## Routing Basics

### ISP Workshops



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## Acknowledgements

- This material originated from the Cisco ISP/IXP Workshop Programme developed by Philip Smith & Barry Greene
- Use of these materials is encouraged as long as the source is fully acknowledged and this notice remains in place
- Bug fixes and improvements are welcomed
  - Please email workshop (at) bgp4all.com

## Routing Concepts

- Routers
- Routing
- Forwarding
- Some definitions
- □ Policy options
- Routing Protocols

### What is a Router?

- A router is a layer 3 device
- Used for interconnecting networks at layer 3
- A router generally has at least two interfaces
  - With VLANs a router can have only one interface (known as "router on a stick")
- A router looks at the destination address in the IP packet, and decides how to forward it



## The Routing Table

- Each router/host has a *routing table*, indicating the path or the next hop for a given destination host or a network
- The router/host tries to match the destination address of a packet against entries in the routing table
- If there is a match, the router forwards it to the corresponding gateway router or directly to the destination host
- Default route is taken if no other entry matches the destination address

## The Routing Table

Destination	Next-Hop	Interface
10.40.0.0/16	192.248.40.60	Ethernet0
192.248.0.140/30	Directly connected	Serial1
192.248.40.0/26	Directly connected	Ethernet0
192.248.0.0/17	192.248.0.141	Serial1
203.94.73.202/32	192.248.40.3	Ethernet0
203.115.6.132/30	Directly connected	Serial0
Default	203.115.6.133	Serial0

Typical routing table on a simple edge router

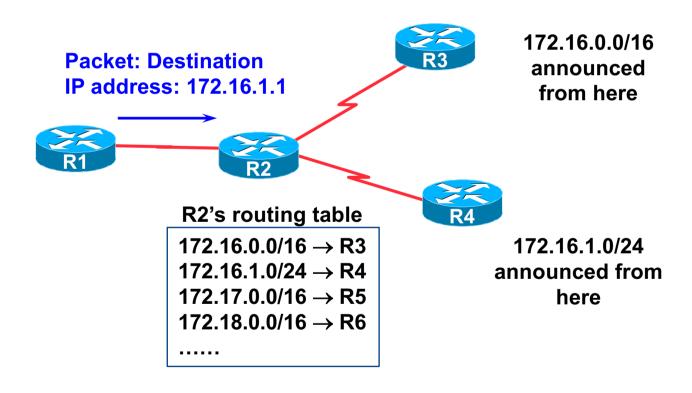
## IP Routing – finding the path

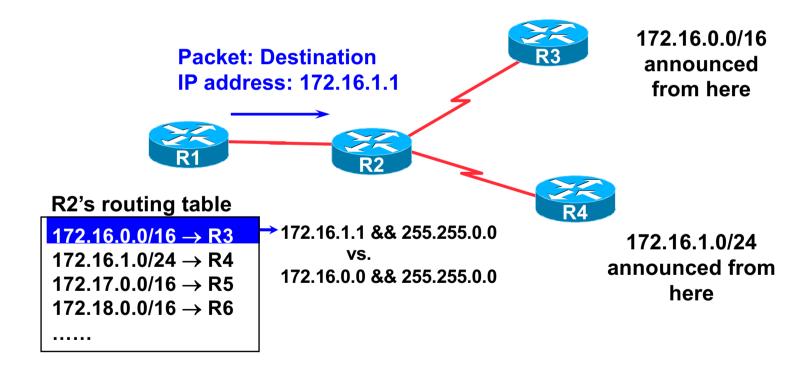
- Routing table entry (the path) is created by the administrator (static) or received from a routing protocol (dynamic)
- More than one routing protocol may run on a router
  - Each routing protocol builds its own routing table (Local RIB)
- Several alternative paths may exist
  - Best path selected for the router's Global routing table (RIB)
- Decisions are updated periodically or as topology changes (event driven)
- Decisions are based on:
  - Topology, policies and metrics (hop count, filtering, delay, bandwidth, etc.)

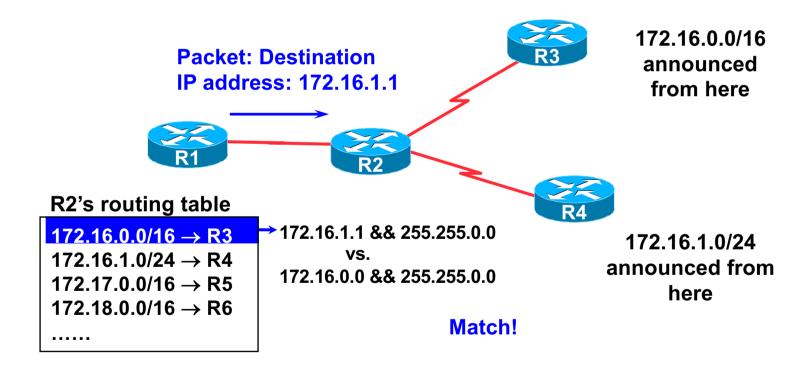
## IP route lookup

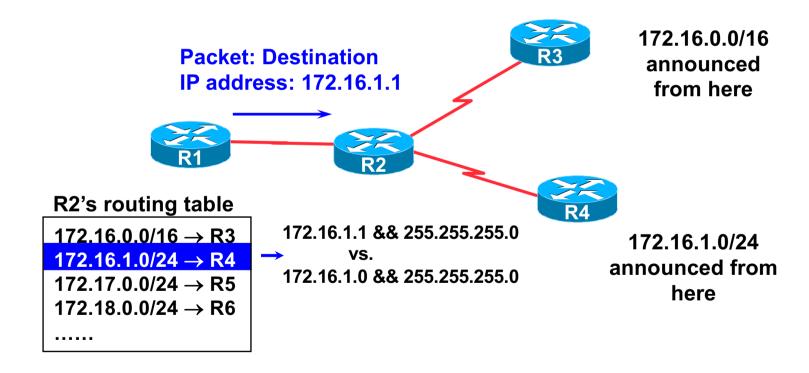
- Based on destination IP address
- "longest match" routing
  - More specific prefix preferred over less specific prefix
  - **Example:** packet with destination of
  - 172.16.1.1/32 is sent to the router announcing 172.16.1.0/24 rather than the router announcing 172.16.0.0/16.

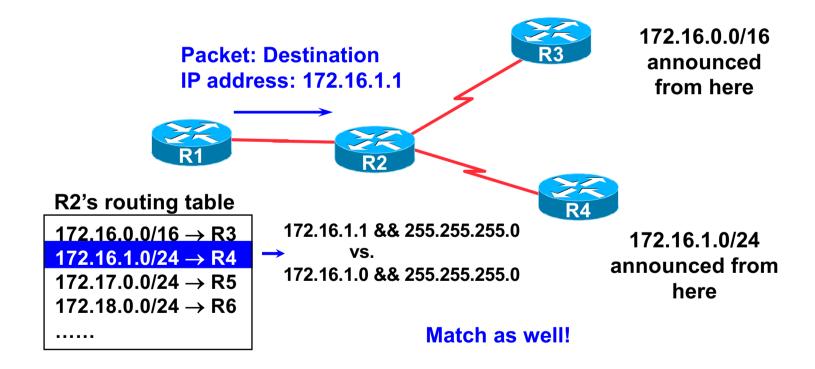
## IP route lookup

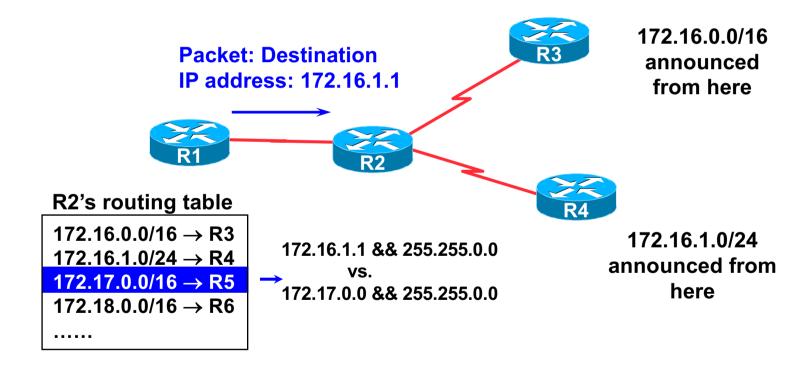


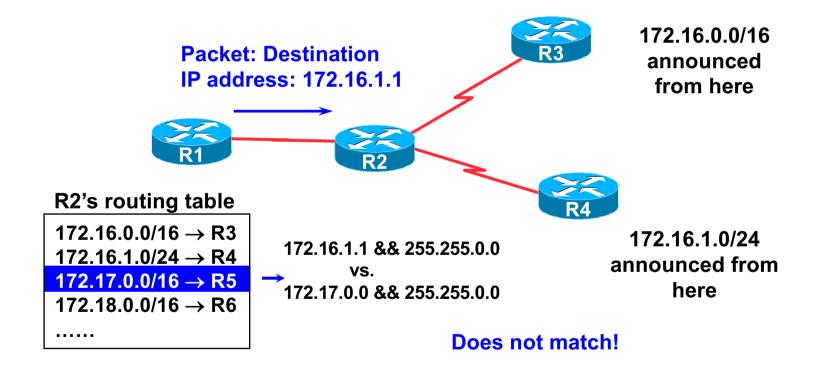


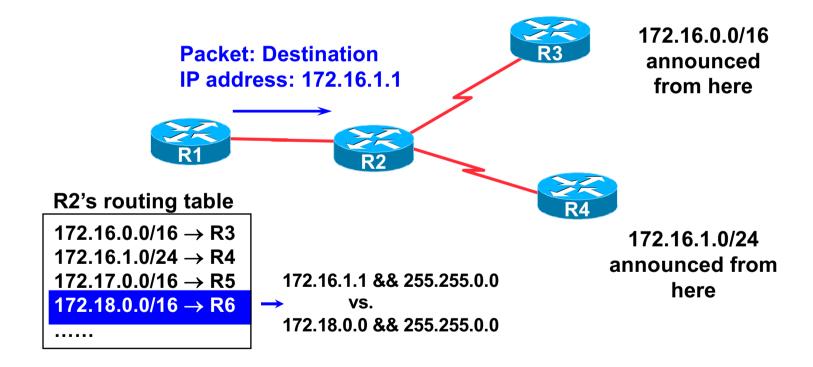


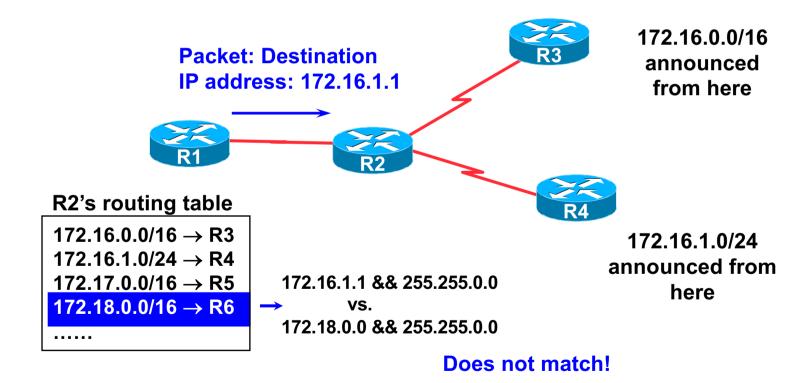


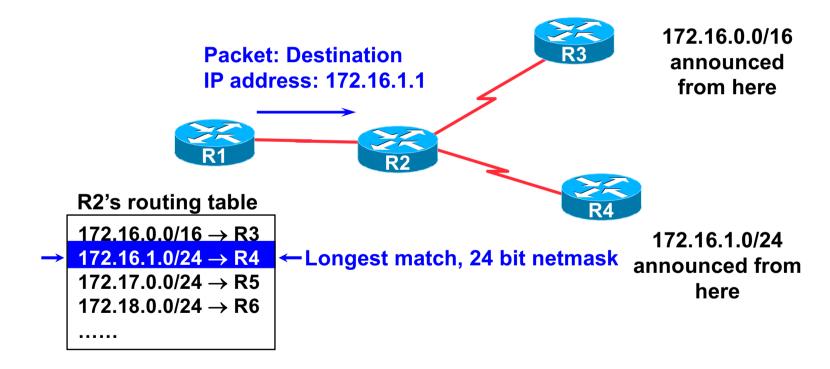


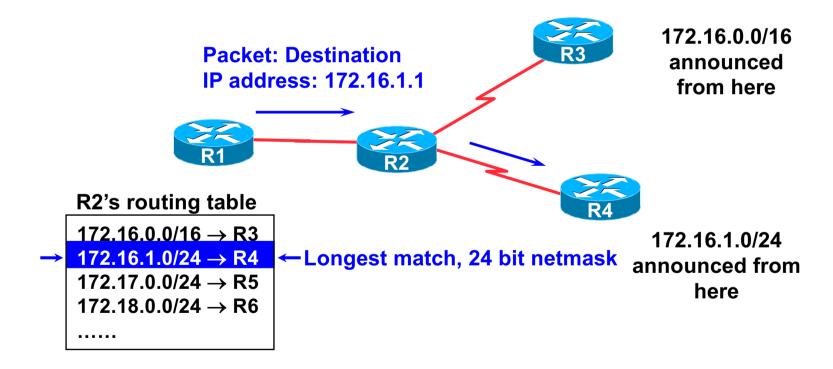












## Routing versus Forwarding

- Routing = building maps and giving directions
- Forwarding = moving packets between interfaces according to the "directions"

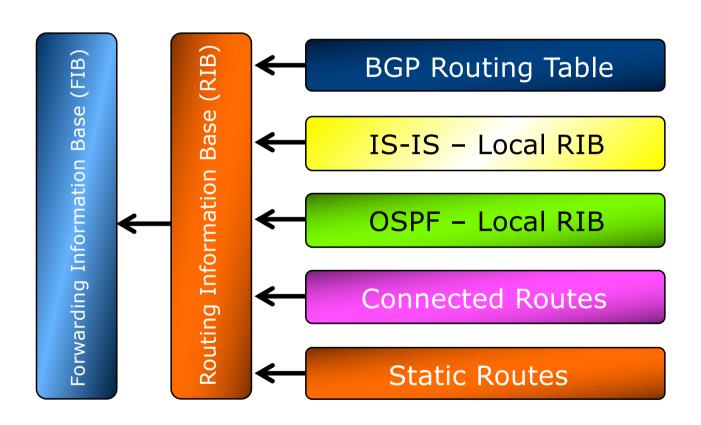




## IP Forwarding

- Router decides which interface a packet is sent to
- Forwarding table populated by routing process
- Forwarding decisions:
  - destination address
  - class of service (fair queuing, precedence, others)
  - local requirements (packet filtering)
- Forwarding is usually aided by special hardware

## Routing Tables Feed the Forwarding Table



### The FIB

- □ FIB is the Forwarding Table
  - It contains destinations, the interfaces and the next-hops to get to those destinations
  - It is built from the router's Global RIB
  - Used by the router to figure out where to send the packet
  - Cisco IOS: "show ip cef"

### The Global RIB

- The Global RIB is the Routing Table
  - Built from the routing tables/RIBs of the routing protocols and static routes on the router
    - Routing protocol priority varies per vendor see addendum
  - It contains all the known destinations and the next-hops used to get to those destinations
  - One destination can have lots of possible next-hops only the best next-hop goes into the Global RIB
  - The Global RIB is used to build the FIB
  - Cisco IOS: "show ip route"

## Explicit versus Default Routing

- Default:
  - Simple, cheap (CPU, memory, bandwidth)
  - No overhead
  - Low granularity (metric games)
- Explicit: (default free zone)
  - Complex, expensive (CPU, memory, bandwidth)
  - High overhead
  - High granularity (every destination known)
- Hybrid:
  - Minimise overhead
  - Provide useful granularity
  - Requires some filtering knowledge

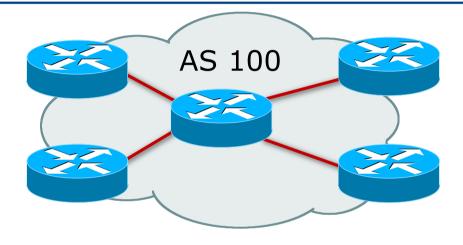
## Egress Traffic

- How packets leave your network
- Egress traffic depends on:
  - Route availability (what others send you)
  - Route acceptance (what you accept from others)
  - Policy and tuning (what you do with routes from others)
  - Peering and transit agreements

## Ingress Traffic

- How packets get to your network and your customers' networks
- Ingress traffic depends on:
  - What information you send and to whom
  - Based on your addressing and AS's
  - Based on others' policy (what they accept from you and what they do with it)

## Autonomous System (AS)



- Collection of networks with same routing policy
- Single routing protocol
- Usually under single ownership, trust and administrative control

### Definition of terms

#### Neighbours

- AS's which directly exchange routing information
- Routers which exchange routing information

#### Announce

send routing information to a neighbour

#### Accept

receive and use routing information sent by a neighbour

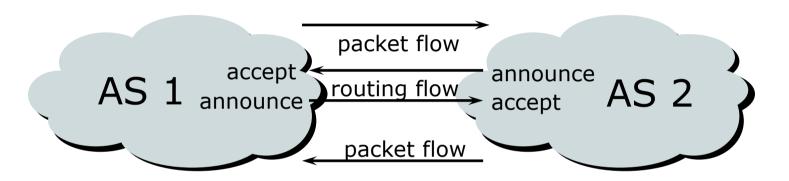
#### Originate

insert routing information into external announcements (usually as a result of the IGP)

#### Peers

 routers in neighbouring AS's or within one AS which exchange routing and policy information

## Routing flow and packet flow



For networks in AS1 and AS2 to communicate:

AS1 must announce to AS2

AS2 must accept from AS1

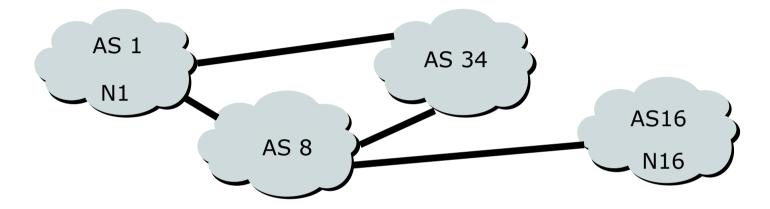
AS2 must announce to AS1

AS1 must accept from AS2

## Routing flow and Traffic flow

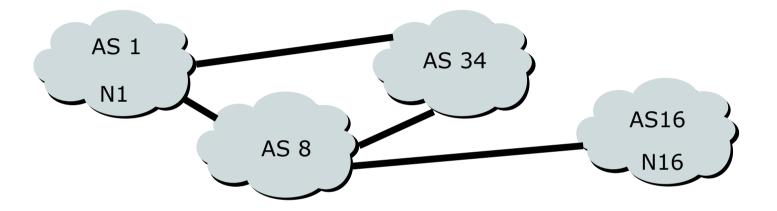
- Traffic flow is always in the opposite direction of the flow of Routing information
  - Filtering outgoing routing information inhibits traffic flow inbound
  - Filtering inbound routing information inhibits traffic flow outbound

# Routing Flow/Packet Flow: With multiple ASes



- □ For net N1 in AS1 to send traffic to net N16 in AS16:
  - AS16 must originate and announce N16 to AS8.
  - AS8 must accept N16 from AS16.
  - AS8 must announce N16 to AS1 or AS34.
  - AS1 must accept N16 from AS8 or AS34.
- For two-way packet flow, similar policies must exist for N1

# Routing Flow/Packet Flow: With multiple ASes

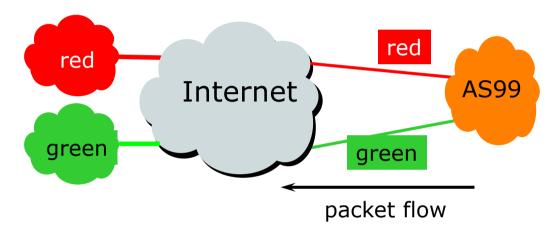


■ As more and more paths are implemented between sites it is easy to see how policies can become quite complex.

## Routing Policy

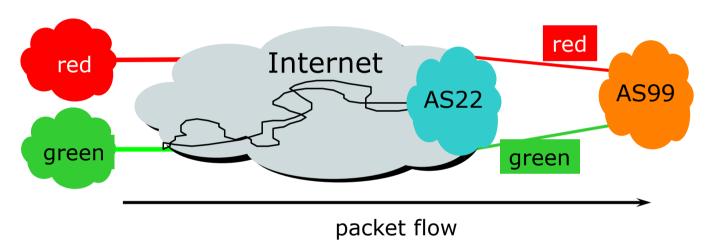
- Used to control traffic flow in and out of a network
- The network operator makes decisions on what routing information to accept and discard from its neighbours
  - Individual routes
  - Routes originated by specific ASes
  - Routes traversing specific ASes
  - Routes belonging to other groupings
    - Groupings which you define as you see fit

## Routing Policy Limitations



- AS99 uses red link for traffic to the red AS and the green link for remaining traffic
- To implement this policy, AS99 has to:
  - Accept routes originating from the red AS on the red link
  - Accept all other routes on the green link

## Routing Policy Limitations



- AS99 would like packets coming from the green AS to use the green link.
- But unless AS22 cooperates in pushing traffic from the green AS down the green link, there is very little that AS99 can do to achieve this aim

### Routing Policy Issues

- □ July 2019:
  - 67000 IPv6 prefixes & 763000 IPv4 prefixes
    - Not realistic to set policy on all of them individually
  - 65000 origin AS's
    - Too many to try and create individual policies for
- Routes tied to a specific AS or path may be unstable regardless of connectivity
- Solution: Groups of AS's are a natural abstraction for filtering purposes

## Routing Protocols

We now know what routing means...
...but what do the routers get up to?
And why are we doing this anyway?

- Internet is made up of the Network Operators who connect to each other's networks
- How does an operator in Kenya tell an operator in Japan what end-site customers they have?
- And how does that operator send data packets to the customers of the Japanese operator, and get responses back
  - After all, as on a local ethernet, two way packet flow is needed for communication between two devices

- The operator in Kenya could buy a direct connection to the operator in Japan
  - But this doesn't scale there are thousands of distinct networks, would need thousands of connections, and cost would be astronomical
- Instead, the operator in Kenya tells his neighbouring operators what end-sites they have
  - And the neighbouring operators pass this information on to their neighbours, and so on
  - This process repeats until the information reaches the operator in Japan

- This process is called "Routing"
- The mechanisms used are called "Routing Protocols"
- Routing and Routing Protocols ensures that
  - The Internet can scale
  - Thousands of network operators can provide connectivity to each other
  - We have the Internet we see today

- The Network Operator in Kenya doesn't actually tell its neighbouring operators the names of the end-sites
  - (network equipment does not understand names)
- Instead, it has received an IP address block as a member of the Regional Internet Registry serving Kenya
  - Its customers have received address space from this address block as part of their "Internet service"
  - And it announces this address block to its neighbouring operators – this is called announcing a "route"

### Routing Protocols

- Routers use "routing protocols" to exchange routing information with each other
  - IGP is used to refer to the process running on routers inside an ISP's network
  - EGP is used to refer to the process running between routers bordering directly connected ISP networks

#### What Is an IGP?

- Interior Gateway Protocol
- Within an Autonomous System
- Carries information about internal infrastructure prefixes
- Two widely used IGPs:
  - OSPF
  - IS-IS

#### Why Do We Need an IGP?

- Network Operator backbone scaling
  - Hierarchy
  - Limiting scope of failure
  - Only used for operator's infrastructure addresses, not customers or anything else
  - Design goal is to minimise number of prefixes in IGP to aid scalability and rapid convergence

#### What Is an EGP?

- Exterior Gateway Protocol
- Used to convey routing information between Autonomous Systems
- De-coupled from the IGP
- □ Current EGP is BGP

### Why Do We Need an EGP?

- Scaling to large network
  - Hierarchy
  - Limit scope of failure
- Define Administrative Boundary
- Policy
  - Control reachability of prefixes
  - Merge separate organisations
  - Connect multiple IGPs

# Interior versus Exterior Routing Protocols

#### Interior

- Automatic neighbour discovery
- Generally trust your IGP routers
- Prefixes go to all IGP routers
- Binds routers in one AS together

#### Exterior

- Specifically configured peers
- Connecting with outside networks
- Set administrative boundaries
- Binds AS's together

# Interior versus Exterior Routing Protocols

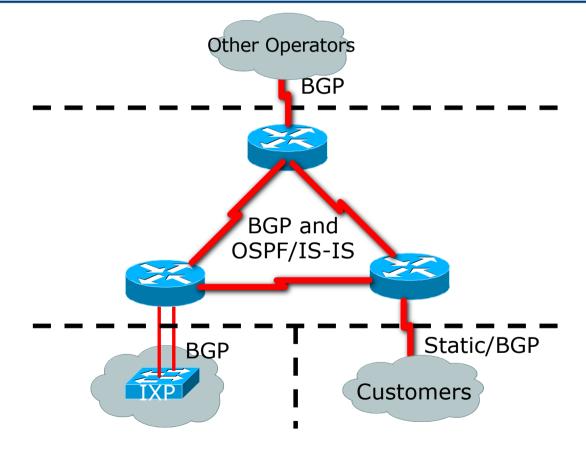
#### Interior

- Carries network infrastructure addresses only
- Network operators aim to keep the IGP small for efficiency and scalability

#### Exterior

- Carries customer prefixes
- Carries Internet prefixes
- EGPs are independent of the operator's network topology

## Hierarchy of Routing Protocols



## FYI: Default Administrative Distances

Route Source	Cisco	Juniper	Huawei	Brocade	Nokia/ALU	Mikrotik
<b>Connected Interface</b>	0	0	0	0	0	0
<b>Static Route</b>	1	5	60	1	1	1
EIGRP Summary Route	5	N/A	?	N/A	N/A	N/A
External BGP	20	170	255	20	170	20
Internal EIGRP Route	90	N/A	?	N/A	N/A	N/A
IGRP	100	N/A	?	N/A	N/A	N/A
OSPF	110	10	10	110	10	110
IS-IS	115	18	15	115	18	N/A
RIP	120	100	100	120	100	120
EGP	140	N/A	N/A	N/A	N/A	N/A
External EIGRP	170	N/A	?	N/A	N/A	N/A
Internal BGP	200	170	255	200	130	200
Unknown	255	255	?	255	?	

## Routing Basics

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