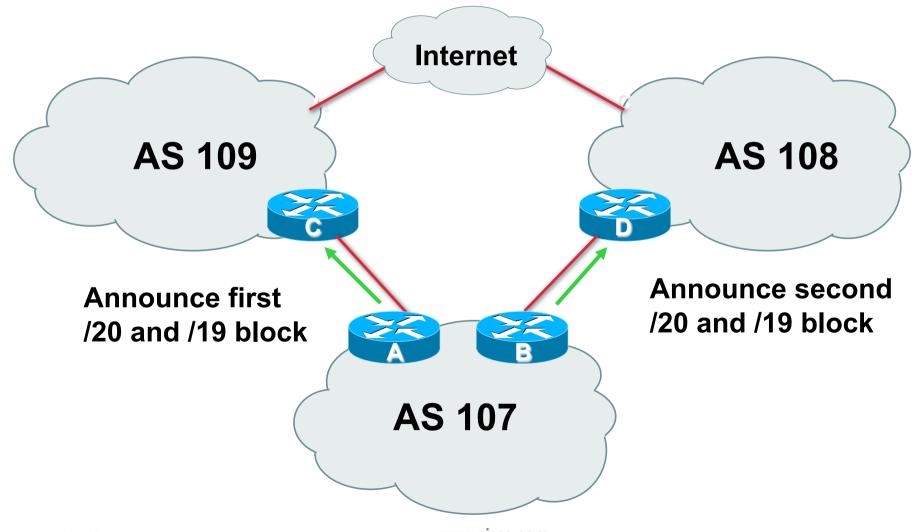
# Two links to different ISPs (with redundancy)

- Announce /19 aggregate on each link
- Split /19 and announce as two /20s, one on each link

**basic** inbound loadsharing

 When one link fails, the announcement of the /19 aggregate via the other ISP ensures continued connectivity

## Two links to different ISPs (with redundancy)



## Two links to different ISPs (with redundancy)

#### Router A Configuration

```
router bgp 107
network 221.10.0.0 mask 255.255.224.0
network 221.10.0.0 mask 255.255.240.0
neighbor 222.222.10.1 remote-as 109
neighbor 222.222.10.1 prefix-list firstblock out
neighbor 222.222.10.1 prefix-list default in
ip prefix-list default permit 0.0.0/0
ip prefix-list firstblock permit 221.10.0.0/20
ip prefix-list firstblock permit 221.10.0.0/19
```

# Two links to different ISPs (with redundancy)

#### Router B Configuration

```
router bop 107
network 221.10.0.0 mask 255.255.224.0
network 221.10.16.0 mask 255.255.240.0
neighbor 220.1.5.1 remote-as 108
neighbor 220.1.5.1 prefix-list secondblock out
neighbor 220.1.5.1 prefix-list default in
ip prefix-list default permit 0.0.0.0/0
ip prefix-list secondblock permit 221.10.16.0/20
ip prefix-list secondblock permit 221.10.0.0/19
```

## Two links to different ISPs

**More Controlled Loadsharing** 

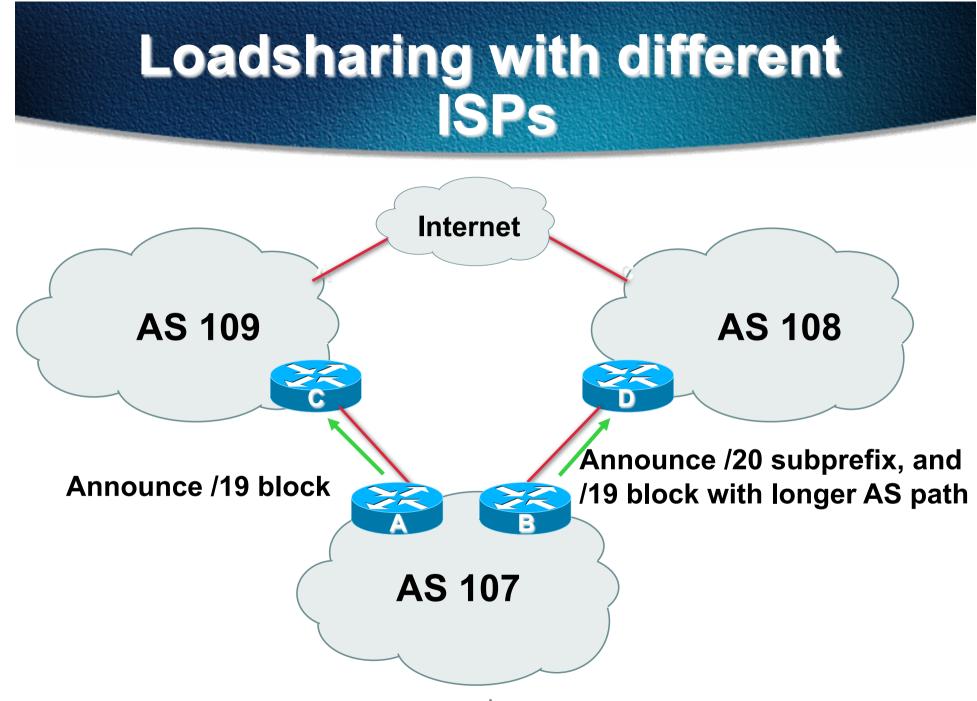
Announce /19 aggregate on each link

On first link, announce /19 as normal

On second link, announce /19 with longer AS PATH, and announce one /20 subprefix

controls loadsharing between upstreams and the Internet

- Vary the subprefix size and AS PATH length until "perfect" loadsharing achieved
- Still require redundancy!



#### Router A Configuration

router bgp 107
network 221.10.0.0 mask 255.255.224.0
neighbor 222.222.10.1 remote-as 109
neighbor 222.222.10.1 prefix-list default in
neighbor 222.222.10.1 prefix-list aggregate out
!
ip prefix-list aggregate permit 221.10.0.0/19

#### Router B Configuration

router bgp 107
network 221.10.0.0 mask 255.255.224.0
network 221.10.16.0 mask 255.255.240.0
neighbor 220.1.5.1 remote-as 108
neighbor 220.1.5.1 prefix-list default in
neighbor 220.1.5.1 prefix-list subblocks out
neighbor 220.1.5.1 route-map routerD out
!

..next slide..

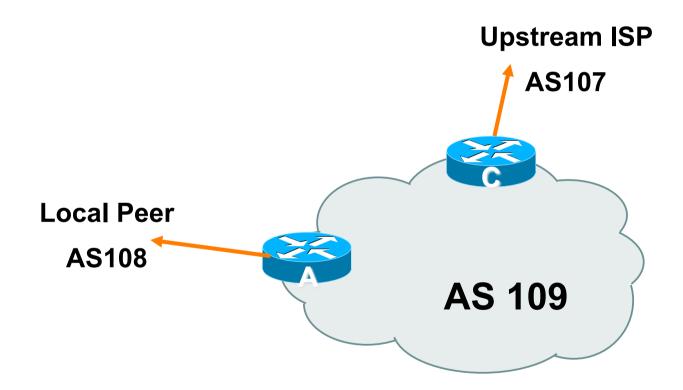
```
route-map routerD permit 10
match ip address prefix-list aggregate
set as-path prepend 107 107
route-map routerD permit 20
!
ip prefix-list subblocks permit 221.10.0.0/19 le 20
ip prefix-list aggregate permit 221.10.0.0/19
```

## Service Provider Multihoming

One Upstream, One local peer

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- Announce /19 aggregate on each link
- Accept default route only from upstream
   Either 0.0.0.0/0 or a network which can be used as default
- Accept all routes from local peer
- Border routers talk iBGP with each other



#### Router A Configuration

router bgp 109

network 221.10.0.0 mask 255.255.224.0

neighbor 222.222.10.2 remote-as 108

neighbor 222.222.10.2 prefix-list my-block out

neighbor 222.222.10.2 prefix-list AS108-peer in

```
ip prefix-list AS108-peer permit 222.5.16.0/19
ip prefix-list AS108-peer permit 221.240.0.0/20
ip prefix-list my-block permit 221.10.0.0/19
!
```

ip route 221.10.0.0 255.255.224.0 null0

#### Router C Configuration

```
router bgp 109
```

network 221.10.0.0 mask 255.255.224.0

neighbor 222.222.10.1 remote-as 107

neighbor 222.222.10.1 prefix-list default in

neighbor 222.222.10.1 prefix-list my-block out

```
ip prefix-list my-block permit 221.10.0.0/19
ip prefix-list default permit 0.0.0.0/0
!
```

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 Two configurations possible for Router A Filtering on ASes assumes peer knows what they are doing (never do this)

**Prefix-list higher maintenance, but safer** 

- Local traffic goes to and from local peer, everything else goes to upstream
- Routers A and C have minimum memory and CPU requirements

## Service Provider Multihoming Two Upstreams, One local peer

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## Two Upstreams, One Local Peer

 Two configuration options: Accept full routing from both upstreams Expensive!

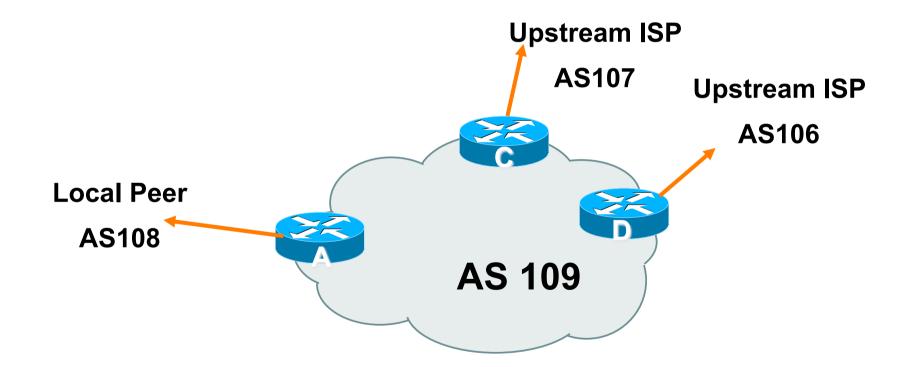
But this is the popular choice today?!!

Accept default from one upstream and some routes from the other upstream

**Best compromise, not expensive!** 

Better convergence rate and stability

## Two Upstreams, One Local Peer



#### Router A configuration is as previously

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#### Router C Configuration

```
router bgp 109
network 221.10.0.0 mask 255.255.224.0
neighbor 222.222.10.1 remote-as 107
neighbor 222.222.10.1 prefix-list rfc1918-deny in
neighbor 222.222.10.1 prefix-list my-block out
neighbor 222.222.10.1 route-map AS107-loadshare in
!
ip prefix-list my-block permit 221.10.0.0/19
! See earlier in tutorial for RFC1918 list
!
```

ip route 221.10.0.0 255.255.224.0 null0

..next slide

```
ip as-path access-list 10 permit ^(107_)+$
ip as-path access-list 10 permit ^(107_)+_[0-9]+$
!
route-map AS107-loadshare permit 10
match ip as-path 10
set local-preference 120
route-map AS107-loadshare permit 20
set local-preference 80
!
```

#### Router C configuration:

Accept full routes from AS107

Tag prefixes originated by AS107 and AS107's neighbouring ASes with local preference 120

Remaining prefixes tagged with local preference of 80

Traffic to those ASes will go over AS107 link

Traffic to other all other ASes will go over the link to AS106

Router D configuration same as Router C without the route-map

Hears full routing table!

- Full routes from upstreams
  - Expensive needs lots of memory today
  - **Expensive contributes to network instability**
  - Need to play preference games
  - Previous example is only an example real life will need improved fine-tuning!
  - Previous example doesn't consider inbound traffic see earlier slides for examples

#### Router C Configuration

```
router bop 109
network 221.10.0.0 mask 255.255.224.0
neighbor 222.222.10.1 remote-as 107
neighbor 222.222.10.1 prefix-list rfc1918-nodef-deny in
neighbor 222.222.10.1 prefix-list my-block out
neighbor 222.222.10.1 filter-list 10 in
neighbor 222.222.10.1 route-map tag-default-low in
I
ip prefix-list my-block permit 221.10.0.0/19
ip prefix-list default permit 0.0.0.0/0
! See earlier in tutorial for RFC1918 list
1
ip route 221.10.0.0 255.255.224.0 null0
```

```
ip as-path access-list 10 permit ^(107_)+$
ip as-path access-list 10 permit ^(107_)+_[0-9]+$
!
route-map tag-default-low permit 10
match ip address prefix-list default
set local-preference 80
route-map tag-default-low permit 20
```

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#### Router D Configuration

```
router bgp 109
```

network 221.10.0.0 mask 255.255.224.0

neighbor 222.222.10.5 remote-as 106

neighbor 222.222.10.5 prefix-list default in

neighbor 222.222.10.5 prefix-list my-block out

```
ip prefix-list my-block permit 221.10.0.0/19
ip prefix-list default permit 0.0.0.0/0
!
```

ip route 221.10.0.0 255.255.224.0 null0

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• Router C configuration:

Accept full routes from AS107

(or get them to send less)

Filter ASNs so only AS107 and AS107's neighbouring ASes are accepted

Allow default, and set it to local preference 80

Traffic to those ASes will go over AS107 link

Traffic to other all other ASes will go over the link to AS106

If AS106 link fails, backup via AS107 – and viceversa

Partial routes from upstreams

Not expensive – only carry the routes necessary for loadsharing

- Not expensive network more stable!
- Need to filter on AS paths

Previous example is only an example – real life will need improved fine-tuning!

Previous example doesn't consider inbound traffic – see earlier slides for examples

## BGP for Internet Service Providers

- BGP Basics (quick recap)
- Scaling BGP
- Deploying BGP in an ISP network
- Trouble & Troubleshooting
- Multihoming Examples
- Using Communities

## Communities





#### • RFC1998

#### Examples of SP applications

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- Informational RFC
- Describes how to implement loadsharing and backup on multiple inter-AS links

BGP communities used to determine local preference in upstream's network

- Gives control to the customer
- Simplifies upstream's configuration simplifies network operation!

#### Community values defined to have particular meanings:

ASx:100 set local pref 100 preferred route

- ASx:90 set local pref 90 backup route if dualhomed on ASx
  - main link is to another ISP with same AS path length
- ASx:70 set local pref 70

set local pref 80

main link is to another ISP

**ASx:80** 

#### Sample Customer Router Configuration

```
router bgp 107
neighbor x.x.x.x remote-as 109
neighbor x.x.x.x description Backup ISP
neighbor x.x.x.x route-map config-community out
neighbor x.x.x.x send-community
ip as-path access-list 20 permit ^$
ip as-path access-list 20 deny .*
route-map config-community permit 10
match as-path 20
 set community 109:90
```

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#### Sample ISP Router Configuration

! Homed to another ISP

```
ip community-list 70 permit 109:70
```

! Homed to another ISP with equal ASPATH length

```
ip community-list 80 permit 109:80
```

```
! Customer backup routes
```

```
ip community-list 90 permit 109:90
```

```
!
```

```
route-map set-customer-local-pref permit 10
match community 70
```

```
set local-preference 70
```

#### Sample ISP Router Configuration

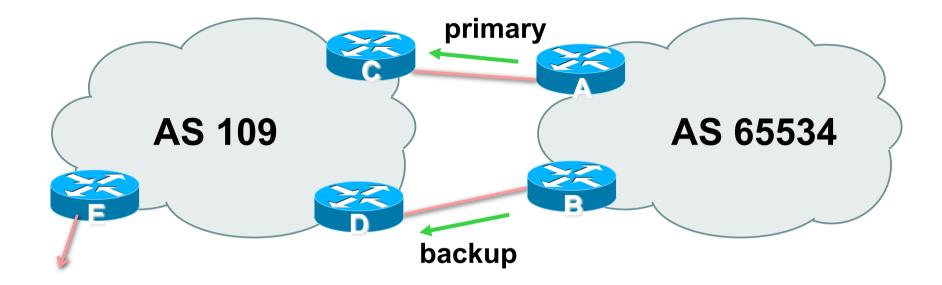
```
route-map set-customer-local-pref permit 20
match community 80
set local-preference 80
route-map set-customer-local-pref permit 30
match community 90
 set local-preference 90
route-map set-customer-local-pref permit 40
set local-preference 100
```

# Supporting RFC1998 many ISPs do, more should check AS object in the Internet Routing Registry if you do, insert comment in AS object

## Two links to the same ISP

One link primary, the other link backup only

## Two links to the same ISP



#### AS109 proxy aggregates for AS 65534

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Announce /19 aggregate on each link

primary link makes standard announcement

backup link sends community

 When one link fails, the announcement of the /19 aggregate via the other link ensures continued connectivity

#### Router A Configuration

router bgp 65534

network 221.10.0.0 mask 255.255.224.0

neighbor 222.222.10.2 remote-as 109

neighbor 222.222.10.2 description RouterC

neighbor 222.222.10.2 prefix-list aggregate out

neighbor 222.222.10.2 prefix-list default in

```
!
```

ip prefix-list aggregate permit 221.10.0.0/19
ip prefix-list default permit 0.0.0.0/0

#### Router B Configuration

router bgp 65534

network 221.10.0.0 mask 255.255.224.0

neighbor 222.222.10.6 remote-as 109

neighbor 222.222.10.6 description RouterD

neighbor 222.222.10.6 send-community

neighbor 222.222.10.6 prefix-list aggregate out

neighbor 222.222.10.6 route-map routerD-out out

neighbor 222.222.10.6 prefix-list default in

neighbor 222.222.10.6 route-map routerD-in in

!

..next slide

```
ip prefix-list aggregate permit 221.10.0.0/19
ip prefix-list default permit 0.0.0.0/0
!
route-map routerD-out permit 10
match ip address prefix-list aggregate
set community 109:90
route-map routerD-out permit 20
```

```
!
route-map routerD-in permit 10
set local-preference 90
```

#### Router C Configuration (main link)

router bgp 109
neighbor 222.222.10.1 remote-as 65534
neighbor 222.222.10.1 default-originate
neighbor 222.222.10.1 prefix-list Customer in
neighbor 222.222.10.1 prefix-list default out
!
ip prefix-list Customer permit 221.10.0.0/19
ip prefix-list default permit 0.0.0.0/0

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Router D Configuration (backup link)

```
router bgp 109
```

neighbor 222.222.10.5 remote-as 65534

neighbor 222.222.10.5 default-originate

neighbor 222.222.10.5 prefix-list Customer in

neighbor 222.222.10.5 route-map bgp-cust-in in

neighbor 222.222.10.5 prefix-list default out

```
!
```

ip prefix-list Customer permit 221.10.0.0/19
ip prefix-list default permit 0.0.0.0/0
!

..next slide

```
ip prefix-list Customer permit 221.10.0.0/19
  ip prefix-list default permit 0.0.0/0
  I
  ip community-list 90 permit 109:90
  1
<snip>
  route-map bgp-cust-in permit 30
   match community 90
   set local-preference 90
  route-map bgp-cust-in permit 40
   set local-preference 100
```

# Service Providers use of Communities

Some working examples

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#### Background

#### RFC1998 is okay for "simple" multihomed customers

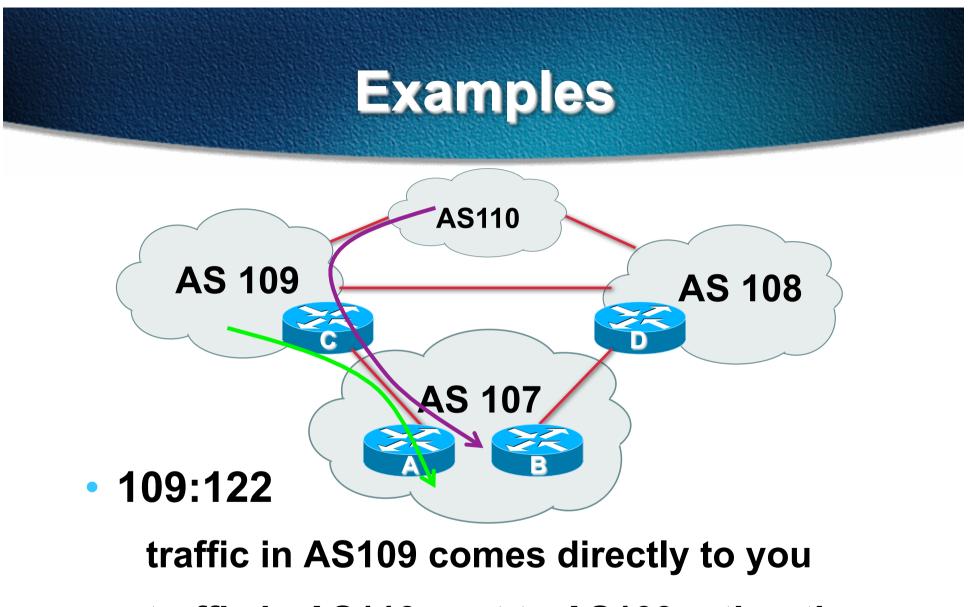
assumes that upstreams are interconnected

 ISPs create many other communities to handle more complex situations

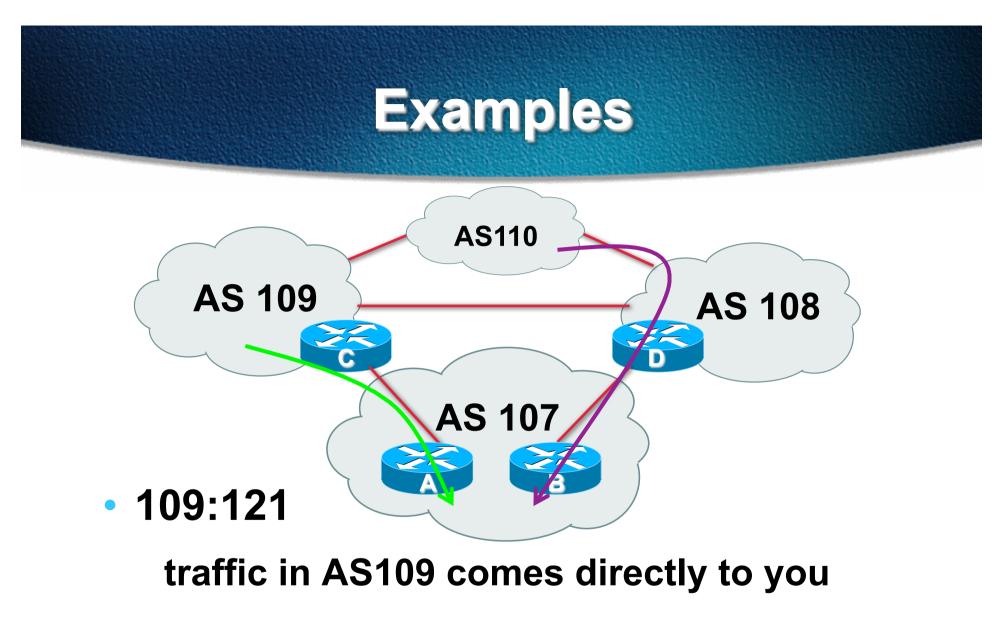
#### More community definitions

- ASx:122 set local pref 120 and set local pref high on upstreams
- ASx:121 set local pref 120 and set local pref low on upstreams
- ASx:120 set local pref 120 (opposite to ASx:80)
- ASx:82 set local pref 80 and set local pref high on upstreams
- ASx:81 set local pref 80 and set local pref low on upstreams
- ASx:21 announce to customers with no-export
- ASx:20 announce only to backbone and customers
- ASx:3 set 3x as-path prepend on peer announcement
- ASx:2 set 2x as-path prepend on peer announcement
- ASx:1 set 1x as-path prepend on peer announcement

(and variations on this theme depending on local conditions, e.g. IXPs, domestic vs. international transit, etc.)



### traffic in AS110 sent to AS109 rather than best path



### traffic in AS110 sent to AS108 rather than best path

#### Examples

• 109:3

### prepend any announcements to peers of AS109 with 109\_109\_109

"AS109 is my backup transit AS"

• 109:20

### Don't announce outside upstream's customer base

"AS109 provides local connections only"

109:21 is very similar

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#### BGP for Internet Service Providers

- BGP Basics (quick recap)
- Scaling BGP
- Deploying BGP in an ISP network
- Trouble & Troubleshooting
- Multihoming Examples
- Using Communities

### BGP for Internet Service Providers End of Tutorial

