

# BGP Multihoming Techniques



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MyNOG 2

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# Presentation Slides

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- Available on
  - <http://thyme.apnic.net/ftp/seminars/MyNOG2-Multihoming.pdf>
  - And on the MyNOG2 website
- Feel free to ask questions any time



# BGP Multihoming Techniques

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- Why Multihome?
- Definition & Options
- How to Multihome
- Principles & Addressing
- Basic Multihoming
- Service Provider Multihoming

# Why Multihome?



It's all about redundancy,  
diversity & reliability

# Why Multihome?

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## □ Redundancy

- One connection to internet means the network is dependent on:
  - Local router (configuration, software, hardware)
  - WAN media (physical failure, carrier failure)
  - Upstream Service Provider (configuration, software, hardware)

# Why Multihome?

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## □ Reliability

- Business critical applications demand continuous availability
- Lack of redundancy implies lack of reliability  
implies loss of revenue

# Why Multihome?

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## □ Supplier Diversity

- Many businesses demand supplier diversity as a matter of course
- Internet connection from two or more suppliers
  - With two or more diverse WAN paths
  - With two or more exit points
  - With two or more international connections
  - **Two of everything**

# Why Multihome?

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- ❑ Not really a reason, but oft quoted...
- ❑ Leverage:
  - Playing one ISP off against the other for:
    - ❑ Service Quality
    - ❑ Service Offerings
    - ❑ Availability



# Why Multihome?

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## □ Summary:

- Multihoming is easy to demand as requirement for any service provider or end-site network
- But what does it really mean:
  - In real life?
  - For the network?
  - For the Internet?
- And how do we do it?



# BGP Multihoming Techniques

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- ❑ Why Multihome?
- ❑ Definition & Options
- ❑ How to Multihome
- ❑ Principles & Addressing
- ❑ Basic Multihoming
- ❑ Service Provider Multihoming

# Multihoming: Definitions & Options



What does it mean, what do we  
need, and how do we do it?

# Multihoming Definition

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- ❑ More than one link external to the local network
  - two or more links to the same ISP
  - two or more links to different ISPs
- ❑ Usually **two** external facing routers
  - one router gives link and provider redundancy only

# Autonomous System Number (ASN)

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- ❑ Two ranges
  - 0-65535 (original 16-bit range)
  - 65536-4294967295 (32-bit range – RFC4893)
- ❑ Usage:
  - 0 and 65535 (reserved)
  - 1-64495 (public Internet)
  - 64496-64511 (documentation – RFC5398)
  - 64512-65534 (private use only)
  - 23456 (represent 32-bit range in 16-bit world)
  - 65536-65551 (documentation – RFC5398)
  - 65552-4294967295 (public Internet)
- ❑ 32-bit range representation specified in RFC5396
  - Defines “asplain” (traditional format) as standard notation

# Autonomous System Number (ASN)

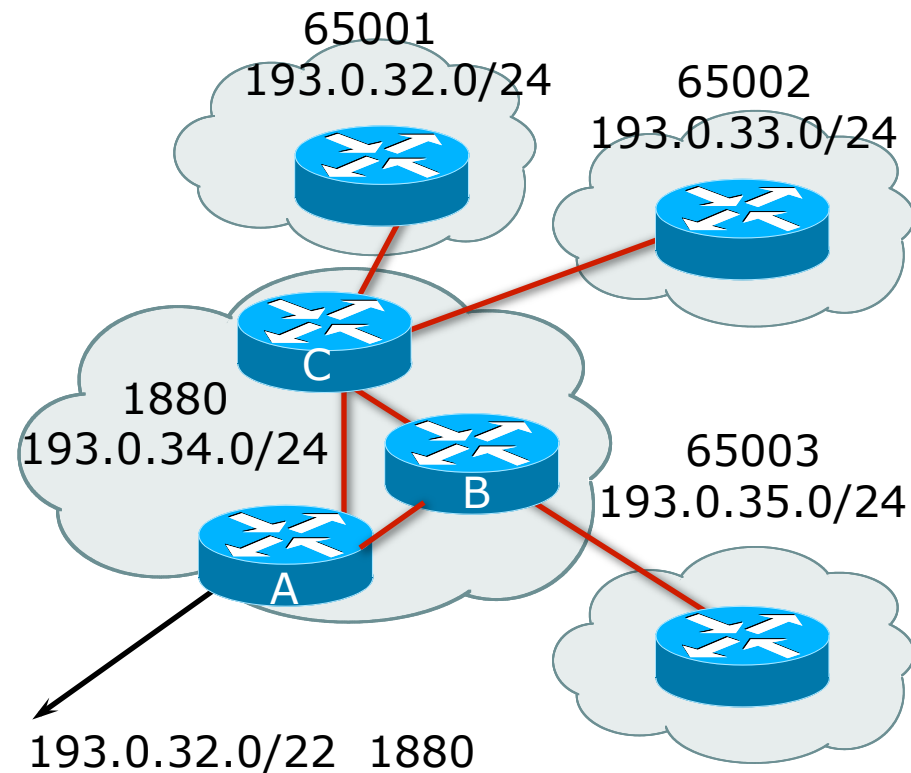
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- ❑ ASNs are distributed by the Regional Internet Registries
  - They are also available from upstream ISPs who are members of one of the RIRs
  - Around 43000 are visible on the Internet
- ❑ Current 16-bit ASN allocations up to 61439 have been made to the RIRs
- ❑ Each RIR has also received a block of 32-bit ASNs
  - Out of 3500 assignments, around 3100 are visible on the Internet
- ❑ See [www.iana.org/assignments/as-numbers](http://www.iana.org/assignments/as-numbers)

# Private-AS – Application

## □ Applications

- An ISP with customers multihomed on their backbone (RFC2270)  
-or-
- A corporate network with several regions but connections to the Internet only in the core  
-or-
- Within a BGP Confederation



# Private-AS – Removal

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- ❑ Private ASNs MUST be removed from all prefixes announced to the public Internet
  - Include configuration to remove private ASNs in the eBGP template
- ❑ As with RFC1918 address space, private ASNs are intended for internal use
  - They should not be leaked to the public Internet



# Transit/Peering/Default

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## □ Transit

- Carrying traffic across a network
- Usually **for a fee**

## □ Peering

- Exchanging locally sourced routing information and traffic
- Usually **for no fee**
- Sometimes called settlement free peering

## □ Default

- Where to send traffic when there is no explicit match in the routing table



# Policy Tools

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- Local preference
  - outbound traffic flows
- Metric (MED)
  - inbound traffic flows (local scope)
- AS-PATH prepend
  - inbound traffic flows (Internet scope)
- Communities
  - specific inter-provider peering

# Originating Prefixes: Assumptions

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- ❑ **MUST** announce assigned address block to Internet
- ❑ MAY also announce subprefixes – reachability is not guaranteed
- ❑ Current minimum IPv4 allocation ranges from /20 to /24 depending on the RIR
  - Several ISPs filter RIR blocks on this boundary
  - Several ISPs filter the rest of address space according to the IANA assignments
  - This activity is called “Net Police” by some

# Originating Prefixes

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- ❑ The RIRs publish their minimum allocation sizes per /8 address block
  - AfriNIC: [www.afrinic.net/docs/policies/afpol-v4200407-000.htm](http://www.afrinic.net/docs/policies/afpol-v4200407-000.htm)
  - APNIC: [www.apnic.net/db/min-alloc.html](http://www.apnic.net/db/min-alloc.html)
  - ARIN: [www.arin.net/reference/ip\\_blocks.html](http://www.arin.net/reference/ip_blocks.html)
  - LACNIC: [lacnic.net/en/registro/index.html](http://lacnic.net/en/registro/index.html)
  - RIPE NCC: [www.ripe.net/ripe/docs/smallest-alloc-sizes.html](http://www.ripe.net/ripe/docs/smallest-alloc-sizes.html)
  - Note that AfriNIC only publishes its current minimum allocation size, not the allocation size for its address blocks
- ❑ IANA publishes the address space it has assigned to end-sites and allocated to the RIRs:
  - [www.iana.org/assignments/ipv4-address-space](http://www.iana.org/assignments/ipv4-address-space)
- ❑ Several ISPs use this published information to filter prefixes on:
  - What should be routed (from IANA)
  - The minimum allocation size from the RIRs

# “Net Police” prefix list issues

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- ❑ Meant to “punish” ISPs who pollute the routing table with specifics rather than announcing aggregates
- ❑ Impacts legitimate multihoming especially at the Internet’s edge
- ❑ Impacts regions where domestic backbone is unavailable or costs \$\$\$ compared with international bandwidth
- ❑ Hard to maintain – requires updating when RIRs start allocating from new address blocks
- ❑ Don’t do it unless consequences understood and you are prepared to keep the list current
  - Consider using the Team Cymru or other reputable bogon BGP feed:
  - [www.team-cymru.org/Services/Bogons/routeserver.html](http://www.team-cymru.org/Services/Bogons/routeserver.html)

# BGP Multihoming Techniques

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- ❑ Why Multihome?
- ❑ Definition & Options
- ❑ How to Multihome
- ❑ Principles & Addressing
- ❑ Basic Multihoming
- ❑ Service Provider Multihoming

# How to Multihome



Scenarios



# Multihoming Scenarios

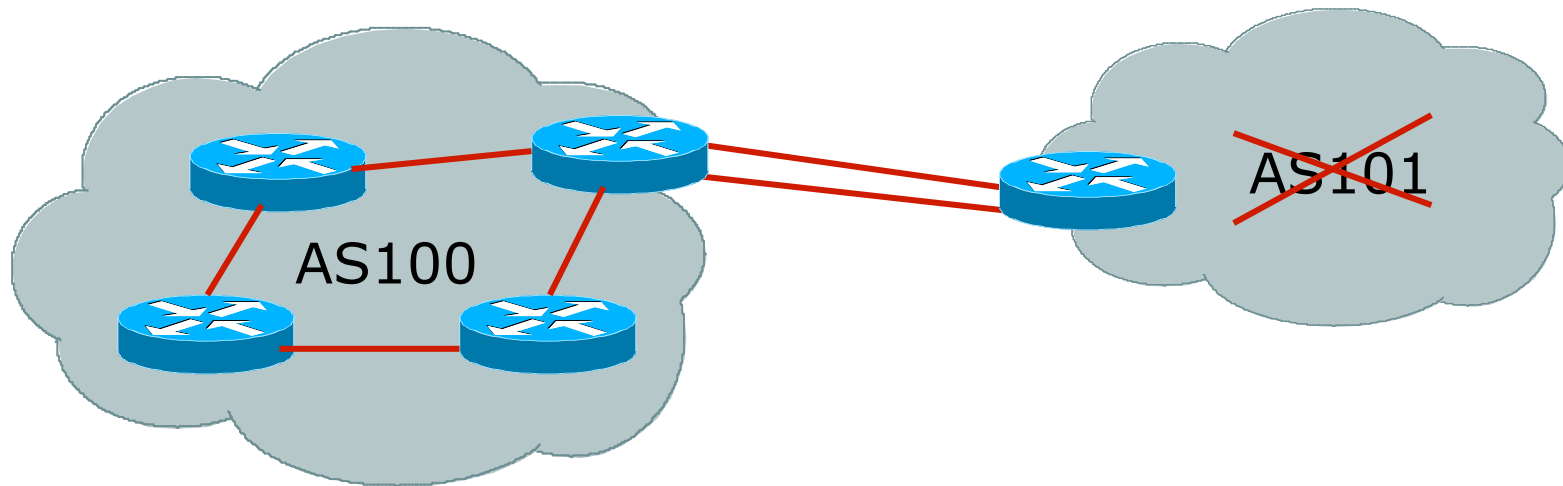
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- ❑ Stub network
- ❑ Multi-homed stub network
- ❑ Multi-homed network
- ❑ Multiple sessions to another AS



# Stub Network

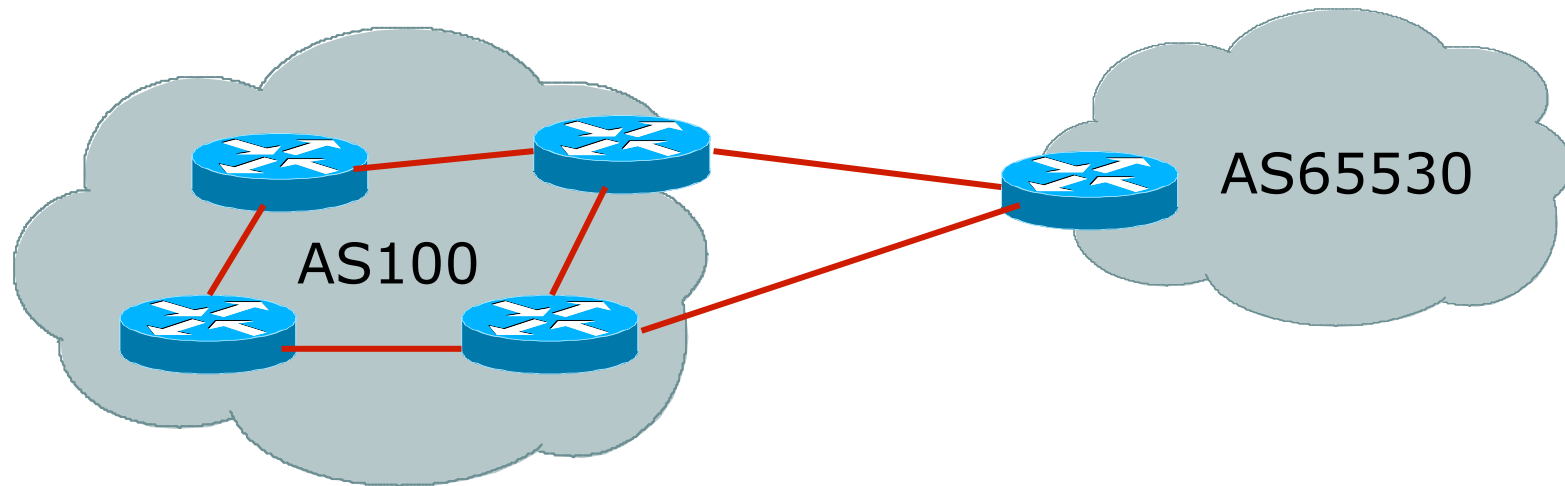
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- ❑ No need for BGP
- ❑ Point static default to upstream ISP
- ❑ Upstream ISP advertises stub network
- ❑ Policy confined within upstream ISP's policy

# Multi-homed Stub Network

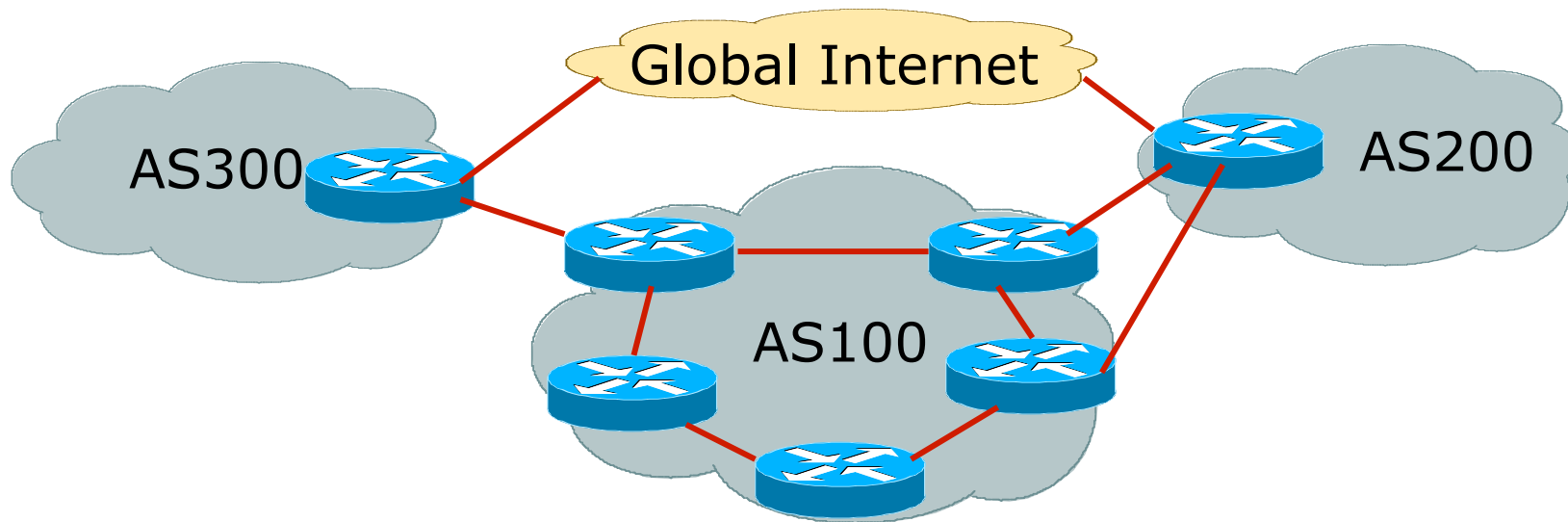
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- ❑ Use BGP (not IGP or static) to loadshare
- ❑ Use private AS (ASN > 64511)
- ❑ Upstream ISP advertises stub network
- ❑ Policy confined within upstream ISP's policy

# Multi-homed Network

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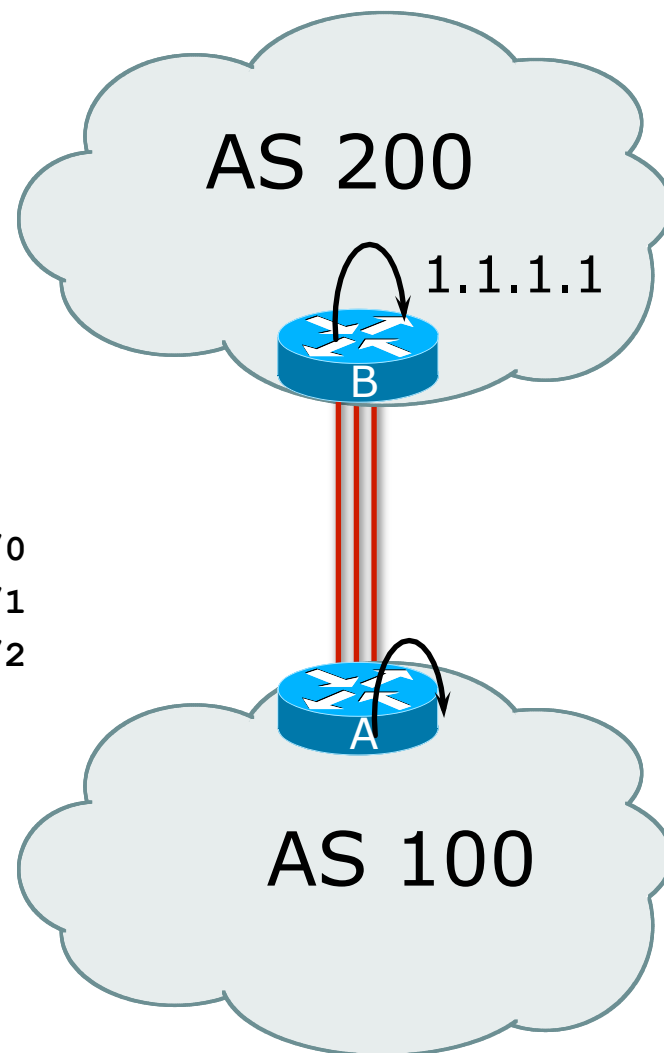
- ❑ Many situations possible
  - multiple sessions to same ISP
  - secondary for backup only
  - load-share between primary and secondary
  - selectively use different ISPs

# Multiple Sessions to an AS

## – ebgp multihop

- ❑ Use ebgp-multihop
  - Run eBGP between loopback addresses
  - eBGP prefixes learned with loopback address as next hop
- ❑ Cisco IOS

```
router bgp 100
  neighbor 1.1.1.1 remote-as 200
  neighbor 1.1.1.1 ebgp-multihop 2
  !
  ip route 1.1.1.1 255.255.255.255 serial 1/0
  ip route 1.1.1.1 255.255.255.255 serial 1/1
  ip route 1.1.1.1 255.255.255.255 serial 1/2
```
- ❑ Common error made is to point remote loopback route at IP address rather than specific link



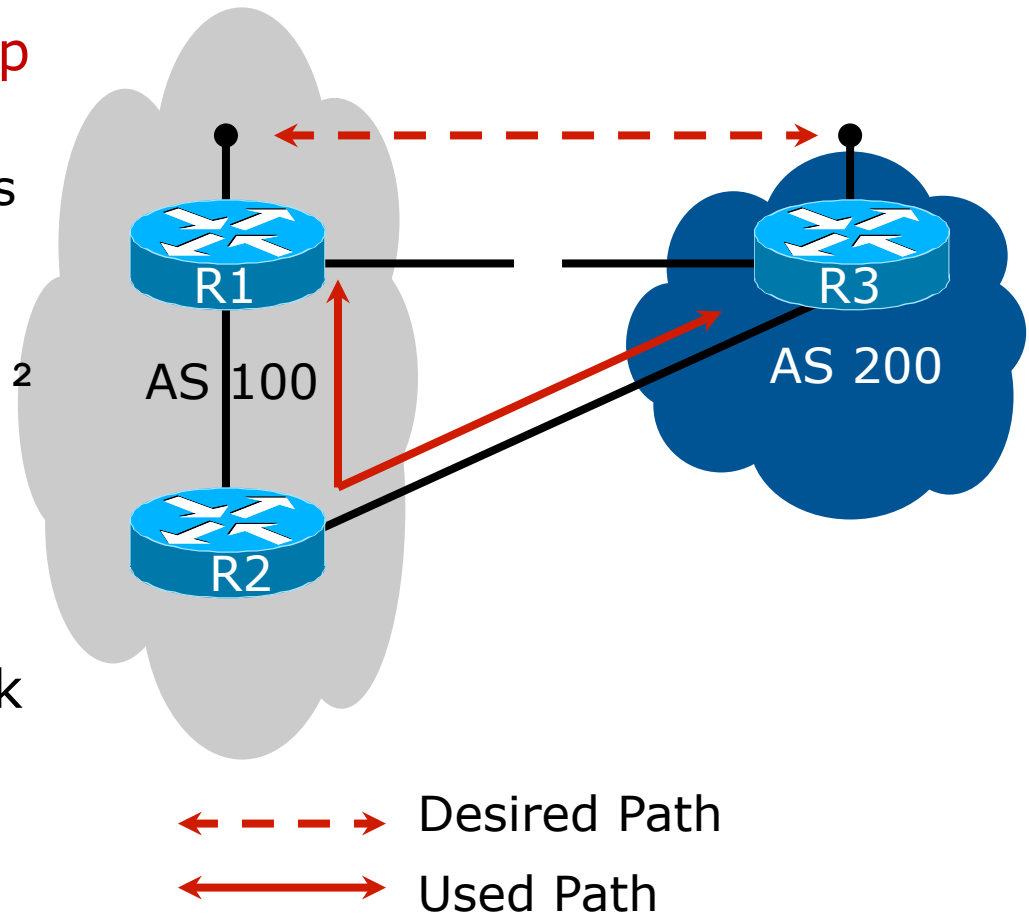
# Multiple Sessions to an AS

## – ebgp multihop

- ❑ One serious eBGP-multihop caveat:

- R1 and R3 are eBGP peers that are loopback peering
- Configured with:  
`neighbor x.x.x.x ebgp-multihop 2`
- If the R1 to R3 link goes down the session could establish via R2

- ❑ Usually happens when routing to remote loopback is dynamic, rather than static pointing at a link



# Multiple Sessions to an ISP

## – ebgp multihop

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- ❑ Try and avoid use of ebgp-multihop unless:
  - It's absolutely necessary –or–
  - Loadsharing across multiple links
- ❑ Many ISPs discourage its use, for example:

We will run eBGP multihop, but do not support it as a standard offering because customers generally have a hard time managing it due to:

- routing loops
- failure to realise that BGP session stability problems are usually due connectivity problems between their CPE and their BGP speaker

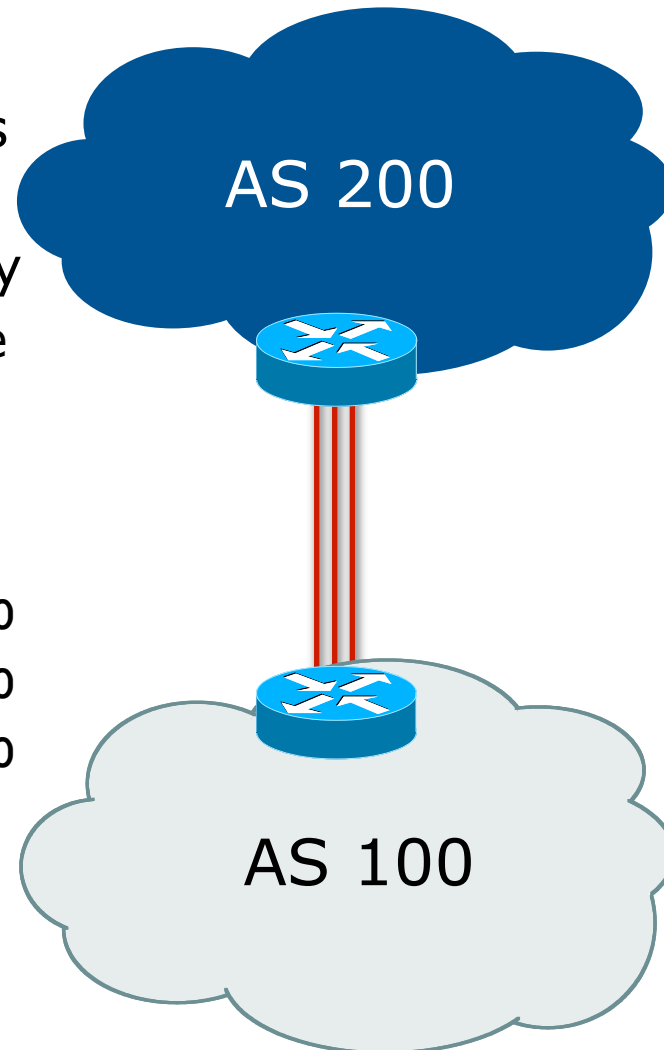
# Multiple Sessions to an AS

## – bgp multi path

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- ❑ Three BGP sessions required
- ❑ Platform limit on number of paths (could be as little as 6)
- ❑ Full BGP feed makes this unwieldy
  - 3 copies of Internet Routing Table goes into the FIB

```
router bgp 100
  neighbor 1.1.2.1 remote-as 200
  neighbor 1.1.2.5 remote-as 200
  neighbor 1.1.2.9 remote-as 200
  maximum-paths 3
```

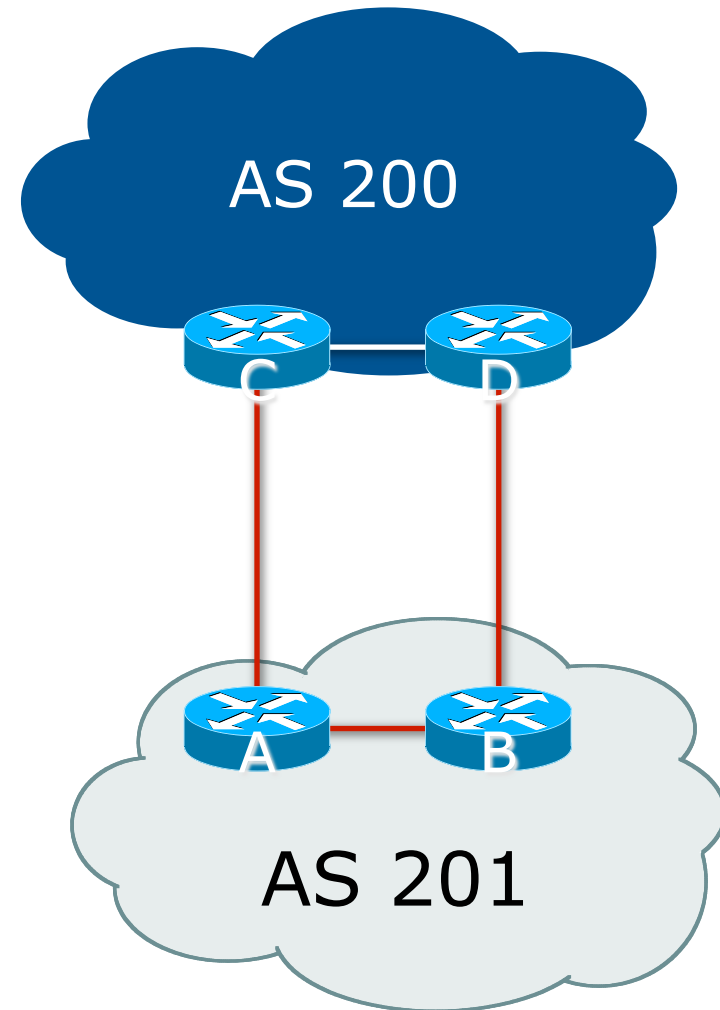


# Multiple Sessions to an AS

## – bgp attributes & filters

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- ❑ Simplest scheme is to use defaults
- ❑ Learn/advertise prefixes for better control
- ❑ Planning and some work required to achieve loadsharing
  - Point default towards one ISP
  - Learn selected prefixes from second ISP
  - Modify the number of prefixes learnt to achieve acceptable load sharing
- ❑ **No magic solution**







# BGP Multihoming Techniques

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- ❑ Why Multihome?
- ❑ Definition & Options
- ❑ How to Multihome
- ❑ Principles & Addressing
- ❑ Basic Multihoming
- ❑ Service Provider Multihoming

# Basic Principles of Multihoming



Let's learn to walk before we try  
running...

# The Basic Principles

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- ❑ Announcing address space attracts traffic
  - (Unless policy in upstream providers interferes)
- ❑ Announcing the ISP aggregate out a link will result in traffic for that aggregate coming in that link
- ❑ Announcing a subprefix of an aggregate out a link means that all traffic for that subprefix will come in that link, even if the aggregate is announced somewhere else
  - The most specific announcement wins!

# The Basic Principles

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- ❑ To split traffic between two links:
  - Announce the aggregate on both links - ensures redundancy
  - Announce one half of the address space on each link
  - (This is the first step, all things being equal)
- ❑ Results in:
  - Traffic for first half of address space comes in first link
  - Traffic for second half of address space comes in second link
  - If either link fails, the fact that the aggregate is announced ensures there is a backup path

# The Basic Principles

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- ❑ The keys to successful multihoming configuration:
  - Keeping traffic engineering prefix announcements independent of customer iBGP
  - Understanding how to announce aggregates
  - Understanding the purpose of announcing subprefixes of aggregates
  - Understanding how to manipulate BGP attributes
  - Too many upstreams/external paths makes multihoming harder (2 or 3 is enough!)

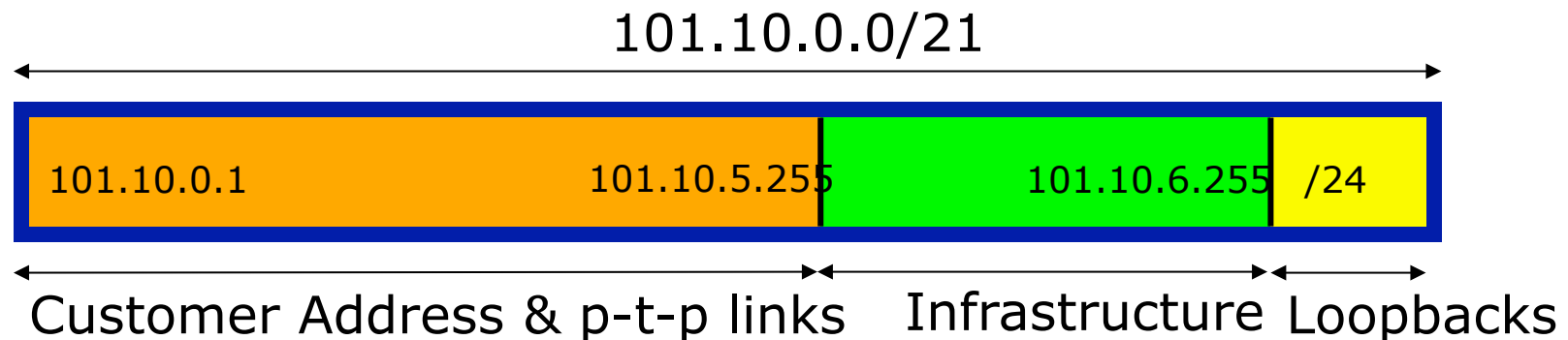
# IP Addressing & Multihoming



How Good IP Address Plans  
assist with Multihoming

# IP Addressing & Multihoming

- ❑ IP Address planning is an important part of Multihoming
- ❑ This means separating:
  - Customer address space
  - Customer p-t-p link address space
  - Infrastructure p-t-p link address space
  - Loopback address space



# IP Addressing & Multihoming

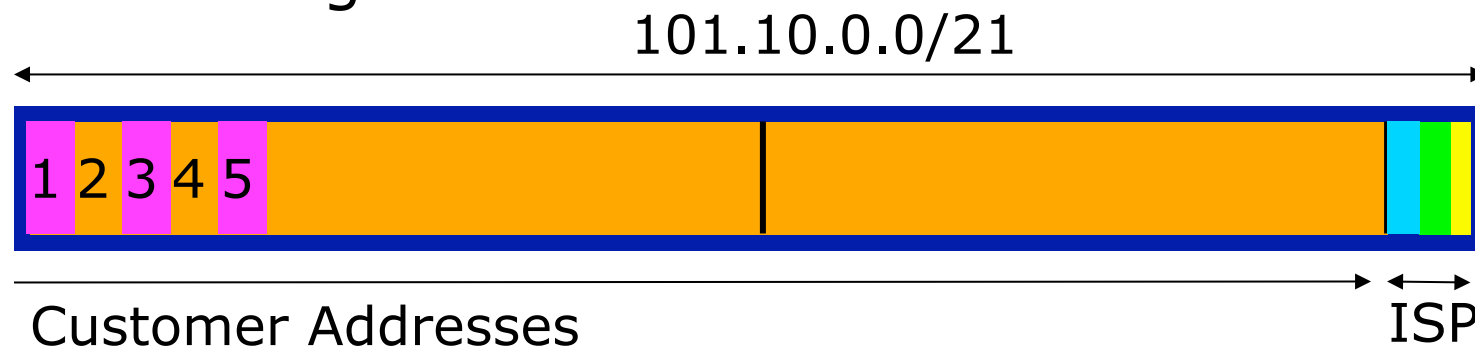
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- ❑ ISP Router loopbacks and backbone point to point links make up a small part of total address space
  - And they don't attract traffic, unlike customer address space
- ❑ Links from ISP Aggregation edge to customer router needs one /30
  - Small requirements compared with total address space
  - Some ISPs use IP unnumbered
- ❑ Planning customer assignments is a very important part of multihoming
  - Traffic engineering involves subdividing aggregate into pieces until load balancing works



# Unplanned IP addressing

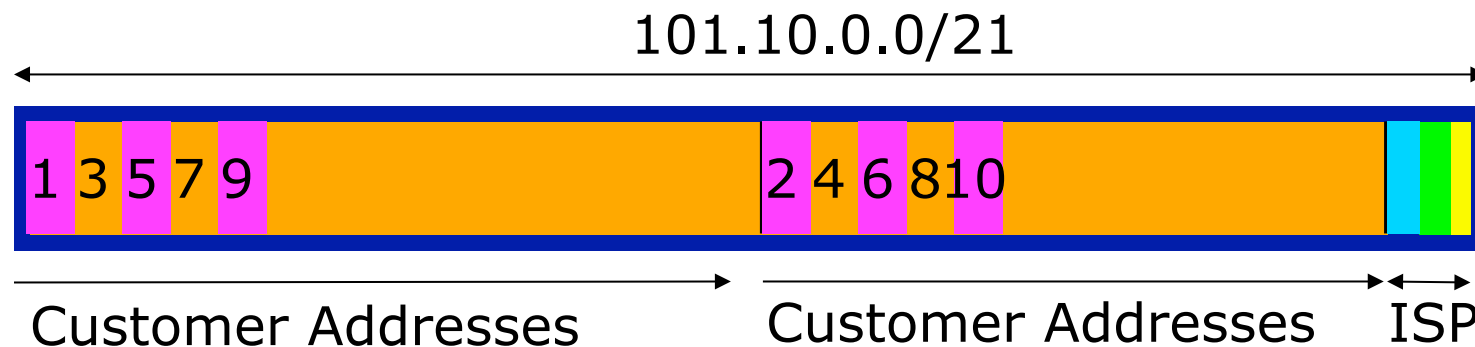
- ❑ ISP fills up customer IP addressing from one end of the range:



- ❑ Customers generate traffic
  - Dividing the range into two pieces will result in one /22 with all the customers, and one /22 with just the ISP infrastructure the addresses
  - No loadbalancing as all traffic will come in the first /22
  - Means further subdivision of the first /22 = harder work

# Planned IP addressing

- If ISP fills up customer addressing from both ends of the range:



- Scheme then is:
  - First customer from first /22, second customer from second /22, third from first /22, etc
- This works also for residential versus commercial customers:
  - Residential from first /22
  - Commercial from second /22

# Planned IP Addressing

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- ❑ This works fine for multihoming between two upstream links (same or different providers)
- ❑ Can also subdivide address space to suit more than two upstreams
  - Follow a similar scheme for populating each portion of the address space
- ❑ Don't forget to always announce an aggregate out of each link



# BGP Multihoming Techniques

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- ❑ Why Multihome?
- ❑ Definition & Options
- ❑ How to Multihome
- ❑ Principles & Addressing
- ❑ **Basic Multihoming**
- ❑ Service Provider Multihoming

# Basic Multihoming



Let's try some simple worked examples...

# Basic Multihoming

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- Will look at two cases:
  - Multihoming with the same ISP
  - Multihoming to different ISPs
- Will keep the examples easy
  - Understanding easy concepts will make the more complex scenarios easier to comprehend
  - All assume that the site multihoming has a /19 address block

# Basic Multihoming

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- This type is most commonplace at the edge of the Internet
  - Networks here are usually concerned with inbound traffic flows
  - Outbound traffic flows being “nearest exit” is usually sufficient
- Can apply to the leaf ISP as well as Enterprise networks

# Basic Multihoming



Multihoming to the Same ISP



# Basic Multihoming:

## Multihoming to the same ISP

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- Use BGP for this type of multihoming
  - use a private AS (ASN > 64511)
  - There is no need or justification for a public ASN
    - Making the nets of the end-site visible gives no useful information to the Internet
- Upstream ISP proxy aggregates
  - in other words, announces only your address block to the Internet from their AS (as would be done if you had one statically routed connection)

# Two links to the same ISP



One link primary, the other link  
backup only

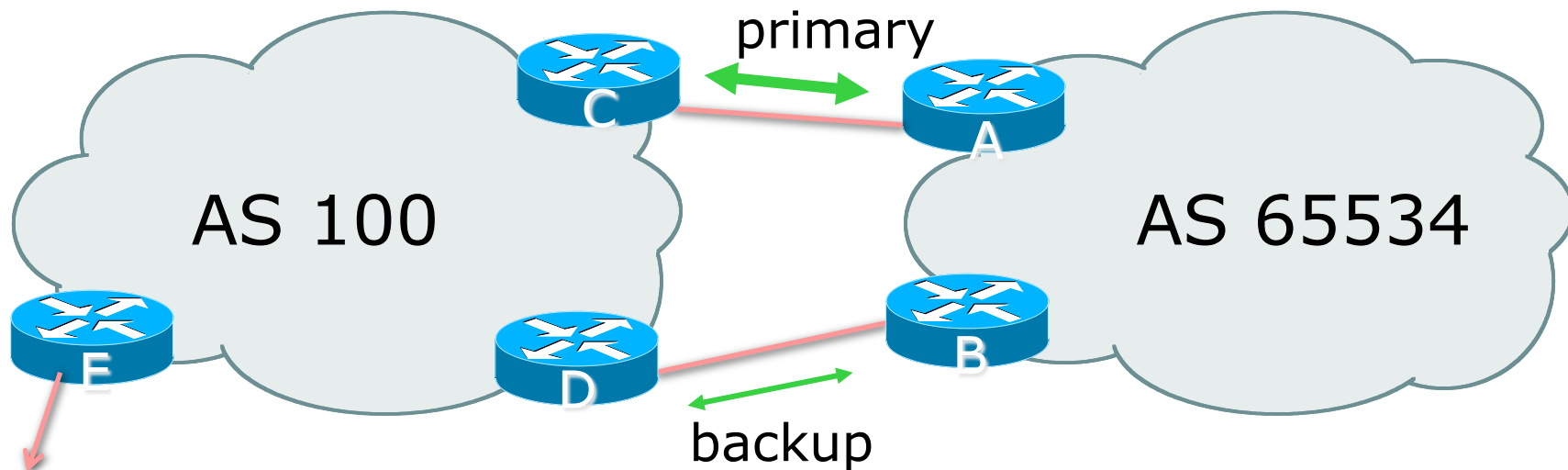
# Two links to the same ISP (one as backup only)

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- Applies when end-site has bought a large primary WAN link to their upstream a small secondary WAN link as the backup
  - For example, primary path might be an E1, backup might be 64kbps

# Two links to the same ISP (one as backup only)

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- ❑ AS100 removes private AS and any customer subprefixes from Internet announcement

# Two links to the same ISP (one as backup only)

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- ❑ Announce /19 aggregate on each link
  - primary link:
    - ❑ Outbound – announce /19 unaltered
    - ❑ Inbound – receive default route
  - backup link:
    - ❑ Outbound – announce /19 with increased metric
    - ❑ Inbound – received default, and reduce local preference
- ❑ When one link fails, the announcement of the /19 aggregate via the other link ensures continued connectivity

# Two links to the same ISP



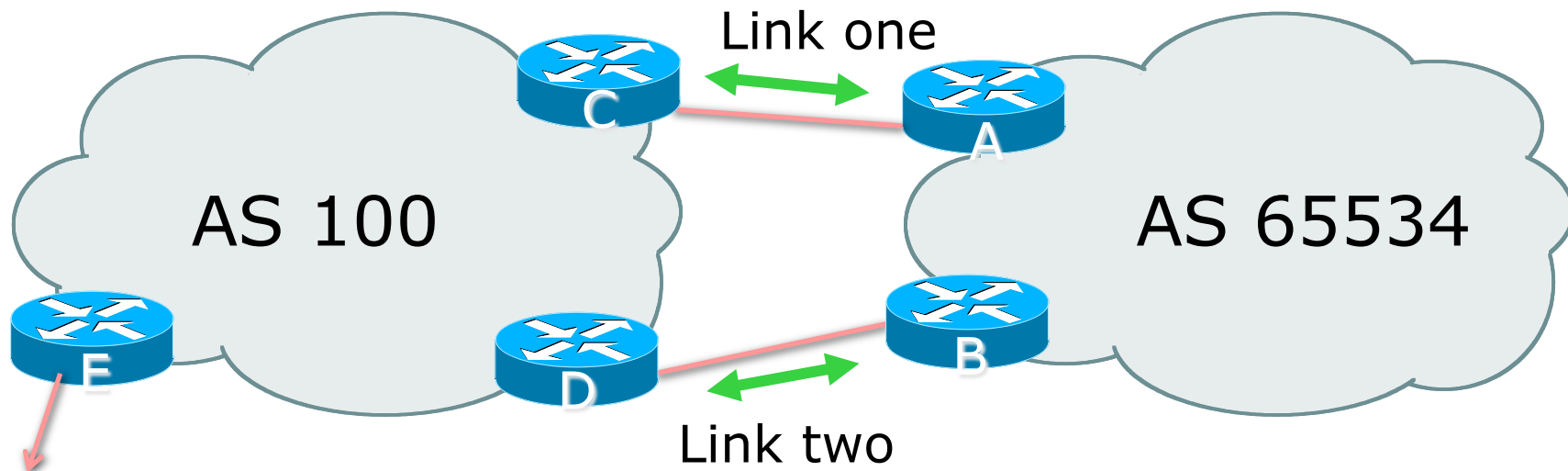
With Loadsharing

# Loadsharing to the same ISP

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- More common case
- End sites tend not to buy circuits and leave them idle, only used for backup as in previous example
- This example assumes equal capacity circuits
  - Unequal capacity circuits requires more refinement – see later

# Loadsharing to the same ISP



- ❑ Border router E in AS100 removes private AS and any customer subprefixes from Internet announcement



# Loadsharing to the same ISP

---

- ❑ Announce /19 aggregate on each link
- ❑ Split /19 and announce as two /20s, one on each link
  - basic inbound loadsharing
  - assumes equal circuit capacity and even spread of traffic across address block
- ❑ Vary the split until “perfect” loadsharing achieved
- ❑ Accept the default from upstream
  - basic outbound loadsharing by nearest exit
  - okay in first approx as most ISP and end-site traffic is inbound

# Loadsharing to the same ISP

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- ❑ Loadsharing configuration is only on customer router
- ❑ Upstream ISP has to
  - remove customer subprefixes from external announcements
  - remove private AS from external announcements
- ❑ Could also use BGP communities

# Two links to the same ISP



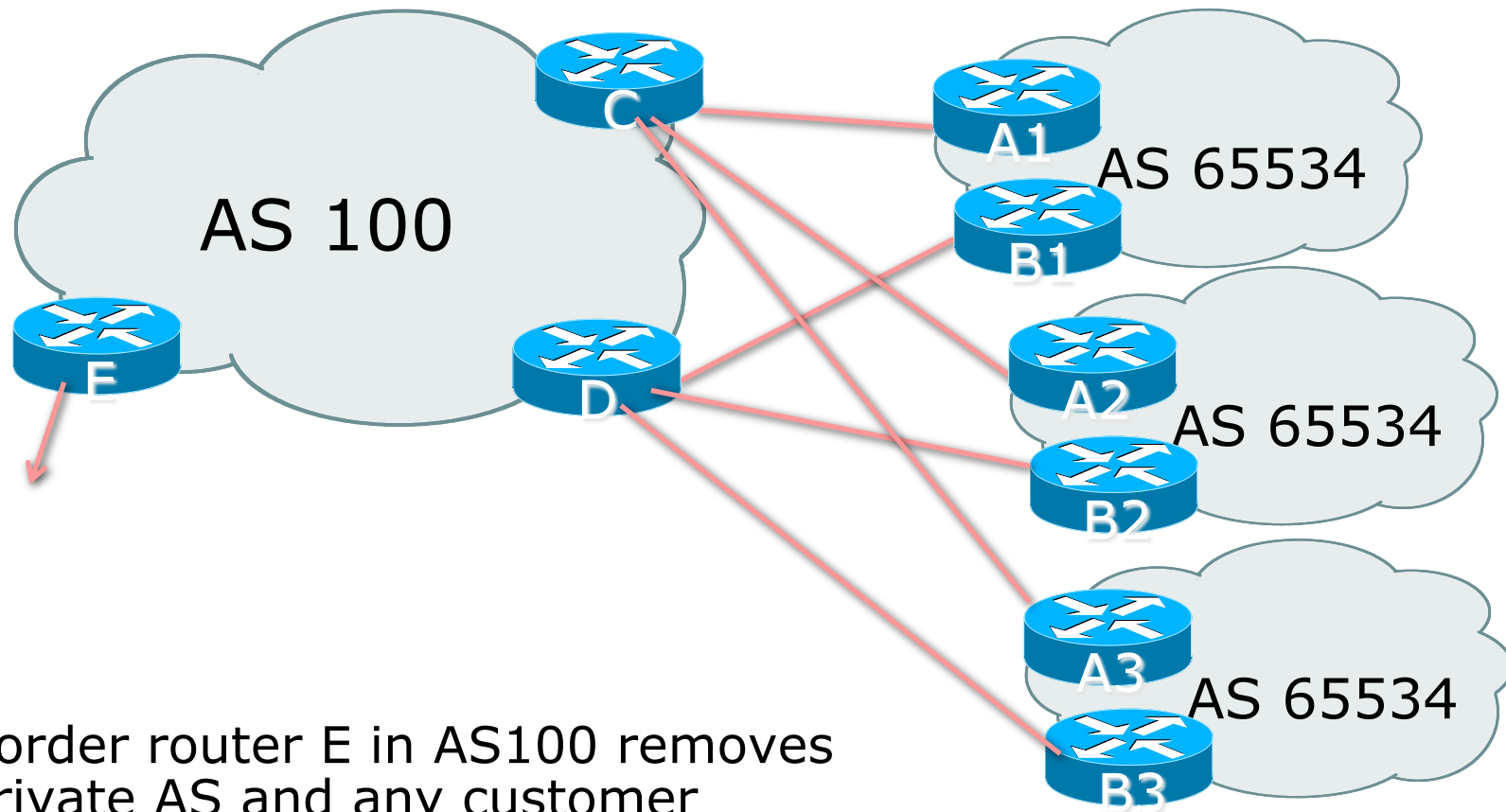
Multiple Dualhomed Customers  
(RFC2270)

# Multiple Dualhomed Customers (RFC2270)

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- ❑ Unusual for an ISP just to have one dualhomed customer
  - Valid/valuable service offering for an ISP with multiple PoPs
  - Better for ISP than having customer multihome with another provider!
- ❑ Look at scaling the configuration
  - ⇒ Simplifying the configuration
  - Using templates, peer-groups, etc
  - Every customer has the same configuration (basically)

# Multiple Dualhomed Customers (RFC2270)



- ❑ Border router E in AS100 removes private AS and any customer subprefixes from Internet announcement

# Multiple Dualhomed Customers (RFC2270)

---

- ❑ Customer announcements as per previous example
- ❑ Use the same private AS for each customer
  - documented in RFC2270
  - address space is not overlapping
  - each customer hears default only
- ❑ Each Router A and B has the same configuration for each instance

# Multihoming Summary

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- ❑ Use private AS for multihoming to the same upstream
- ❑ Leak subprefixes to upstream only to aid loadsharing
- ❑ Upstream router E configuration is identical across all situations

# Basic Multihoming



Multihoming to different ISPs



# Two links to different ISPs

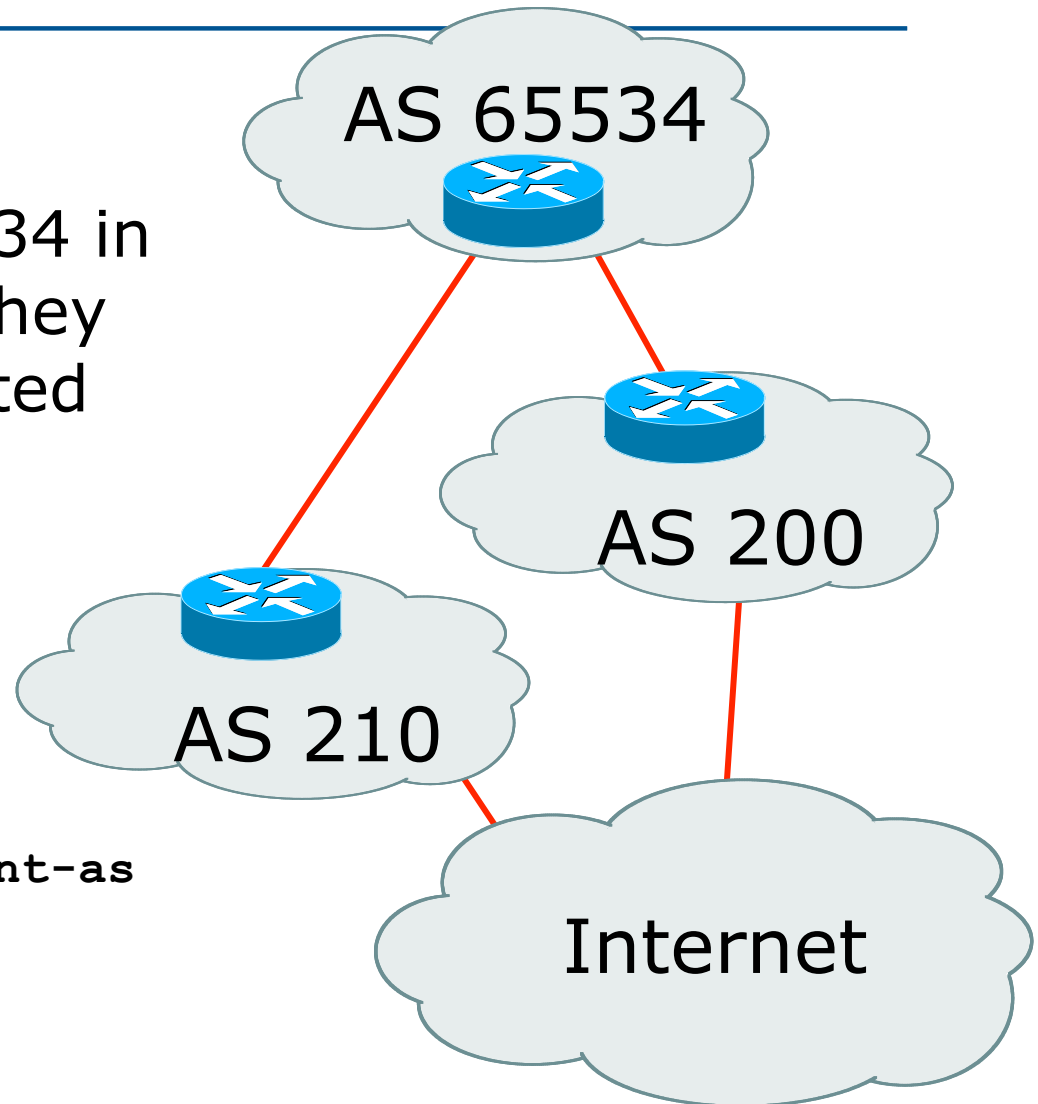
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- Use a Public AS
  - Or use private AS if agreed with the other ISP
  - But some people don't like the "inconsistent-AS" which results from use of a private-AS
- Address space comes from
  - both upstreams or
  - Regional Internet Registry
- Configuration concepts very similar

# Inconsistent-AS?

- ❑ Viewing the prefixes originated by AS65534 in the Internet shows they appear to be originated by both AS210 and AS200
  - This is NOT bad
  - Nor is it illegal
- ❑ IOS command is

```
show ip bgp inconsistent-as
```



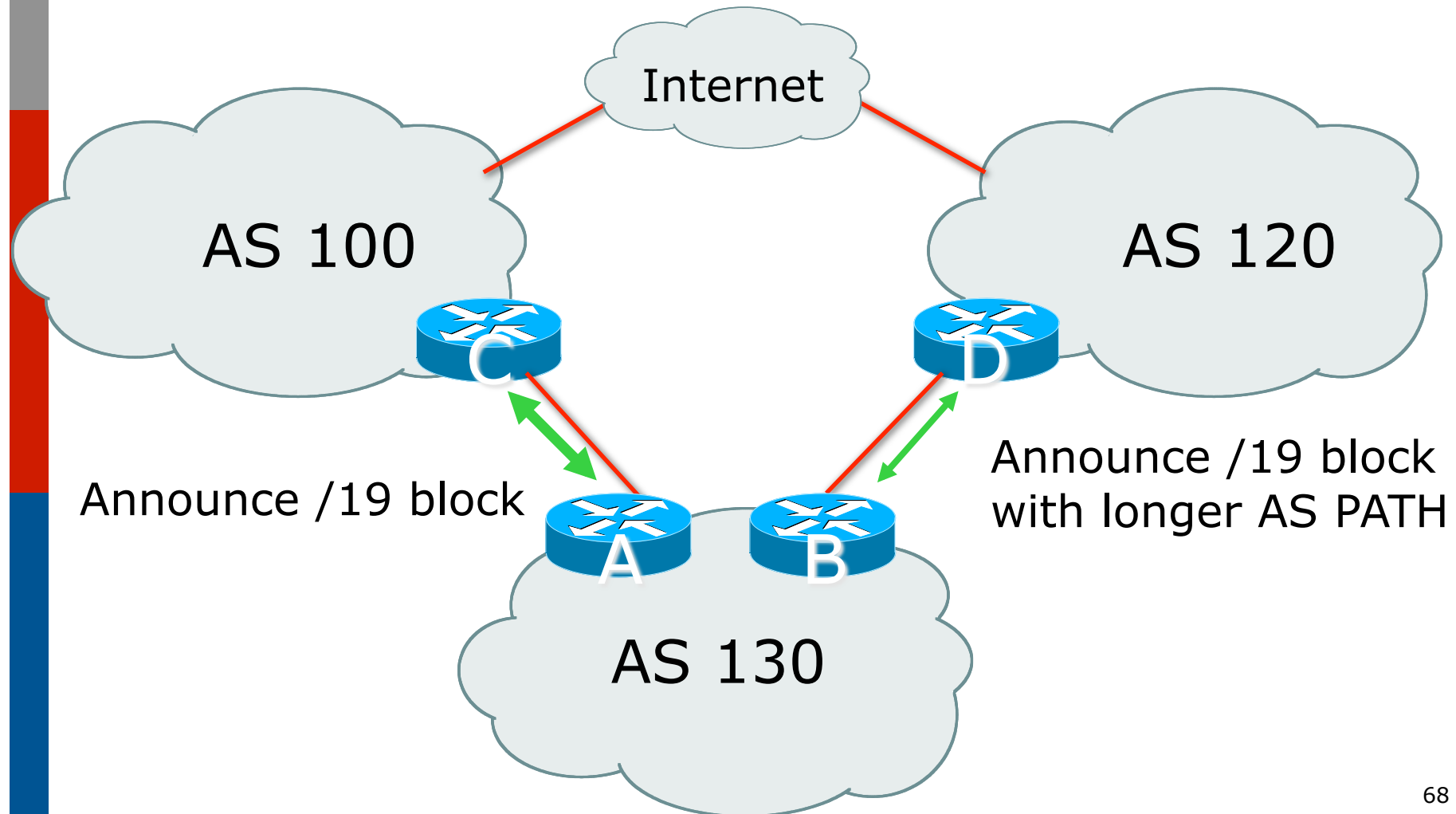
# Two links to different ISPs



One link primary, the other link  
backup only

# Two links to different ISPs (one as backup only)

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# Two links to different ISPs (one as backup only)

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- ❑ Announce /19 aggregate on each link
  - primary link makes standard announcement
  - backup link lengthens the AS PATH by using AS PATH prepend
- ❑ When one link fails, the announcement of the /19 aggregate via the other link ensures continued connectivity

# Two links to different ISPs (one as backup only)

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- ❑ Not a common situation as most sites tend to prefer using whatever capacity they have
  - (Useful when two competing ISPs agree to provide mutual backup to each other)
- ❑ But it shows the basic concepts of using local-prefs and AS-path prepends for engineering traffic in the chosen direction

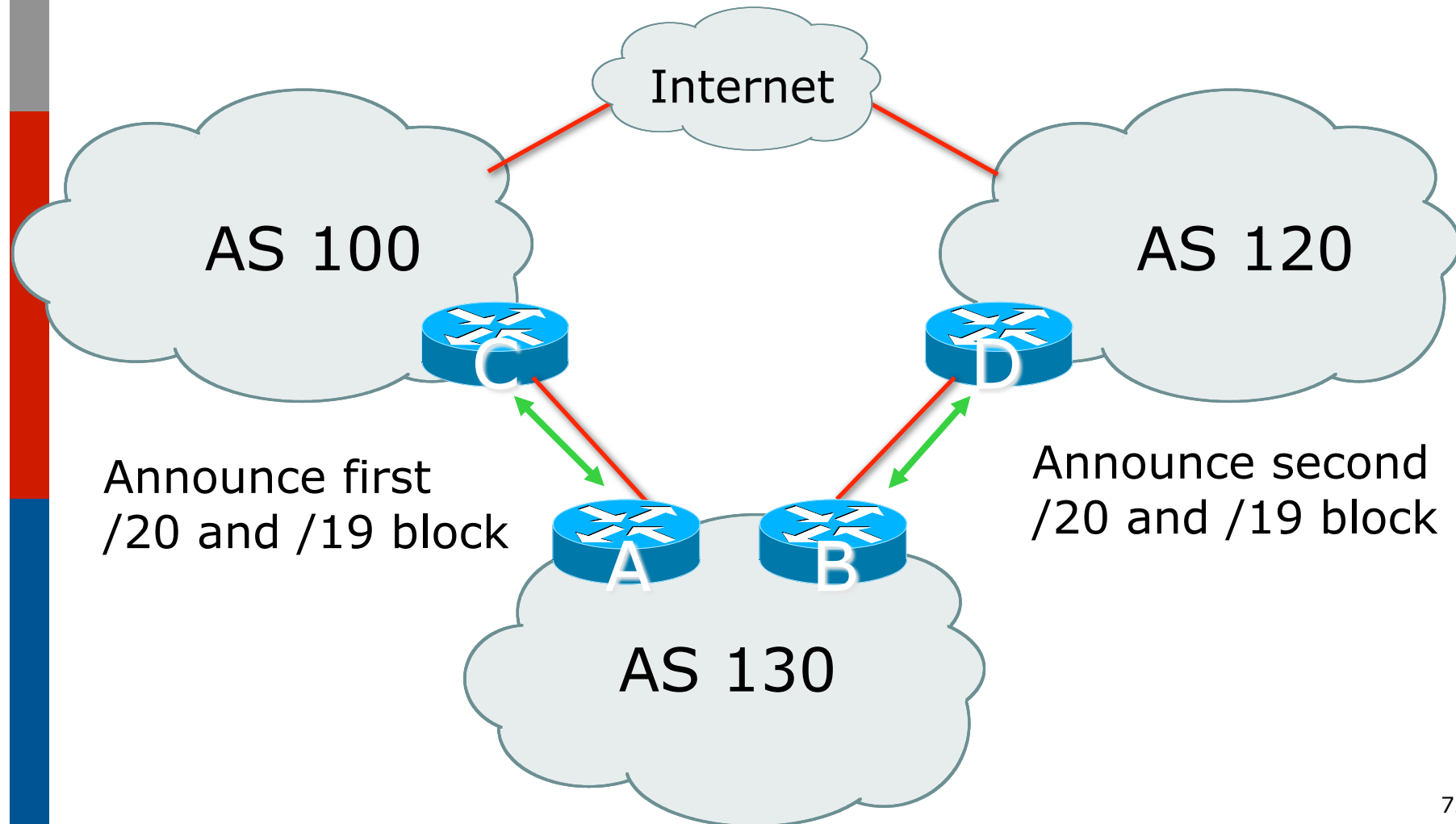
# Two links to different ISPs



With Loadsharing

# Two links to different ISPs (with loadsharing)

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# Two links to different ISPs (with loadsharing)

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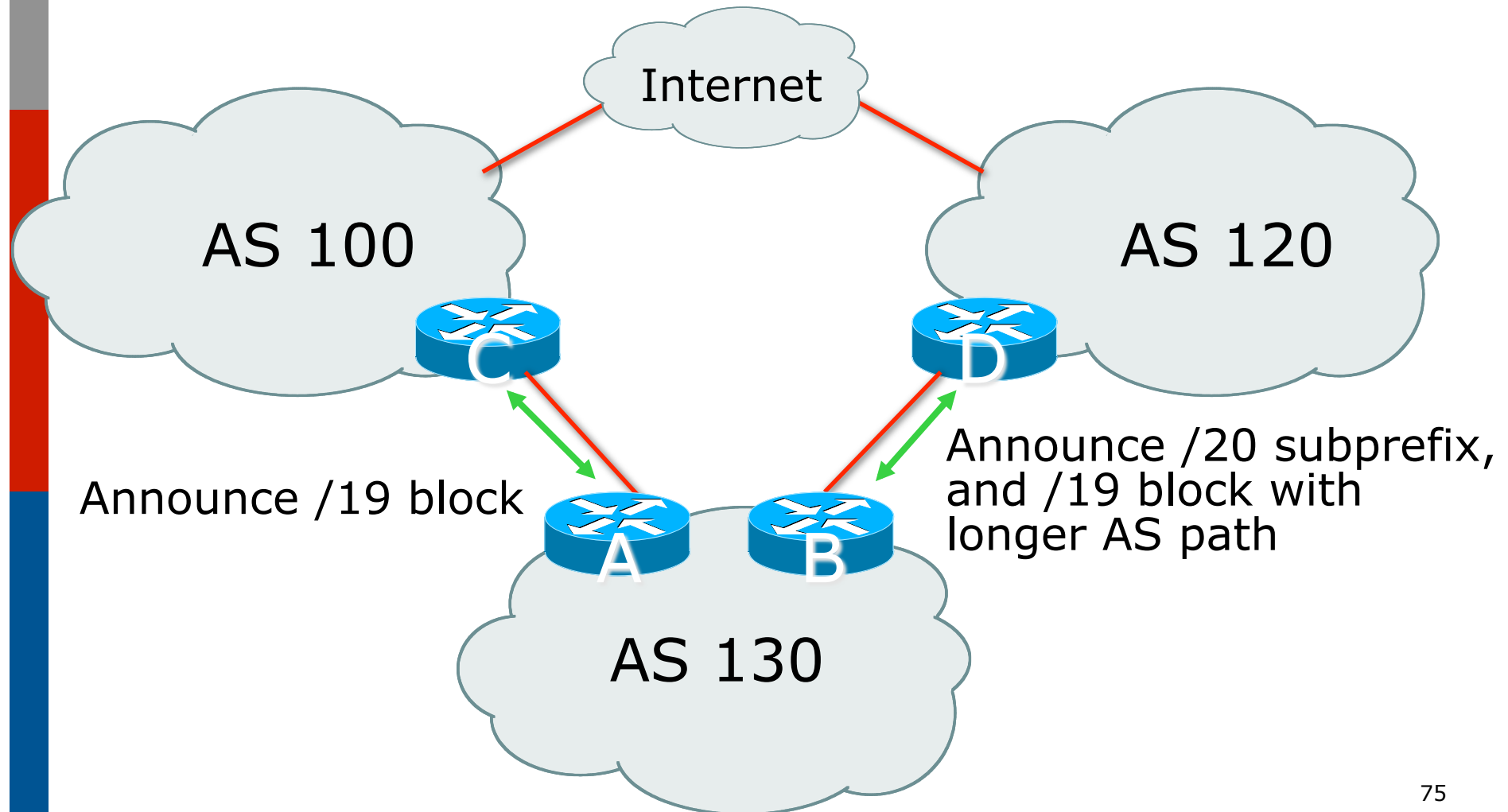
- ❑ Announce /19 aggregate on each link
- ❑ Split /19 and announce as two /20s, one on each link
  - basic inbound loadsharing
- ❑ When one link fails, the announcement of the /19 aggregate via the other ISP ensures continued connectivity

# Two links to different ISPs



More Controlled Loadsharing

# Loadsharing with different ISPs



# Loadsharing with different ISPs

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- Announce /19 aggregate on each link
  - On first link, announce /19 as normal
  - On second link, announce /19 with longer AS PATH, and announce one /20 subprefix
    - controls loadsharing between upstreams and the Internet
- Vary the subprefix size and AS PATH length until “perfect” loadsharing achieved
- Still require redundancy!

# Loadsharing with different ISPs

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- ❑ This example is more commonplace
- ❑ Shows how ISPs and end-sites subdivide address space frugally, as well as use the AS-PATH prepend concept to optimise the load sharing between different ISPs
- ❑ Notice that the /19 aggregate block is **ALWAYS** announced



# BGP Multihoming Techniques

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- ❑ Why Multihome?
- ❑ Definition & Options
- ❑ How to Multihome
- ❑ Principles & Addressing
- ❑ Basic Multihoming
- ❑ “BGP Traffic Engineering”

# Service Provider Multihoming



BGP Traffic Engineering

# Service Provider Multihoming

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- ❑ Previous examples dealt with loadsharing inbound traffic
  - Of primary concern at Internet edge
  - What about outbound traffic?
- ❑ Transit ISPs strive to balance traffic flows in both directions
  - Balance link utilisation
  - Try and keep most traffic flows symmetric
  - Some edge ISPs try and do this too
- ❑ The original “Traffic Engineering”



# Service Provider Multihoming

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- Balancing outbound traffic requires inbound routing information
  - Common solution is “full routing table”
  - Rarely necessary
    - Why use the “routing mallet” to try solve loadsharing problems?
  - “Keep It Simple” is often easier (and \$\$\$ cheaper) than carrying N-copies of the full routing table

# Service Provider Multihoming

## MYTHS!!

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### Common MYTHS

1. **You need the full routing table to multihome**
  - People who sell router memory would like you to believe this
  - Only true if you are a transit provider
  - Full routing table can be a significant hindrance to multihoming
2. **You need a BIG router to multihome**
  - Router size is related to data rates, not running BGP
  - In reality, to multihome, your router needs to:
    - ▣ Have two interfaces,
    - ▣ Be able to talk BGP to at least two peers,
    - ▣ Be able to handle BGP attributes,
    - ▣ Handle at least one prefix
3. **BGP is complex**
  - In the wrong hands, yes it can be! Keep it Simple!

# Service Provider Multihoming: Some Strategies

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- ❑ Take the prefixes you need to aid traffic engineering
  - Look at NetFlow data for popular sites
- ❑ Prefixes originated by your immediate neighbours and their neighbours will do more to aid load balancing than prefixes from ASNs many hops away
  - Concentrate on local destinations
- ❑ Use default routing as much as possible
  - Or use the full routing table with care

# Service Provider Multihoming

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- Two cases:
  - One upstream, one local peer
  - Two upstreams, one local peer
- Require BGP and a public ASN
- Examples assume that the local network has their own /19 address block

# Service Provider Multihoming



One upstream, one local peer

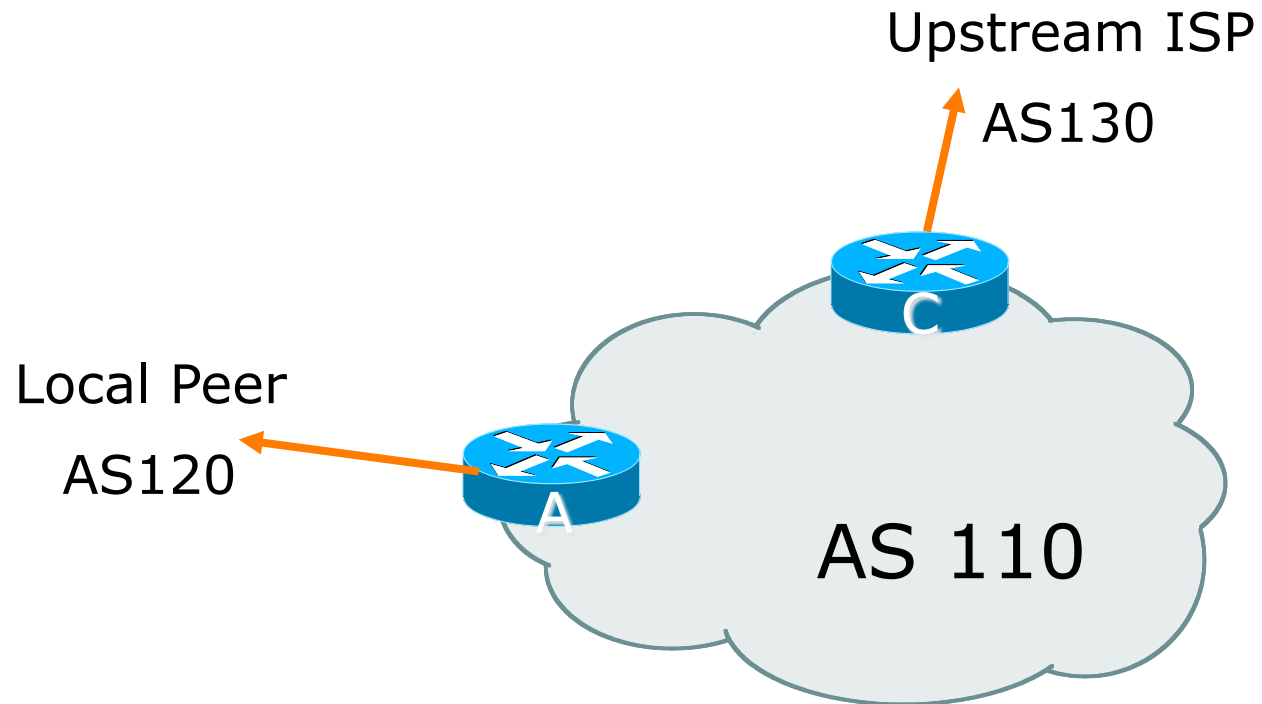
# One Upstream, One Local Peer

---

- ❑ Very common situation in many regions of the Internet
- ❑ Connect to upstream transit provider to see the “Internet”
- ❑ Connect to the local competition so that local traffic stays local
  - Saves spending valuable \$ on upstream transit costs for local traffic

# One Upstream, One Local Peer

---



# One Upstream, One Local Peer

---

- ❑ Announce /19 aggregate on each link
- ❑ Accept default route only from upstream
  - Either 0.0.0.0/0 or a network which can be used as default
- ❑ Accept all routes from local peer
- ❑ This example is easily extendable for multiple local peers and/or an Internet Exchange Point





# Aside:

## Configuration Recommendations

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### □ Private Peers

- The peering ISPs exchange prefixes they originate
- Sometimes they exchange prefixes from neighbouring ASNs too

### □ Be aware that the private peer eBGP router should carry only the prefixes you want the private peer to receive

- Otherwise they could point a default route to you and unintentionally transit your backbone

# Service Provider Multihoming



Two Upstreams, One local peer

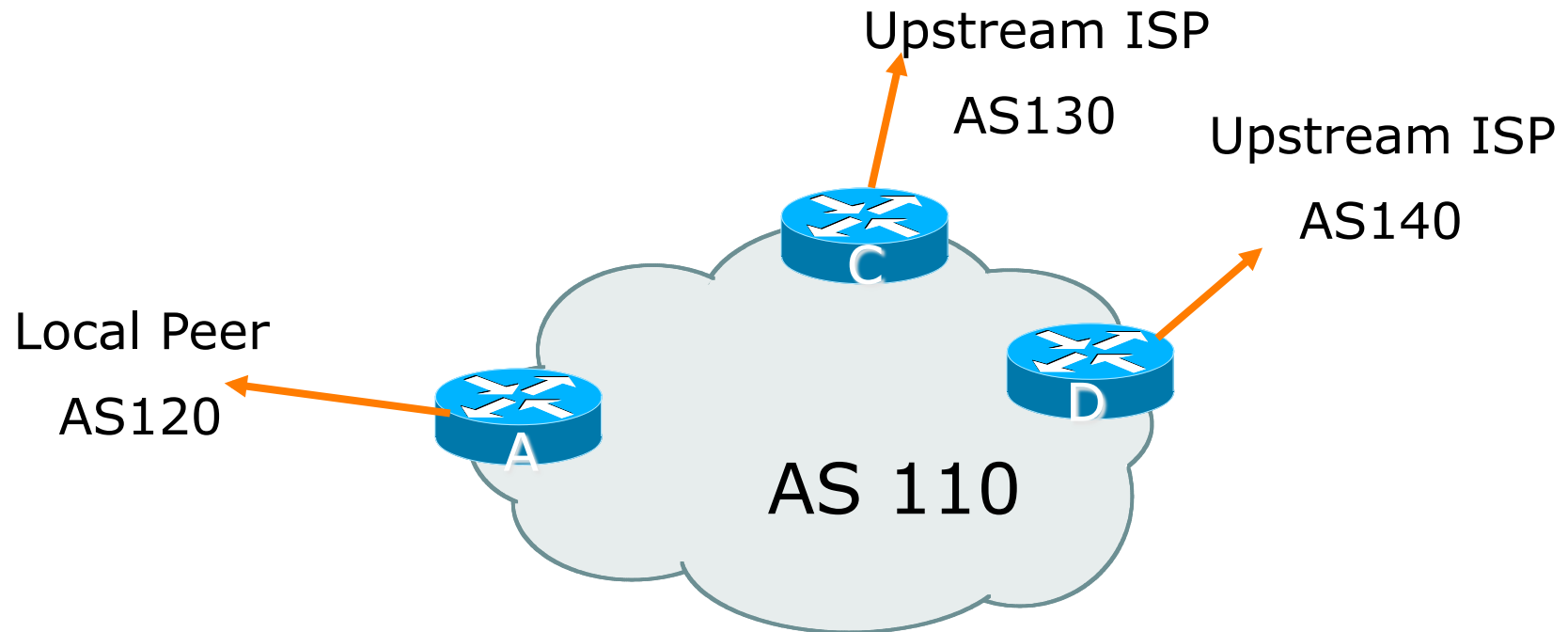
# Two Upstreams, One Local Peer

---

- Connect to both upstream transit providers to see the “Internet”
  - Provides external redundancy and diversity – the reason to multihome
- Connect to the local peer so that local traffic stays local
  - Saves spending valuable \$ on upstream transit costs for local traffic

# Two Upstreams, One Local Peer

---



# Two Upstreams, One Local Peer

---

- ❑ Announce /19 aggregate on each link
- ❑ Accept default route only from upstreams
  - Either 0.0.0.0/0 or a network which can be used as default
- ❑ Accept all routes from local peer
- ❑ Note separation of Router C and D
  - Single edge router means no redundancy
- ❑ Router A
  - Same routing configuration as in example with one upstream and one local peer

# Two Upstreams, One Local Peer

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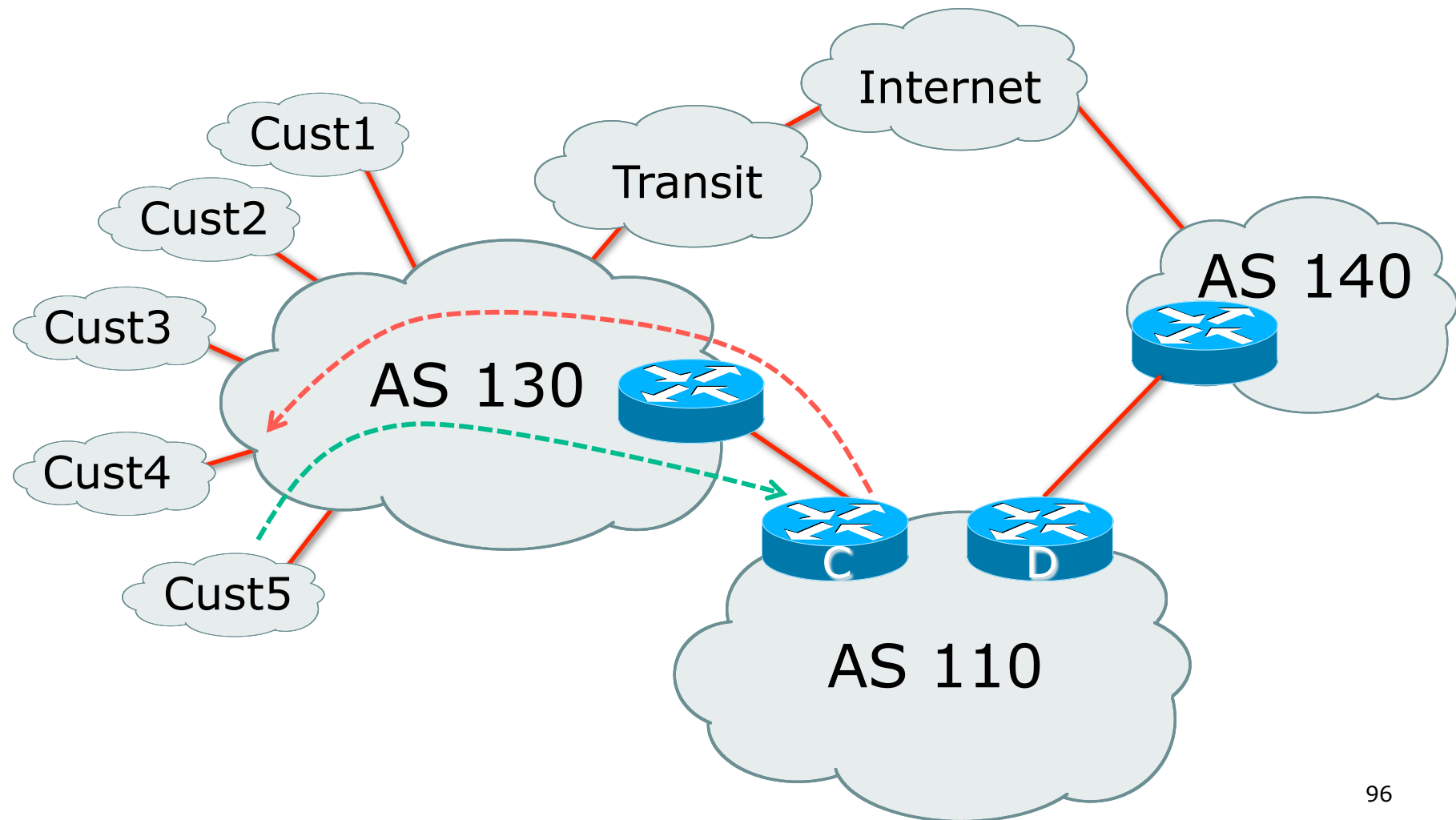
- Traffic out to the two upstreams will take nearest exit
  - Inexpensive routers required
  - This is not useful in practice especially for international links
  - Loadsharing needs to be better

# Two Upstreams, One Local Peer

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- Better configuration options:
  - Accept full routing from both upstreams
    - **Expensive & unnecessary!**
  - Accept default from one upstream and some routes from the other upstream
    - **The way to go!**

# Loadsharing with different ISPs





# Two Upstreams, One Local Peer

## Full Routes

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- Full routes from upstreams
  - Summary of routes received:

ASN	Full Routes	Partial Routes
AS140	430000 @ Ip100	
AS130	30000 @ Ip 120 400000 @ Ip 80	
Total	860000	

# Two Upstreams, One Local Peer

## Full Routes

---

- Full routes from upstreams
  - Expensive – needs lots of memory and CPU
  - Need to play preference games
  - Previous example is only an example – real life will need improved fine-tuning!
  - Previous example doesn't consider inbound traffic – see earlier in presentation for examples

# Two Upstreams, One Local Peer

## Partial Routes: Strategy

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- ❑ Ask one upstream for a default route
  - Easy to originate default towards a BGP neighbour
- ❑ Ask other upstream for a full routing table
  - Then filter this routing table based on neighbouring ASN
  - E.g. want traffic to their neighbours to go over the link to that ASN
  - Most of what upstream sends is thrown away
  - Easier than asking the upstream to set up custom BGP filters for you

# Two Upstreams, One Local Peer

## Partial Routes

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- ❑ Router C configuration:
  - Accept full routes from AS130
    - ❑ (or get them to send less)
  - Filter ASNs so only AS130 and AS130's neighbouring ASes are accepted
  - Allow default, and set it to local preference 80
  - Traffic to those ASes will go over AS130 link
  - Traffic to other all other ASes will go over the link to AS140
  - If AS140 link fails, backup via AS130 – and vice-versa

# Two Upstreams, One Local Peer

## Partial Routes

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- Partial routes from upstreams
  - Summary of routes received:

ASN	Full Routes	Partial Routes
AS140	430000 @ lp100	1 @ lp 100
AS130	30000 @ lp 120 400000 @ lp 80	30000 @ lp 100 1 @ lp 80
Total	860000	30002

# Two Upstreams, One Local Peer

## Partial Routes

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- Partial routes from upstreams
  - Not expensive – only carry the routes necessary for loadsharing
  - Need to filter on AS paths
  - Previous example is only an example – real life will need improved fine-tuning!
  - Previous example doesn't consider inbound traffic – see earlier in presentation for examples



# Aside:

## Configuration Recommendation

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- ❑ When distributing internal default by iBGP or OSPF/ISIS
  - Make sure that routers connecting to private peers or to IXPs do NOT carry the default route
  - Otherwise they could point a default route to you and unintentionally transit your backbone
  - Simple fix for Private Peer/IXP routers:

```
ip route 0.0.0.0 0.0.0.0 null0
```

# Summary





# Summary

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- ❑ Multihoming is not hard, really...
  - **Keep It Simple & Stupid!**
- ❑ Full routing table is *rarely* required
  - A default is often just as good
  - If customers want 430k prefixes, charge them money for the privilege

# BGP Multihoming Techniques



End of Tutorial