## IPv6 Address Planning

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### Presentation Slides

#### Available on

- http://thyme.apnic.net/ftp/seminars/MyNOG2-IPv6-Address-Planning.pdf
- And on the MyNOG2 website

#### Feel free to ask questions any time

## Address Planning

IPv6 address space available to each network operator is very large compared with IPv4

- Design a scalable plan
- Be aware of industry current practices
- Separation of infrastructure and customer addressing
- Distribution of address space according to function

### Why Create an Addressing Plan?

- The options for an IPv4 addressing plan were severely limited:
  - Because of scarcity of addresses
  - Every address block has to be used efficiently
- IPv6 allows for a scalable addressing plan:
  - Security policies are easier to implement
  - Addresses are easier to trace
  - Enables more efficient network management

### Nibble Boundaries

- IPv6 offers network operators more flexibility with addressing plans
  - Network addressing can now be done on nibble boundaries

• For ease of operation

Rather than making maximum use of a very scarce resource

**•** With the resulting operational complexity

- A nibble boundary means subdividing address space based on the address numbering
  - Each number in IPv6 represents 4 bits
  - Which means that IPv6 addressing can be done on 4-bit boundaries

### Nibble Boundaries – example

Consider the address block 2001:db8:0:10::/61

The range of addresses in this block are:

2001:0db8:0000:0010:0000:0000:0000 to 2001:0db8:0000:0017:ffff:ffff:ffff

- Note that this subnet only runs from 0010 to 0017.
- The adjacent block is 2001:db8:0:18::/61

2001:0db8:0000:0018:0000:0000:0000:0000 to 2001:0db8:0000:001f:ffff:ffff:ffff

The address blocks don't use the entire nibble range

### Nibble Boundaries – example

- Now consider the address block 2001:db8:0:10::/60
  - The range of addresses in this block are:

2001:0db8:0000:0010:0000:0000:0000:0000 to 2001:0db8:0000:001f:ffff:ffff:ffff

- Note that this subnet uses the entire nibble range, 0 to f
- Which makes the numbering plan for IPv6 simpler
  - This range can have a particular meaning within the ISP block (for example, infrastructure addressing for a particular PoP)

### Addressing Plans – Infrastructure

- Network Operators should procure a /32 from their RIR
- Address block for infrastructure
  - /48 allows 65k subnets in the backbone
- Address block for router loop-back interfaces
  - Number all loopbacks out of one infrastructure /64
  - /128 per loopback
- Point-to-point links
  - /64 reserved for each, address as a /127
- LANs
  - /64 for each LAN

### Addressing Plans – Customer

Customers get one /48

 Unless they have more than 65k subnets in which case they get a second /48 (and so on)

#### In typical deployments today:

- Several ISPs are giving small customers a /56 and single LAN end-sites a /64, e.g.:
  - /64 if end-site will only ever be a LAN
  - /56 for small end-sites (e.g. home/office/small business)
  - /48 for large end-sites
- This is another very active discussion area
- Observations:
  - Don't assume that a mobile endsite needs only a /64
  - Some operators are distributing /60s to their smallest customers!!

Documentation

- IPv4 addresses are probably short enough to memorise
- IPv6 addresses are unlikely to be memorable at all
- Document the address plan
  - What is used for infrastructure
  - What goes to customers
  - Flat file, spreadsheet, database, etc
  - But documentation is vital
  - Especially when coming to populating the DNS later on

### Addressing Tools

#### Examples of IP address planning tools:

- NetDot netdot.uoregon.edu (recommended!!)
- HaCi sourceforge.net/projects/haci
- IPAT nethead.de/index.php/ipat
- freeipdb home.globalcrossing.net/~freeipdb/
- Examples of IPv6 subnet calculators:
  - ipv6gen code.google.com/p/ipv6gen/
  - sipcalc www.routemeister.net/projects/sipcalc/

Pick the first /48 for our infrastructure

- Reason: keeps the numbers short
- Short numbers: less chance of transcription errors
- Compare:

2001:db8:ef01:d35c::1/128

with

2001:db8::1/128

For Loopback interface addresses

Out of this /48, pick the first /64 for loopbacks

- Reason: keeps the numbers short
- Some operators use first /64 for anycast services

- Pick the second /48 for point-to-point links to customers
  - Addresses not a trusted part of Operator's infrastructure
- Divide the /48 between PoPs
  - e.g. 10 PoPs  $\rightarrow$  split into /52s  $\rightarrow$  4096 links per /52
  - Gives 65536 /64s for 65536 customer links
     /64 per link, number as /127 as previously
  - Adjust number of /48s to suit PoP size (one /48 per PoP?)
  - Aggregate per router or per PoP and carry in iBGP
- Alternative is to use unnumbered interfaces

#### ■ For the infrastructure /48:

- First /64 for loopbacks
- Maybe reserve the final /60 for the NOC
  - Gives 16 possible subnets for the Network Operations Centre (part of the Infrastructure)
- Remaining 65519 /64s used for internal pointto-point links

More than any network needs today

### Example: Loopback addresses

- 2001:db8:0::/48 is used for infrastructure
- Out of this, 2001:db8:0:0::/64 is used for loopbacks
- Network Operator has 20 PoPs
  - Scheme adopted is 2001:db8::XXYY/128
  - Where X is the PoP number (1 through FF)
  - Where Y is the router number (1 through FF)
  - Scheme is good for 255 PoPs with 255 routers per PoP, and keeps addresses small/short

### Example: Loopback addresses

- Loopbacks in PoP 1:
  - CR1 2001:db8::101/128
  - CR2 2001:db8::102/128
  - BR1 2001:db8::103/128
  - BR2 2001:db8::104/128
  - AR1 2001:db8::110/128
  - AR2 2001:db8::111/128
  - AR3 2001:db8::112/128
  - AR4 2001:db8::113/128

...etc...

Loopbacks in PoP 10:

- CR1 2001:db8::a01/128
- CR2 2001:db8::a02/128
- BR1 2001:db8::a03/128
- BR2 2001:db8::a04/128
- AR1 2001:db8::a10/128
- AR2 2001:db8::a11/128
- AR3 2001:db8::a12/128
- AR4 2001:db8::a13/128

...etc...

### Example: Backbone Point-to-Point links

#### ISP has 20 PoPs

Scheme adopted is 2001:db8:0:XXYY::Z/64

#### Where:

- XX is the PoP number (01 through FF)
- YY is the LAN number (when YY is 00 through 0F)
- YY is the P2P link number (when YY is 10 through FF)
- Z is the interface address
- /64 is reserved, but the link is numbered as a /127
- Scheme is good for 16 LANs and 240 backbone PtP links per PoP, and for 255 PoPs

#### Example: Backbone Point-to-Point links

	PtP	&	LANs	in	PoP	1:	
--	-----	---	------	----	-----	----	--

- LAN1 2001:db8:0:100::/64
- LAN2 2001:db8:0:101::/64
- LAN3 2001:db8:0:102::/64
- PtP1 2001:db8:0:110::/64
- PtP2 2001:db8:0:111::/64
- PtP3 2001:db8:0:112::/64
- PtP4 2001:db8:0:113::/64
- PtP5 2001:db8:0:114::/64

...etc...

- PtP & LANs in PoP 14:
  - LAN1 2001:db8:0:e00::/64
  - LAN2 2001:db8:0:e01::/64
  - LAN3 2001:db8:0:e02::/64
  - LAN4 2001:db8:0:e03::/64
  - LAN5 2001:db8:0:e04::/64
  - PtP1 2001:db8:0:e10::/64
  - PtP2 2001:db8:0:e11::/64
  - PtP3 2001:db8:0:e12::/64

...etc...

### Links to Customers (1)

- Some Network Operators use unnumbered IPv4 interface links
  - So replicate this in IPv6 by using unnumbered IPv6 interface links
  - This will not require one /48 to be taken from the ISP's /32 allocation

### Links to Customers (2)

- Other Network Operators use global unicast addresses
  - So set aside the second /48 for this purpose
    - And divide the /48 amongst the PoPs
  - Or set aside a single/48 per PoP (depending on network size)
  - Each /48 gives 65536 possible customer links, assuming a /64 for each link

Scheme used:

- 2001:db8:00XX::/48 where XX is the PoP number
- Good for 255 PoPs with 65536 point-to-point links each

### Example

#### Customer PtP links

- PoP1
  - Reserved 2001:db8:1:0::/64
    Customer1 2001:db8:1:1::/64
    Customer2 2001:db8:1:2::/64
    Customer3 2001:db8:1:3::/64
    Customer4 2001:db8:1:4::/64
- PoP12
  - Reserved 2001:db8:c:0::/64
     Customer1 2001:db8:c:1::/64
     Customer2 2001:db8:c:2::/64
     Customer3 2001:db8:c:3::/64
- ...etc...

### Example: Customer Allocations

- Master allocation documentation would look like this:
  - 2001:db8:0::/48
  - 2001:db8:1::/48
  - 2001:db8:2::/48
  - 2001:db8:3::/48

. . .

- Infrastructure
- PtP links to customers (PoP1) PtP links to customers (PoP2) PtP links to customers (PoP3)

2001:db8:100::/48

Customer 1 assignment

- 2001:db8:ffff::/48 Customer 65280 assignment
- Infrastructure and Customer PtP links would be documented separately as earlier

### Addressing Plans – Customer

Geographical delegations to Customers:

- Network Operator subdivides /32 address block into geographical chunks
- E.g. into /36s
  - Region 1: 2001:db8:1xxx::/36
  - **Region 2: 2001:db8:2xxx::/36**
  - Region 3: 2001:db8:3xxx::/36
  - etc
- Which gives 4096 /48s per region
- For Operational and Administrative ease
- Benefits for traffic engineering if Network Operator multihomes in each region

### Addressing Plans – Customer

Sequential delegations to Customers:

- After carving off address space for network infrastructure, Network Operator simply assigns address space sequentially
- E.g:

Infrastructure:	2001:db8:0::/48
Customer P2P:	2001:db8:1::/48
Customer 1:	2001:db8:2::/48
Customer 2:	2001:db8:3::/48

etc

 Useful when there is no regional subdivision of network and no regional multihoming needs

## Addressing Plans – Traffic Engineering

#### Smaller providers will be single homed

The customer portion of the ISP's IPv6 address block will usually be assigned sequentially

#### Larger providers will be multihomed

- Two, three or more external links from different providers
- Traffic engineering becomes important
- Sequential assignments of customer addresses will negatively impact load balancing

## Addressing Plans – Traffic Engineering

- ISP Router loopbacks and backbone point-topoint links make up a small part of total address space
  - And they don't attract traffic, unlike customer address space
- Links from ISP Aggregation edge to customer router needs one /64
  - Small requirements compared with total address space
  - Some ISPs use IPv6 unnumbered
- Planning customer assignments is a very important part of multihoming
  - Traffic engineering involves subdividing aggregate into pieces until load balancing works



ISP fills up customer IP addressing from one end of the range:

2001:db8::/32

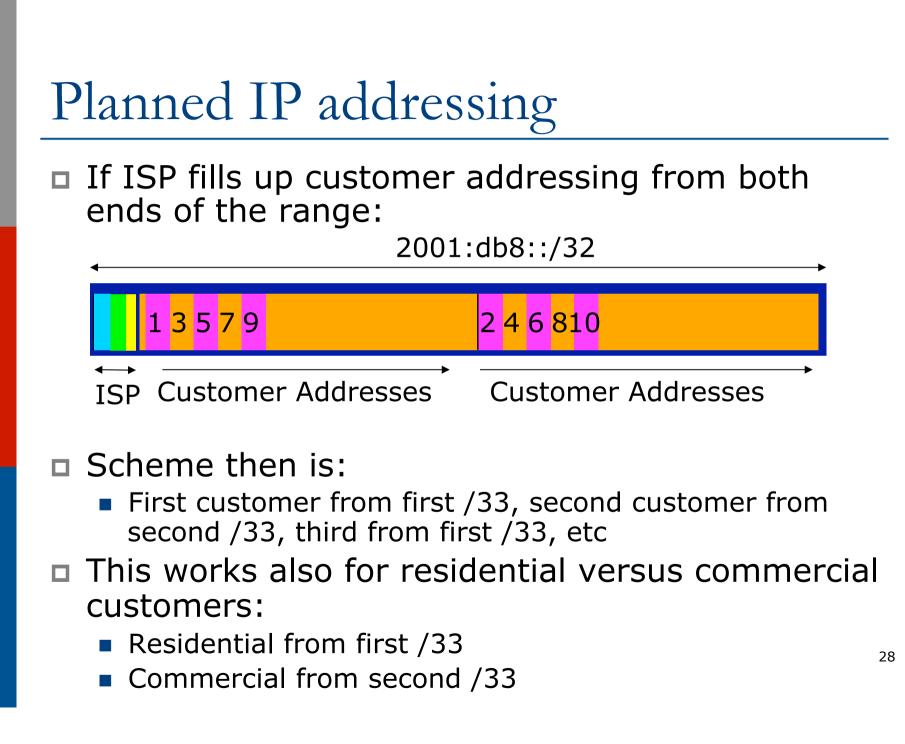
1234

**I**SP

**Customer Addresses** 

#### Customers generate traffic

- Dividing the range into two pieces will result in one /33 with all the customers and the ISP infrastructure the addresses, and one /33 with nothing
- No loadbalancing as all traffic will come in the first /33
- Means further subdivision of the first /33 = harder work



### Planned IP Addressing

- This works fine for multihoming between two upstream links (same or different providers)
- Can also subdivide address space to suit more than two upstreams
  - Follow a similar scheme for populating each portion of the address space
- Consider regional (geographical) distribution of customer delegated address space
- Don't forget to always announce an aggregate out of each link

### Addressing Plans – Advice

Customer address assignments should not be reserved or assigned on a per PoP basis

- Follow same principle as for IPv4
- Subnet aggregate to cater for multihoming needs
- Consider geographical delegation scheme
- ISP iBGP carries customer nets
- Aggregation within the iBGP not required and usually not desirable
- Aggregation in eBGP is very necessary
- Backbone infrastructure assignments:
  - Number out of a single /48
    - Operational simplicity and security
  - Aggregate to minimise size of the IGP

### Summary

# Defined structure within IPv6 addressing is recommended

- Greater flexibility than with IPv4
- Possible to come up with a simple memorable scheme
- Documentation vitally important!